

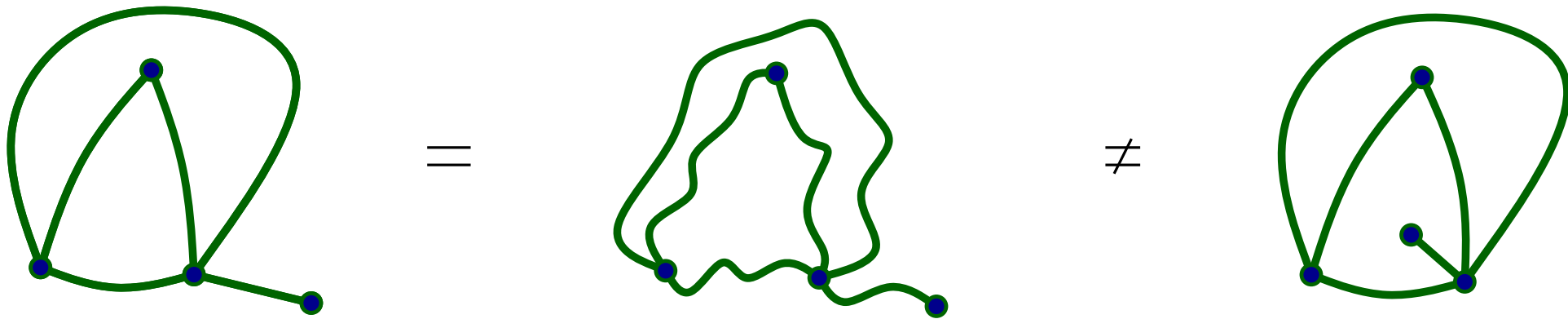
Bipolar orientations and blossoming trees

Marie Albenque (CNRS, Paris)
Joint work with Dominique Poulalhon

Kolkom, November 17th 2012

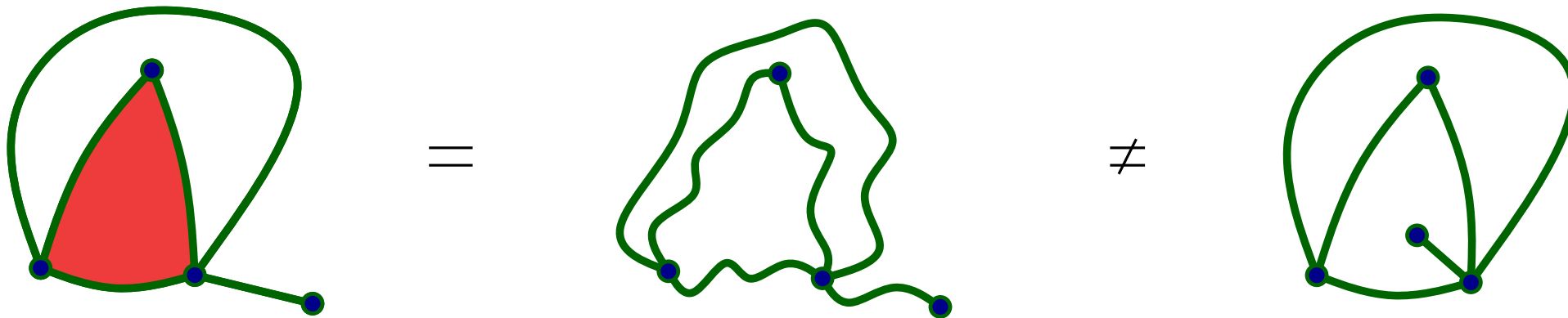
Plane Maps.

A **plane map** is the embedding of a connected graph in the plane up to continuous deformations.



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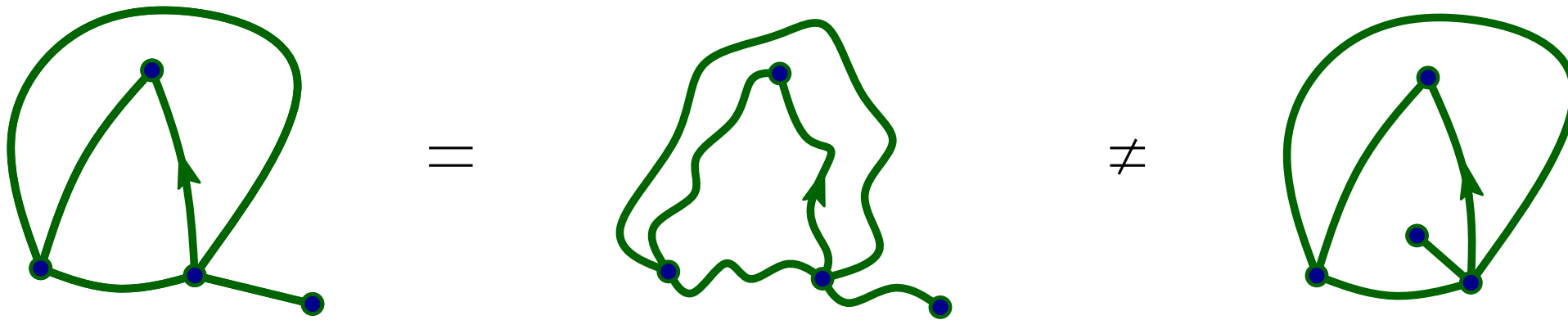
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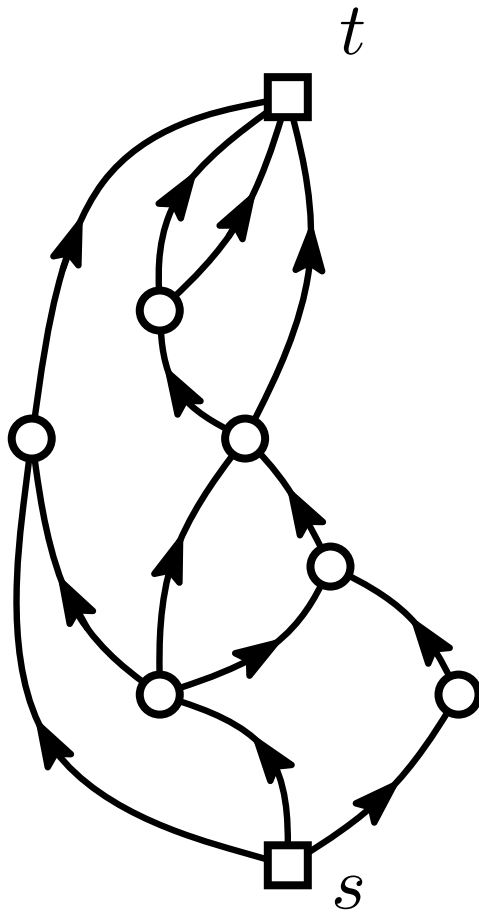
Plane maps are **rooted**.

There is one special face which is infinite: the **outer face**.

Plane Bipolar Orientations

A **plane bipolar orientation** is a plane map:

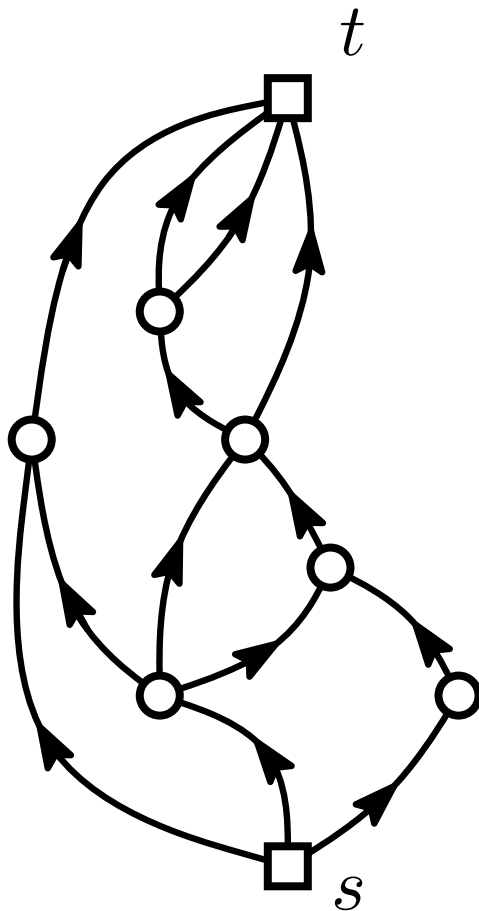
- endowed with an acyclic orientation,
- with a unique **source** vertex (without ingoing edges),
- with a unique **sink** vertex (without outgoing edges).



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- endowed with an acyclic orientation,
- with a unique **source** vertex (without ingoing edges),
- with a unique **sink** vertex (without outgoing edges).



Theorem : The number Θ_{ij} of bipolar orientations with $i + 2$ vertices and $j + 1$ faces is equal to:

$$\Theta_{ij} = \frac{2(i + j)!(i + j + 1)!(i + j + 2)!}{i!(i + 1)!(i + 2)!j!(j + 1)!(j + 2)!}.$$

[Baxter '01]

[Fusy, Poulalhon, Schaeffer '09]

[Bonichon, Bousquet-Mélou, Fusy '10]

[Felsner, Fusy, Noy, Orden '11]

Enumeration

One of the main question when studying some families of maps:

How many maps belong to this family ?

- **Recursive decomposition:** Tutte '60s. Baxter '01
- **Matrix integrals:** t'Hooft '74, Brézin, Itzykson, Parisi and Zuber '78.
- **Representation of the symmetric group:** Goulden and Jackson '87.
- **Bijective approach with labeled trees:** Cori-Vauquelin '81, Schaeffer '98, Bouttier, Di Francesco and Guitter '04, Bernardi, Chapuy, Fusy, Miermont, ...
- **Bijective approach with blossoming trees:** Schaeffer '98, Schaeffer and Bousquet-Mélou '00, Poulalhon and Schaeffer '05, Fusy, Poulalhon and Schaeffer '06.

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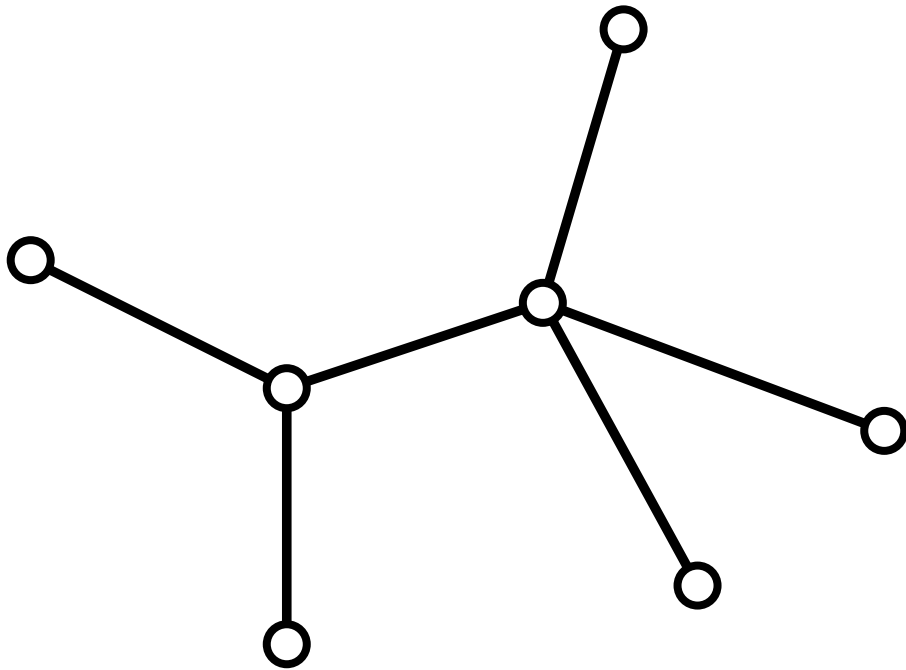
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What is a blossoming tree ?

A **blossoming tree** is a plane tree where vertices can carry **opening stems** or **closing stems**, such that :

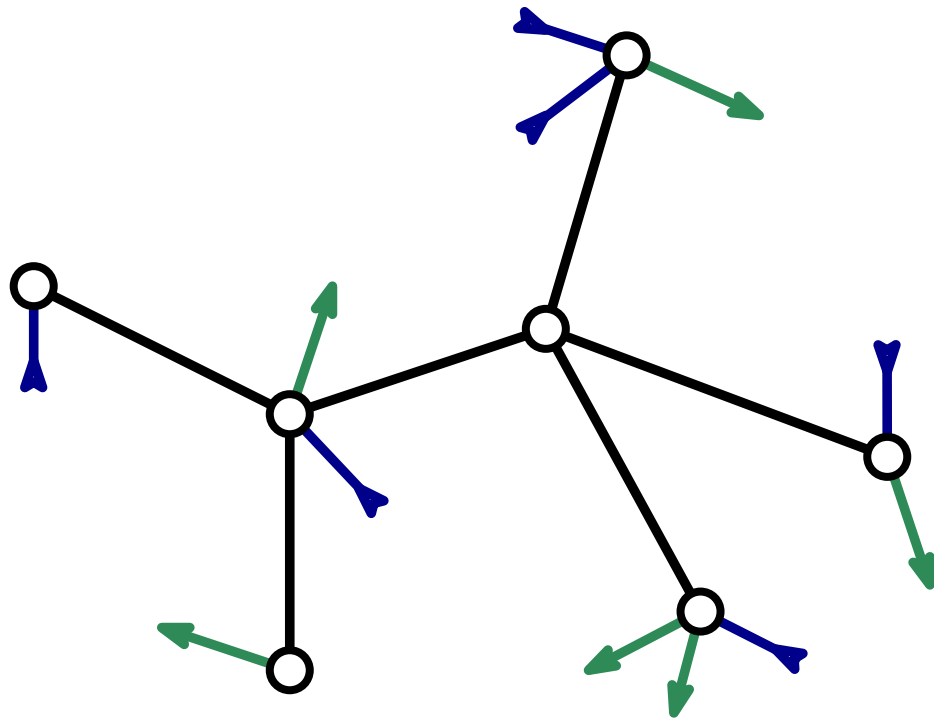
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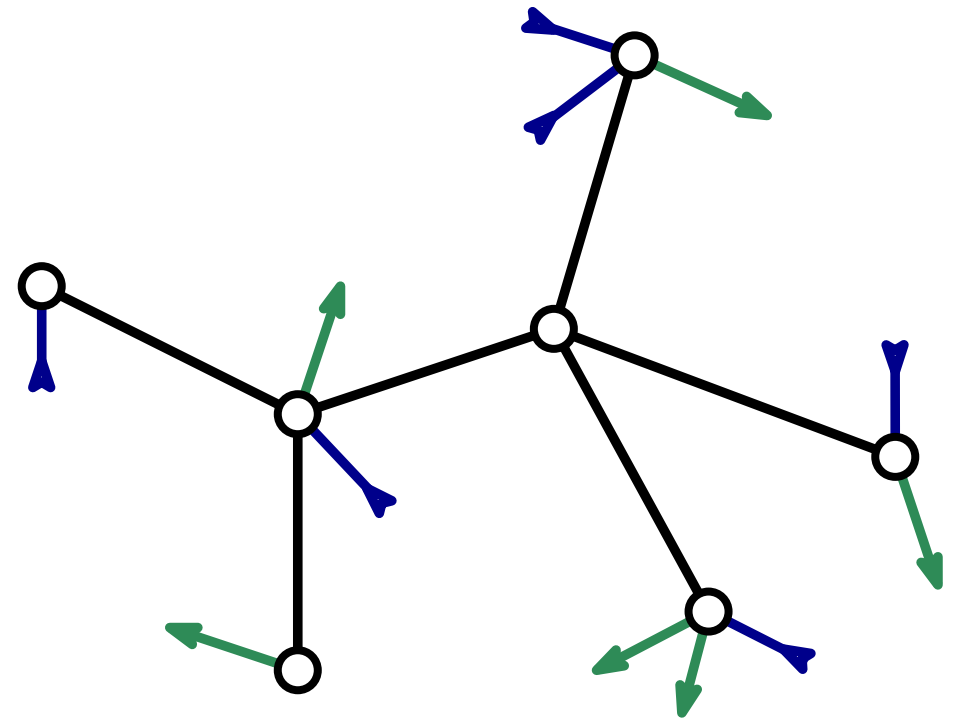
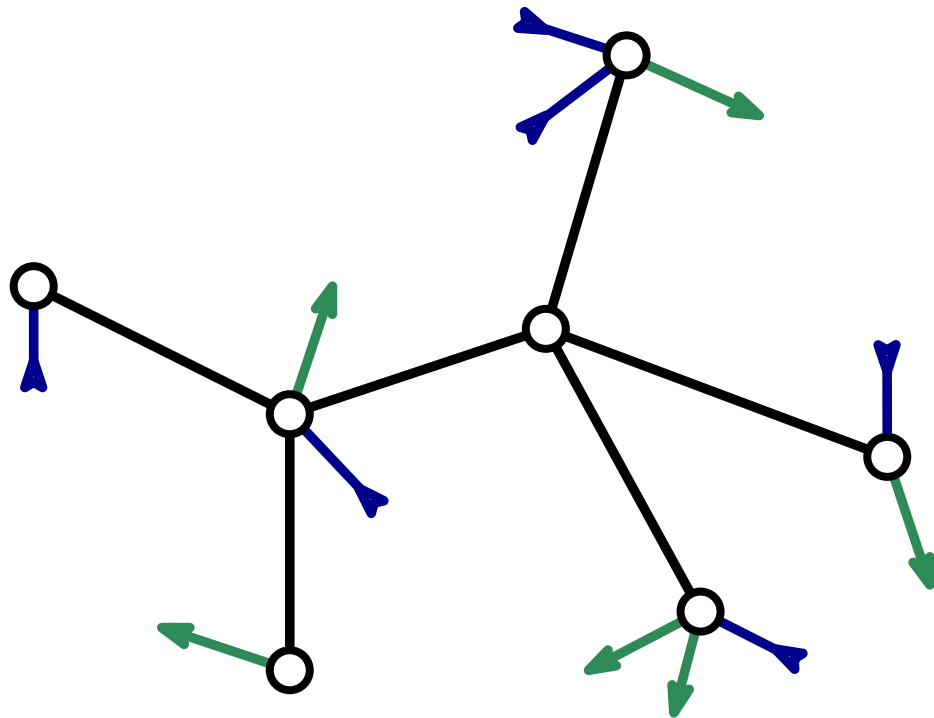
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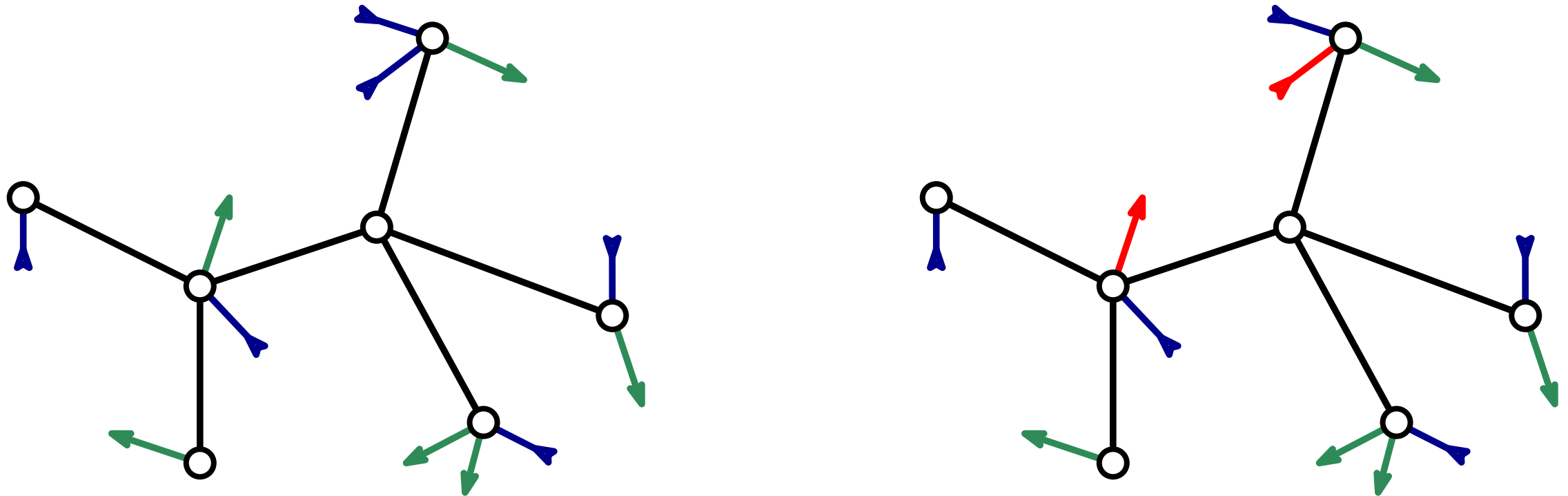
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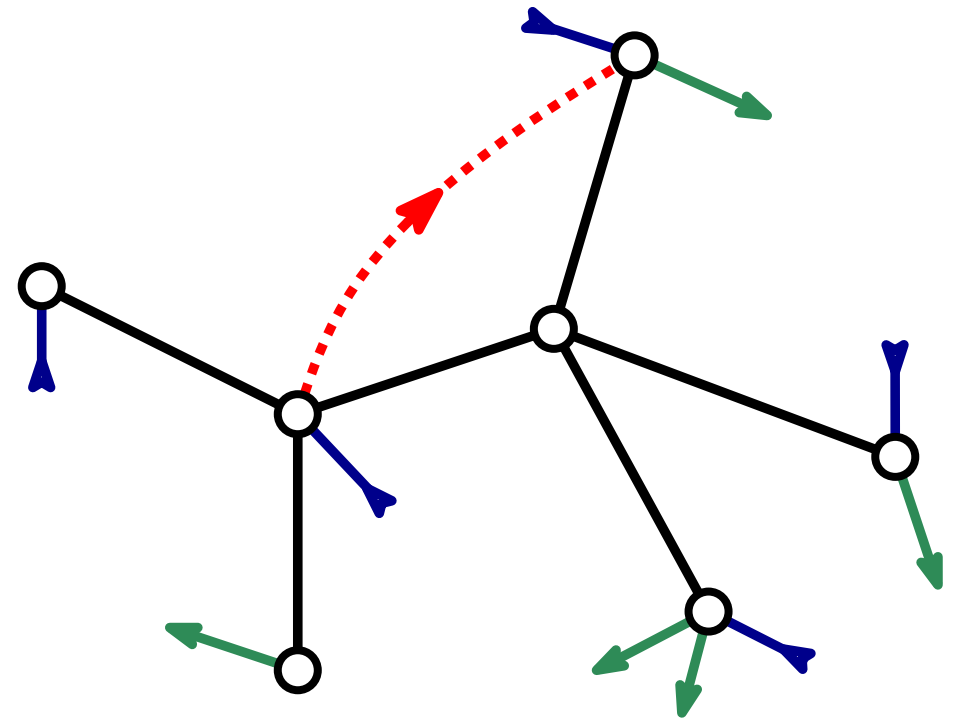
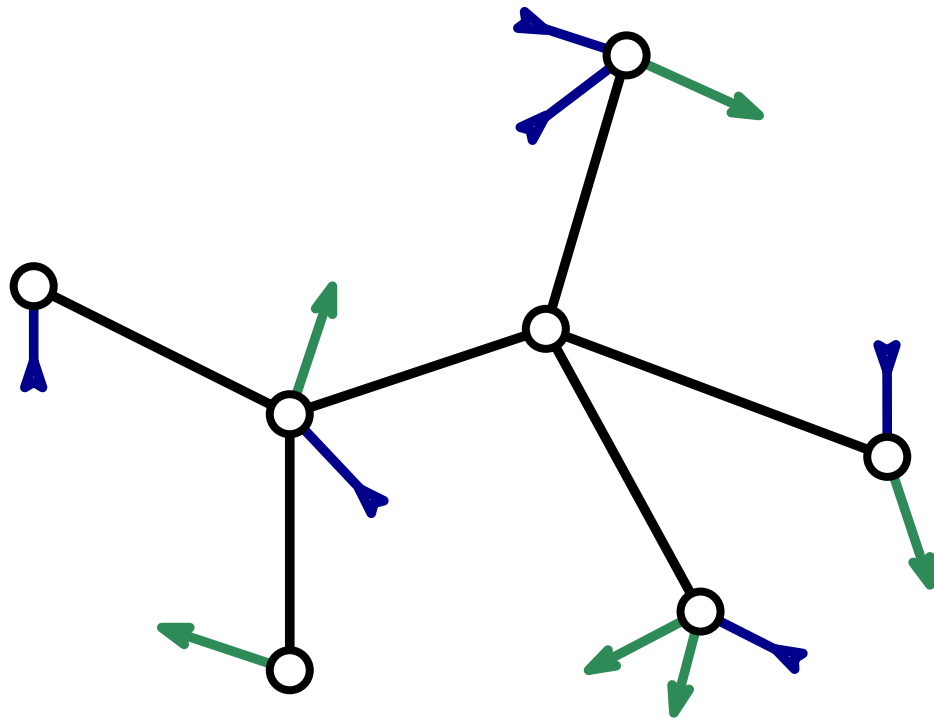
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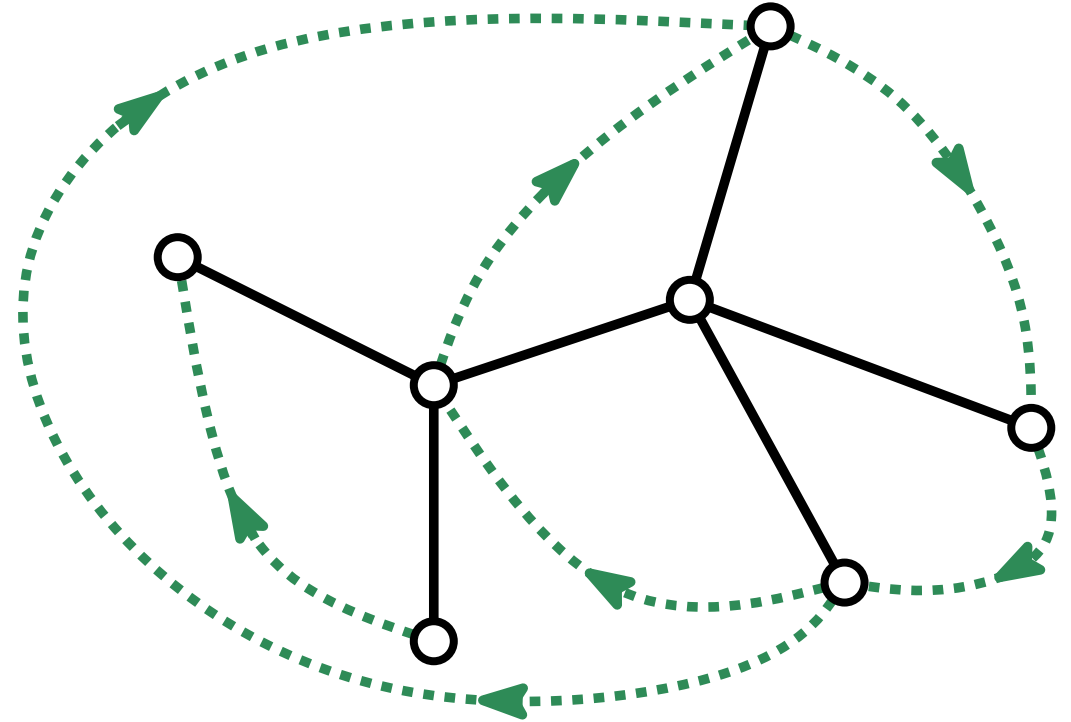
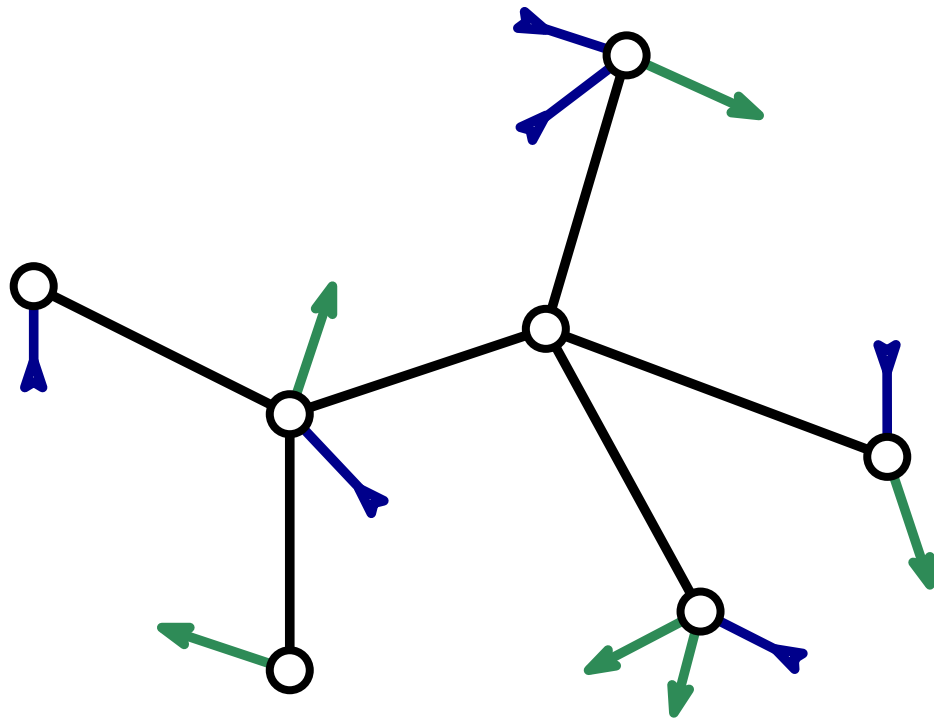
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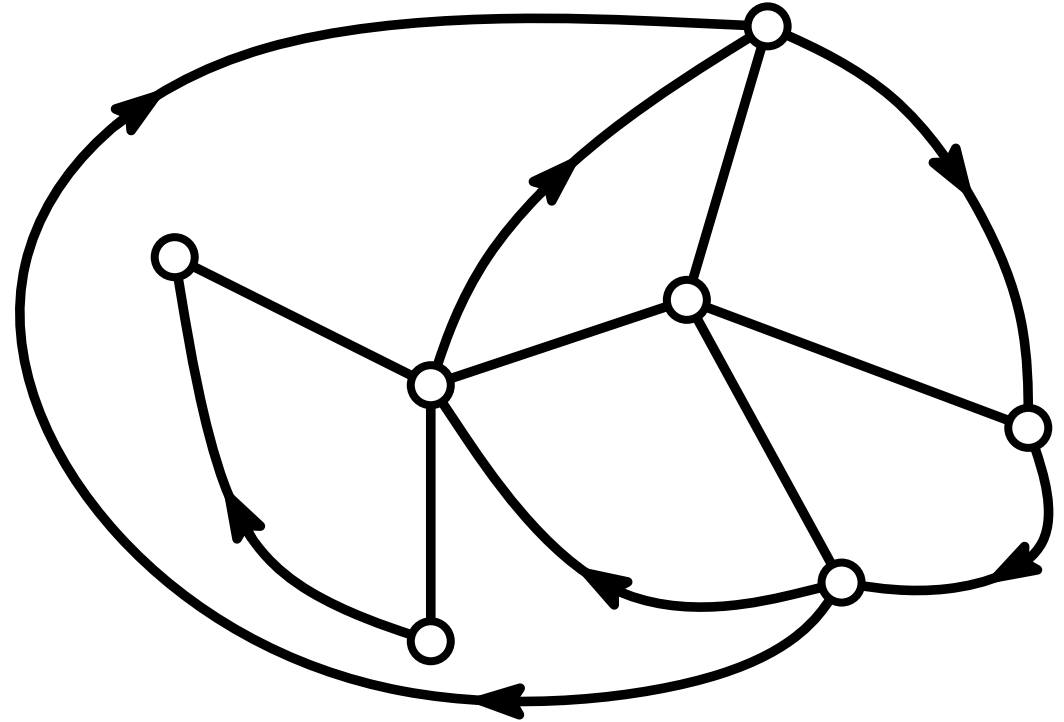
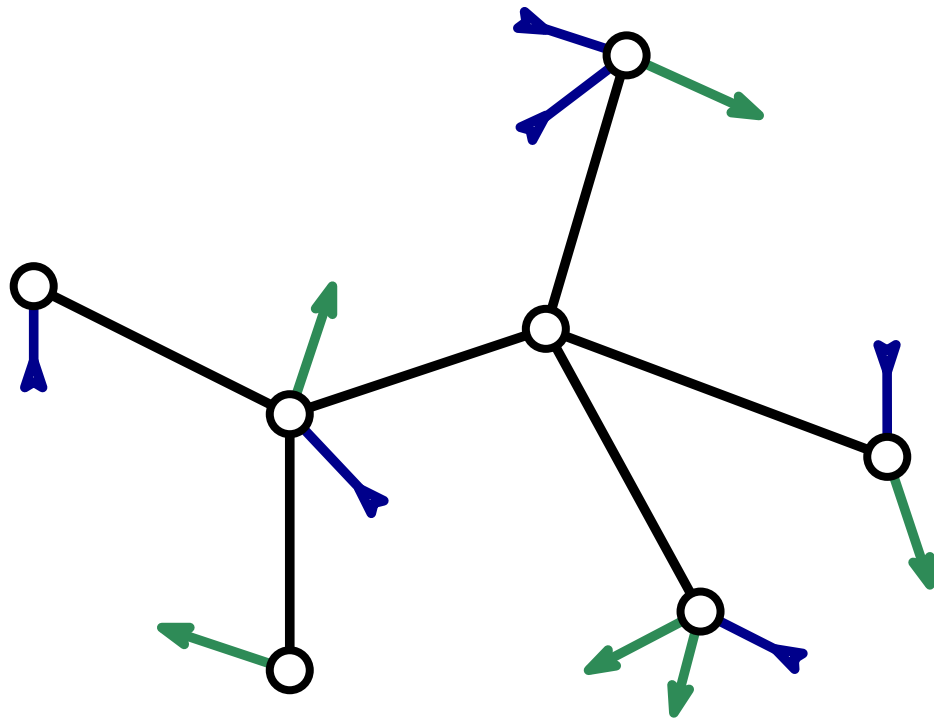
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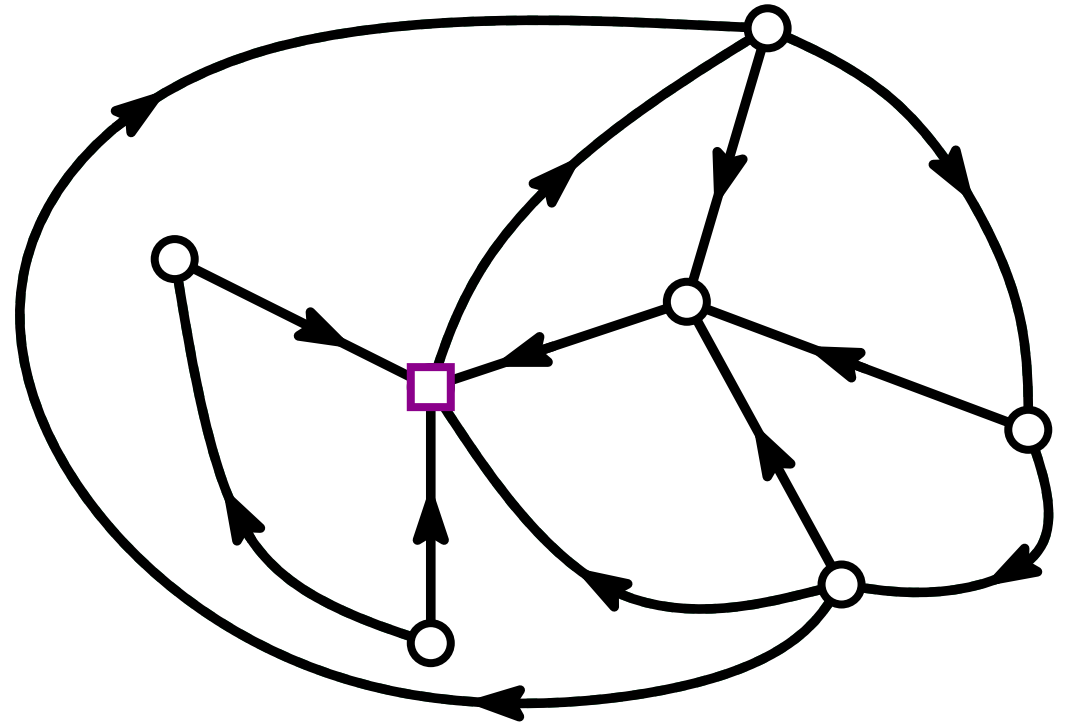
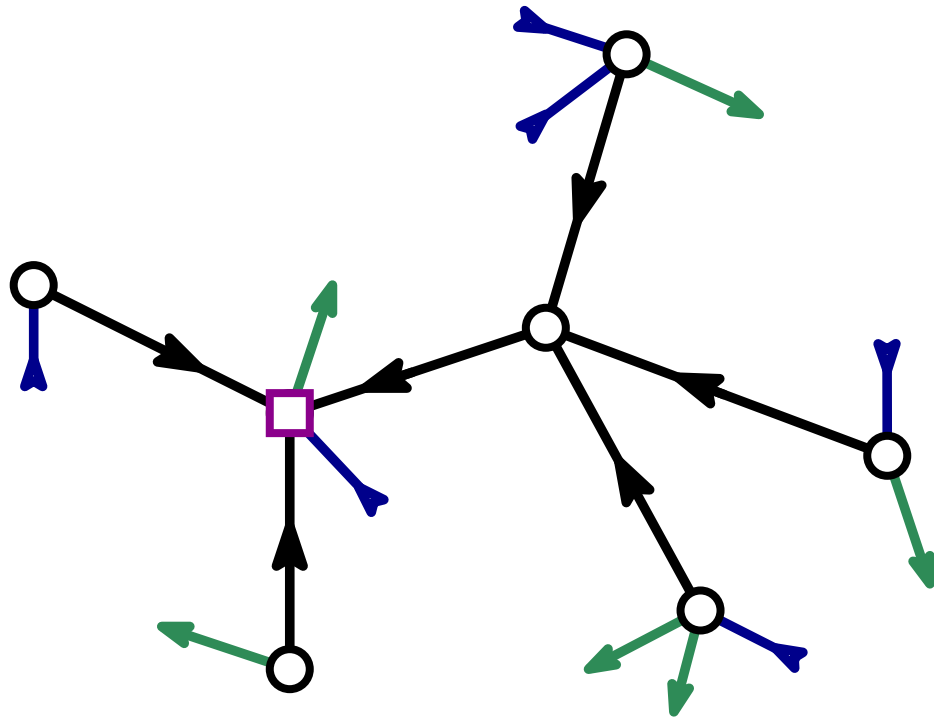
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What is a blossoming tree ?



A plane map can be canonically associated to any blossoming tree by making all closures clockwise.

If the tree is rooted and its edges oriented towards the root + closure edges oriented naturally

⇒ Accessible orientation of the map without ccw cycles.

Can we always find a blossoming tree from a plane map ?

Theorem : [A., Poulalhon]

If a plane map M has a marked vertex v is endowed with an orientation such that :

- there exists a directed path from any vertex to v ,
- there is no counterclockwise cycle,

then there exists a **unique** blossoming tree rooted at v whose closure is M endowed with the same orientation.

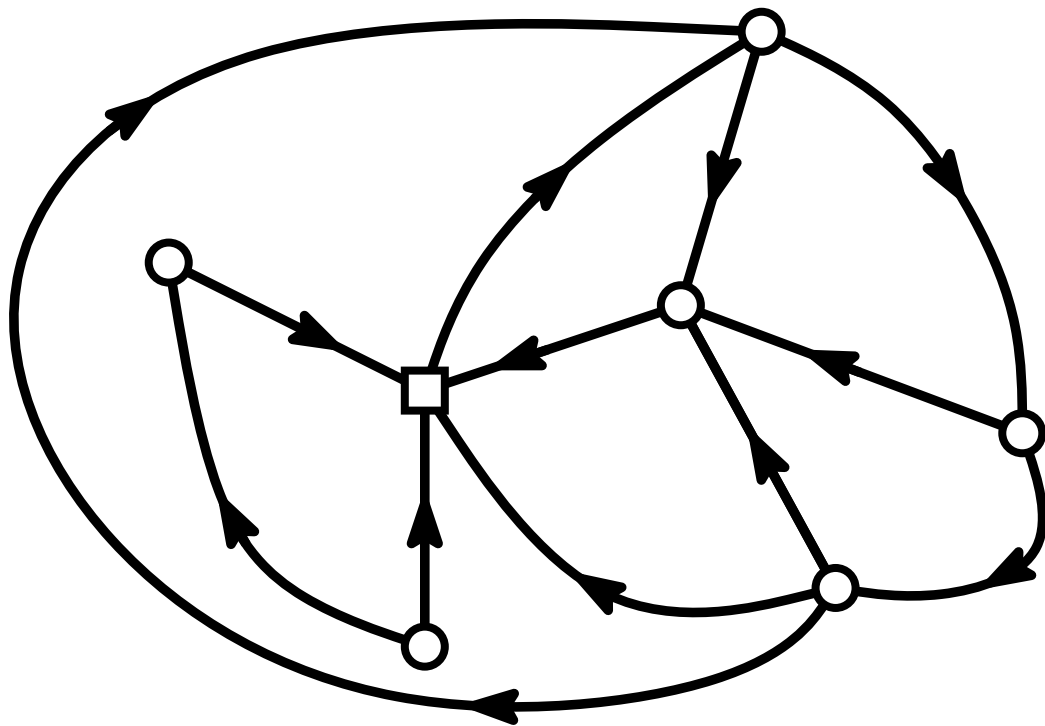
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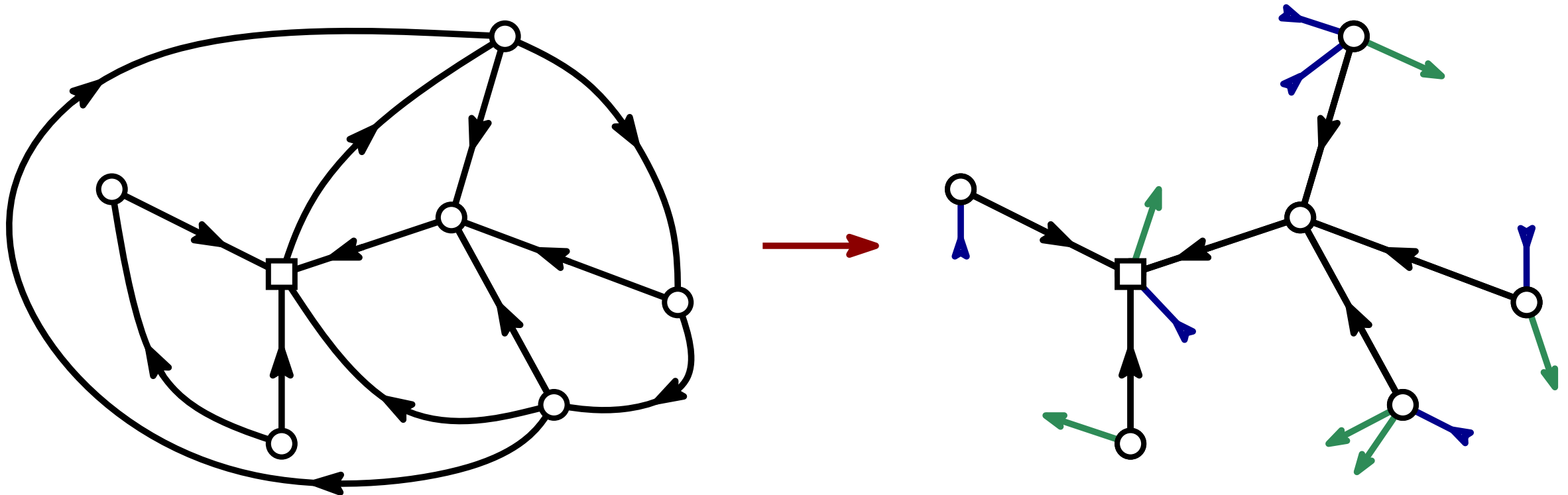
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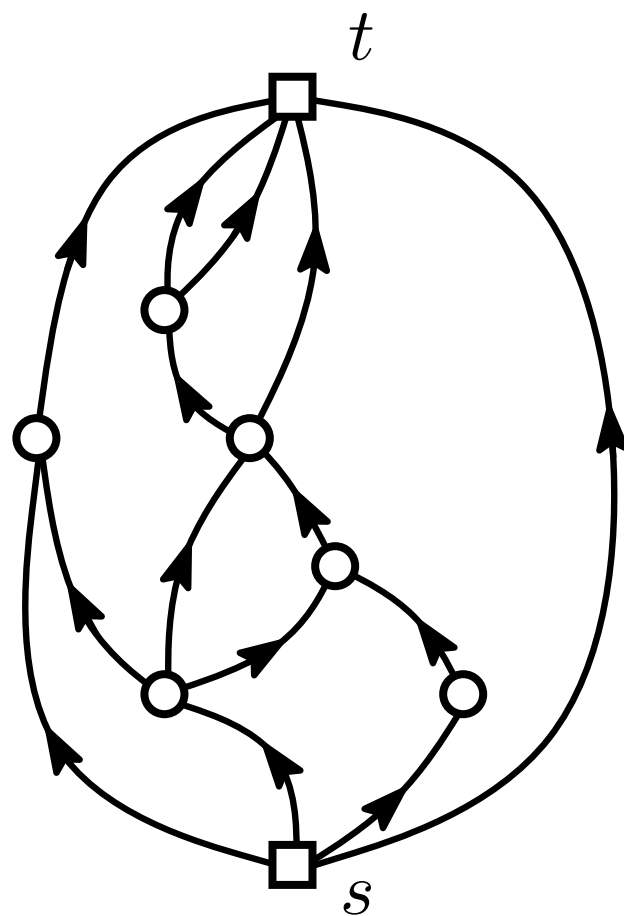
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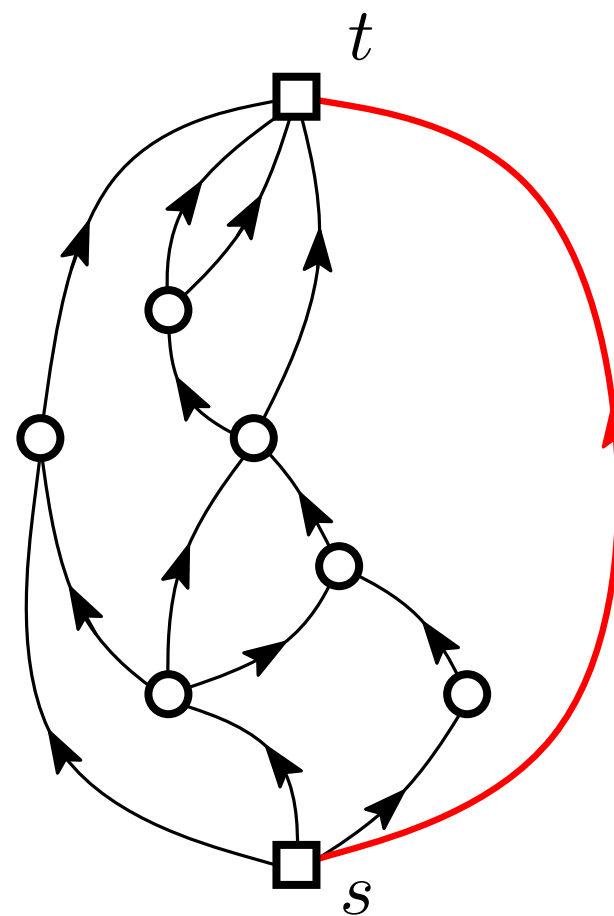
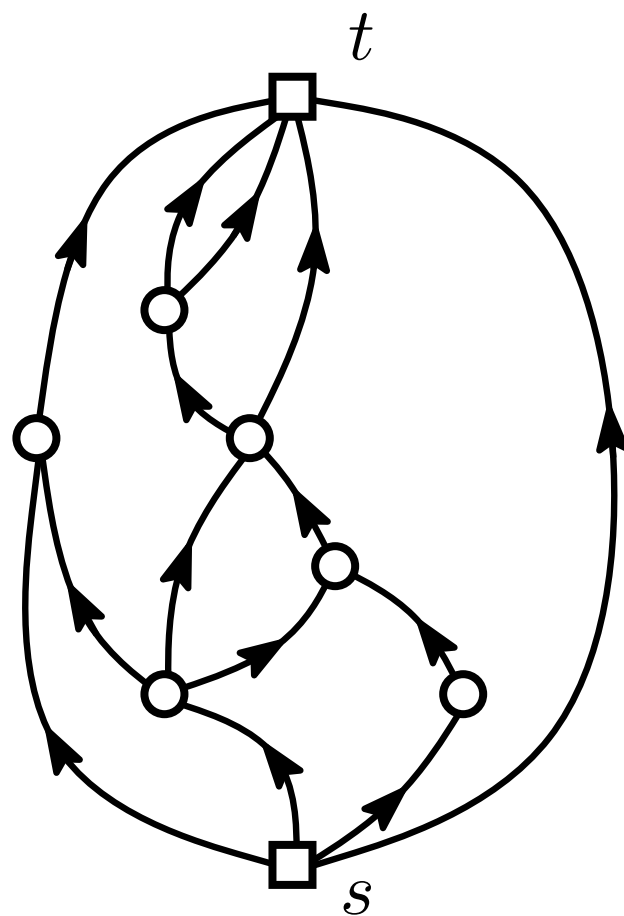
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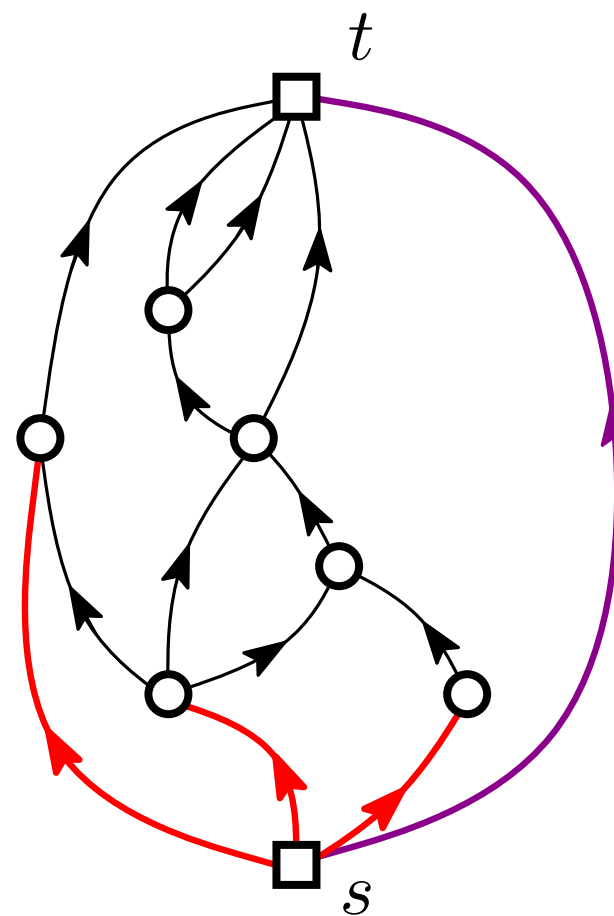
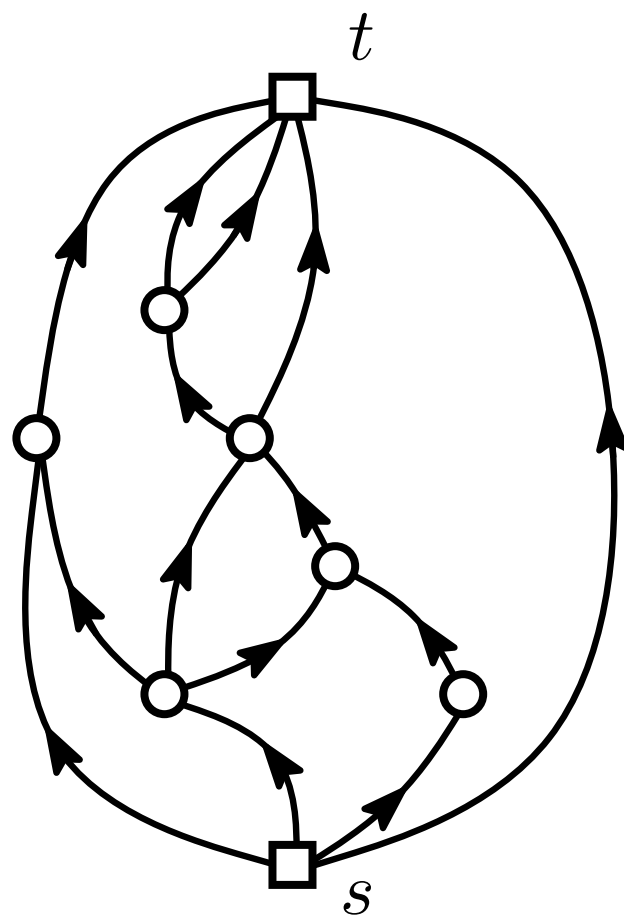
Blossoming trees and bipolar orientations



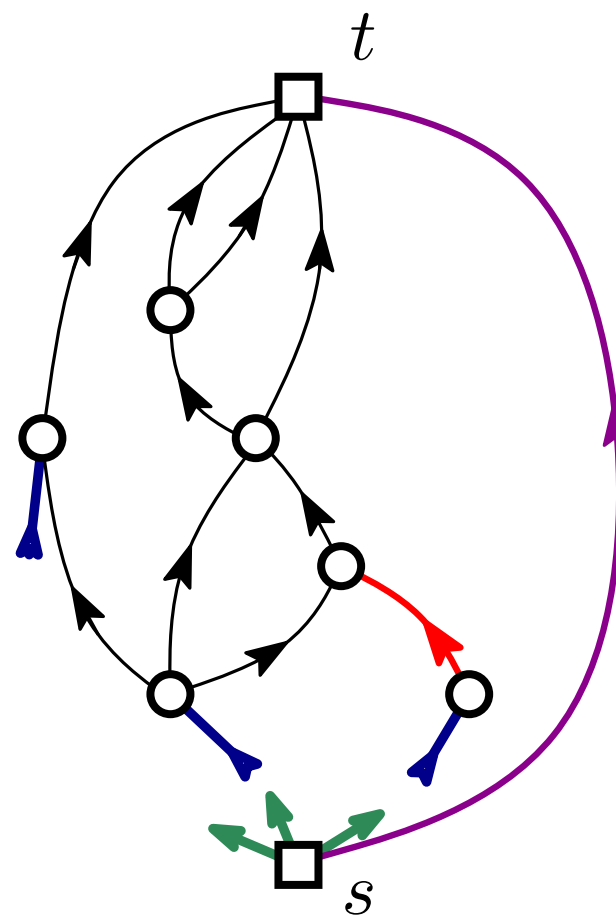
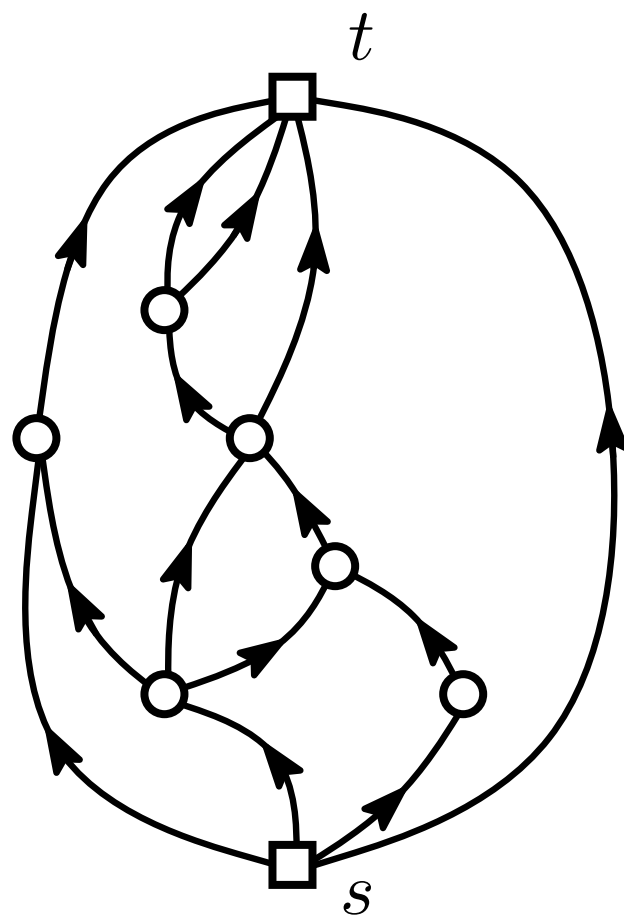
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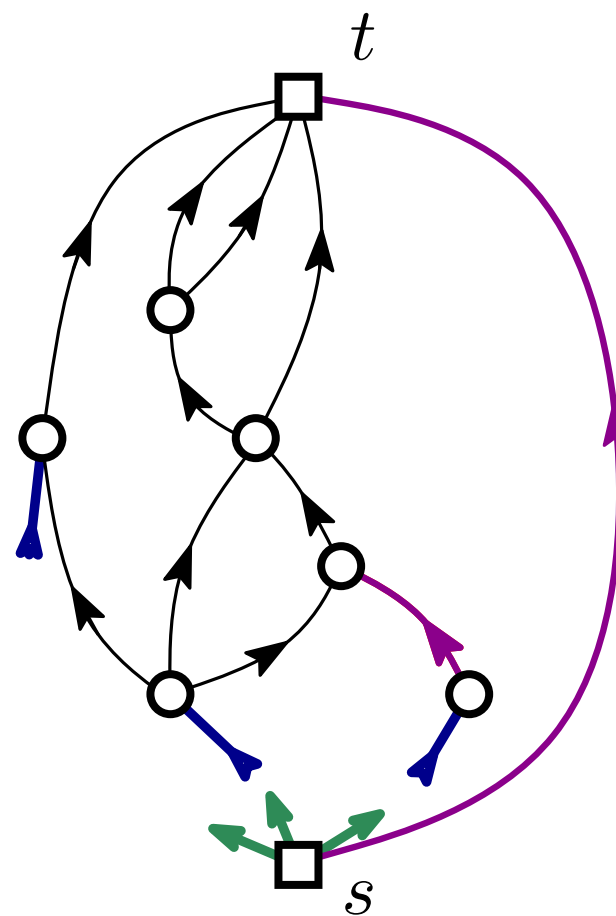
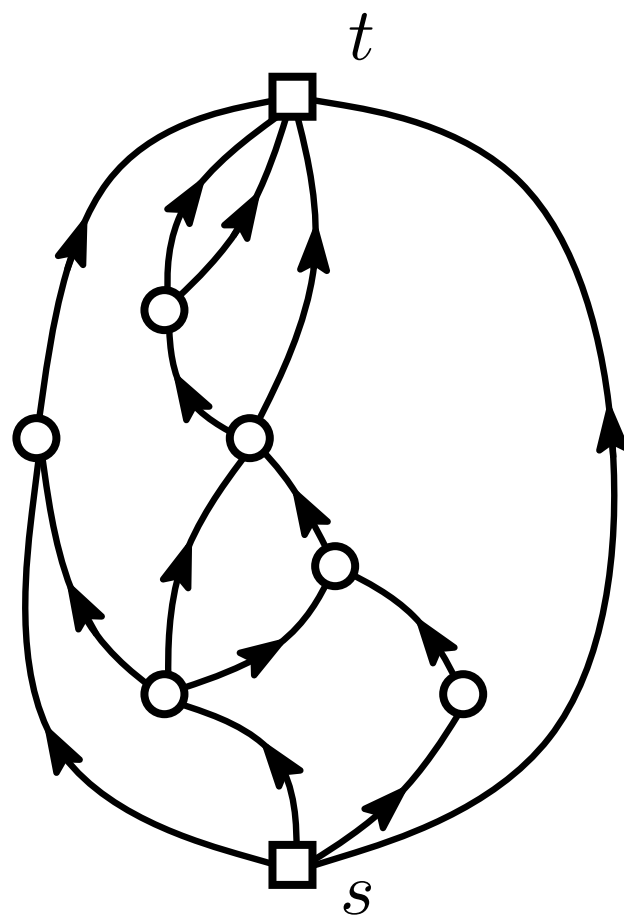
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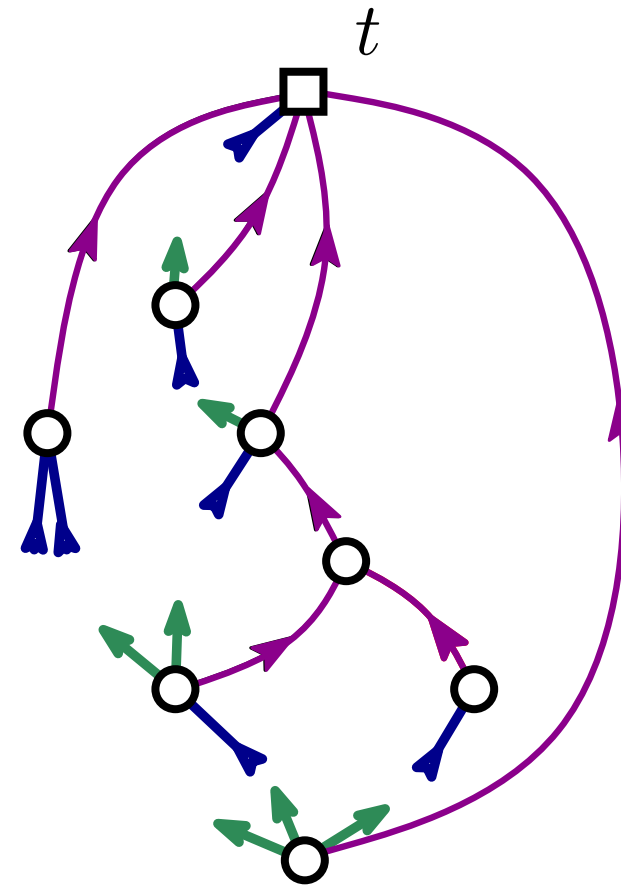
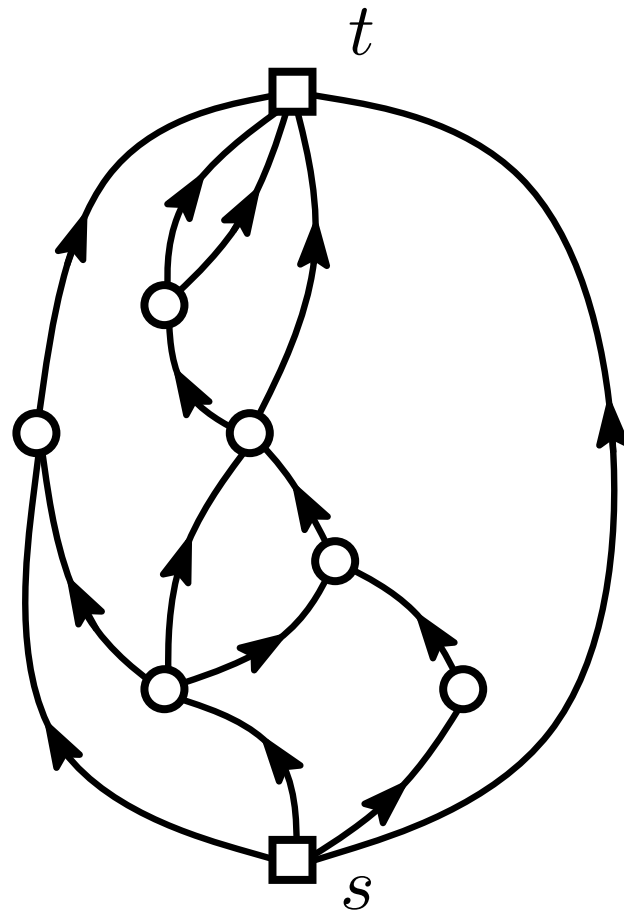
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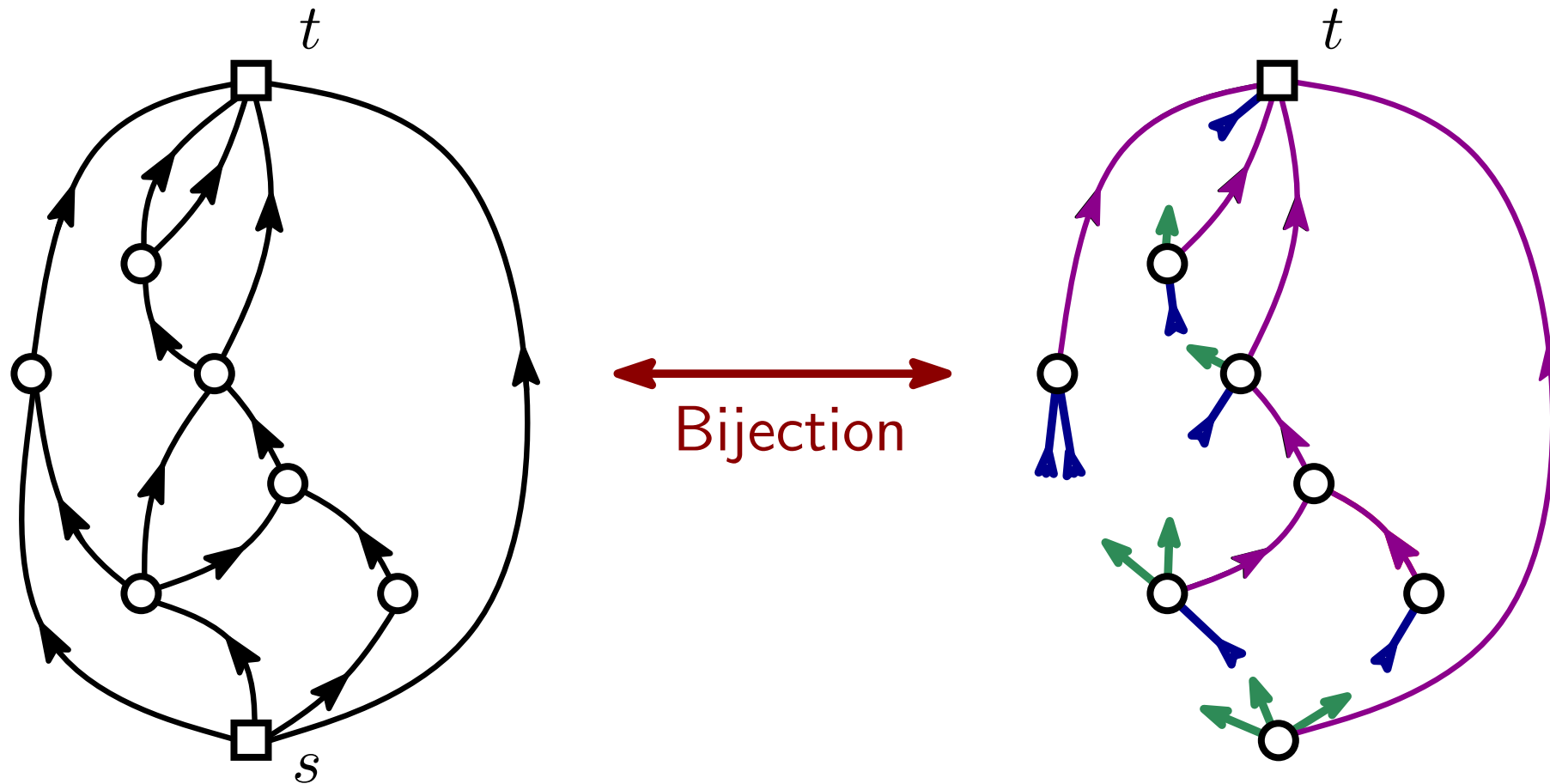
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marked vertex \in outer face \Rightarrow easy to compute the blossoming tree

[Bernardi '07]

Blossoming trees and bipolar orientations



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Description/enumeration of these trees ?

Blossoming trees and triplet of paths

$T_{\text{bip}}(i, j)$ = blossoming trees obtained after opening a bipolar orientation with $i + 2$ vertices and $j + 1$ faces

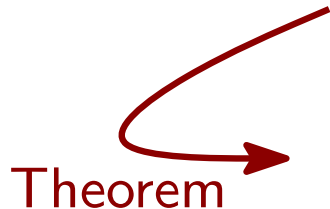
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 **Theorem** = blossoming trees which closes into a bipolar orientation with $i + 2$ vertices and $j + 1$ faces

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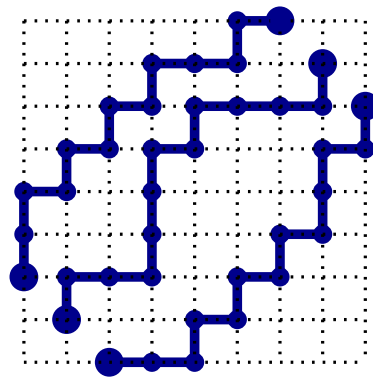
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Proposition: [A., Poulalhon]

There exists a one-to-one correspondence between :

$T_{\text{bip}}(i, j)$ and



triplet of non-intersecting paths
with i  and j 
and fixed first and final points

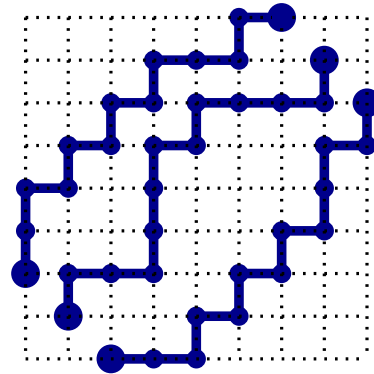
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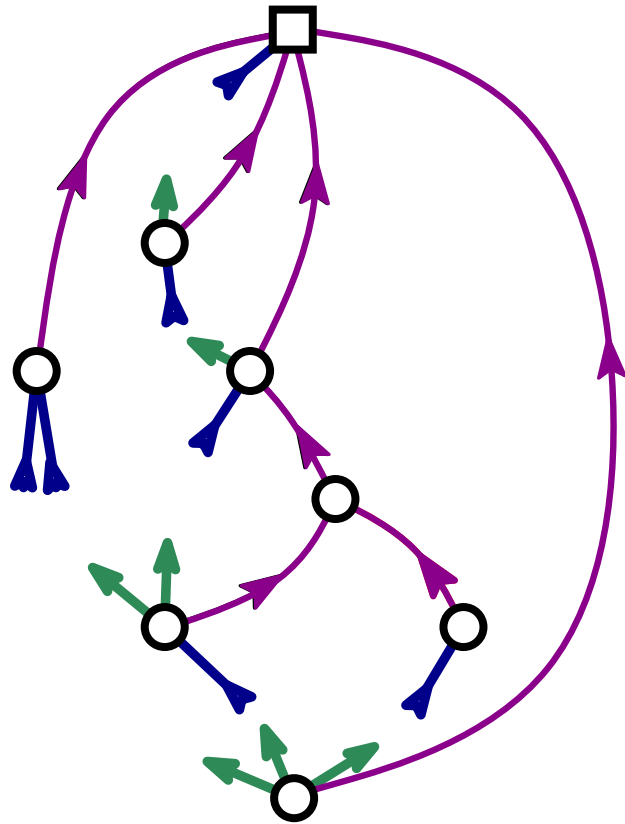
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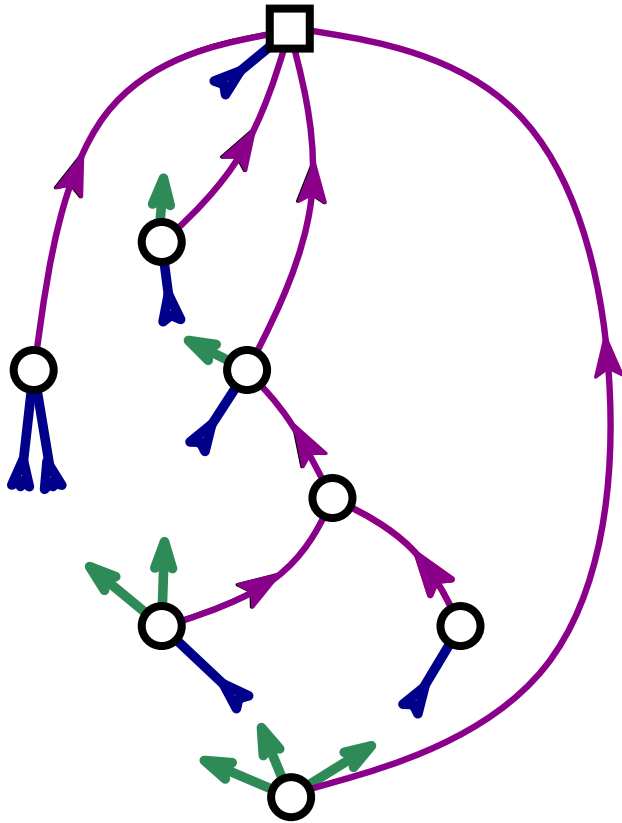
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Trees of T_{bip}

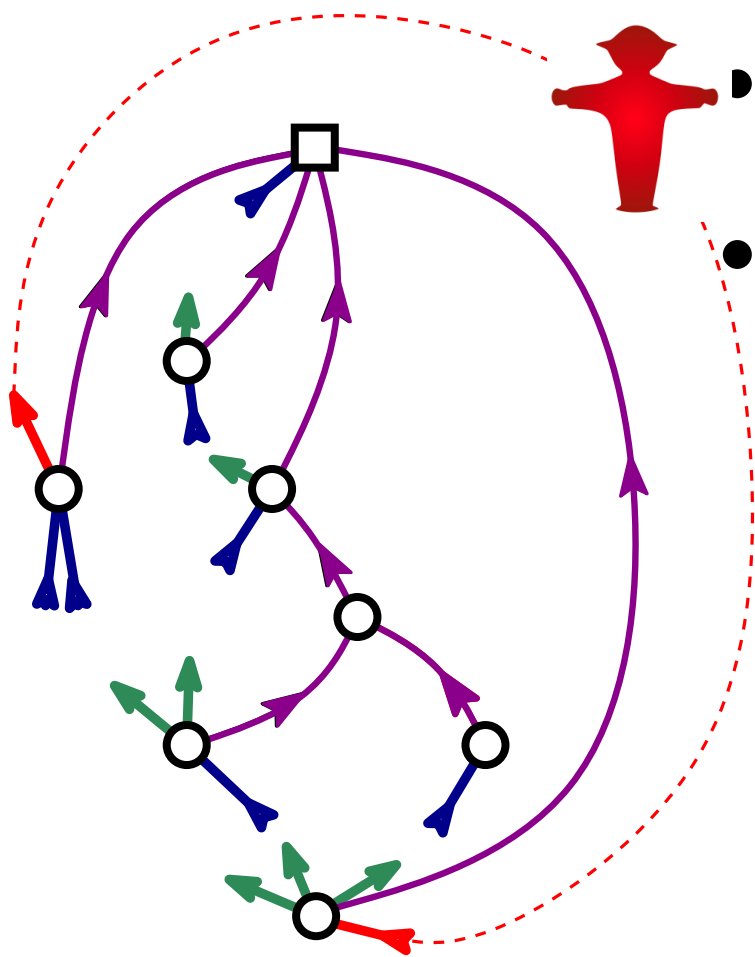


Trees of T_{bip}

- First son of the root = only opening stems

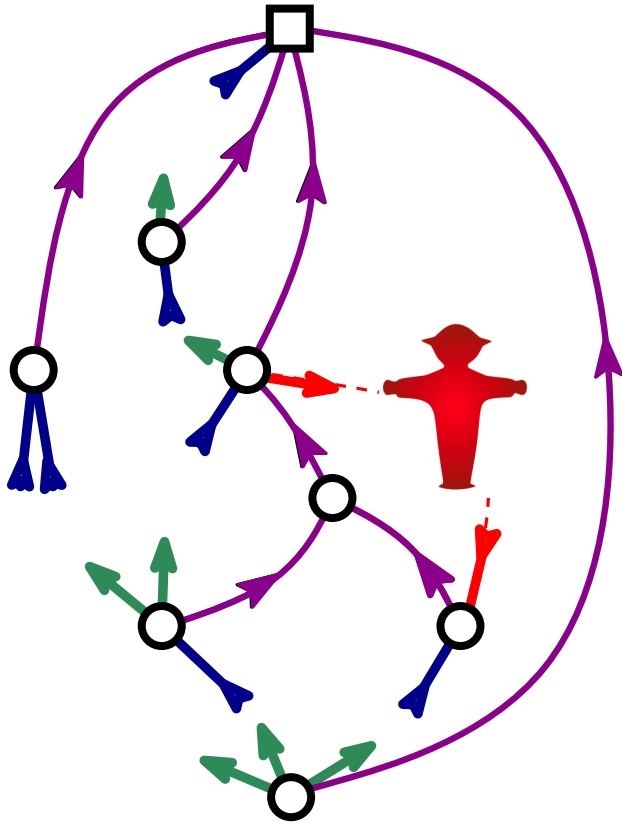


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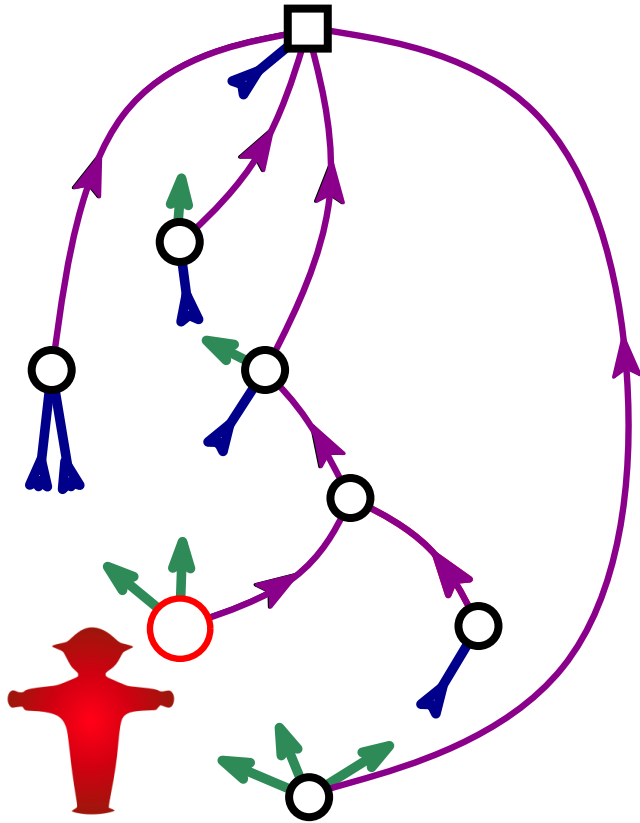
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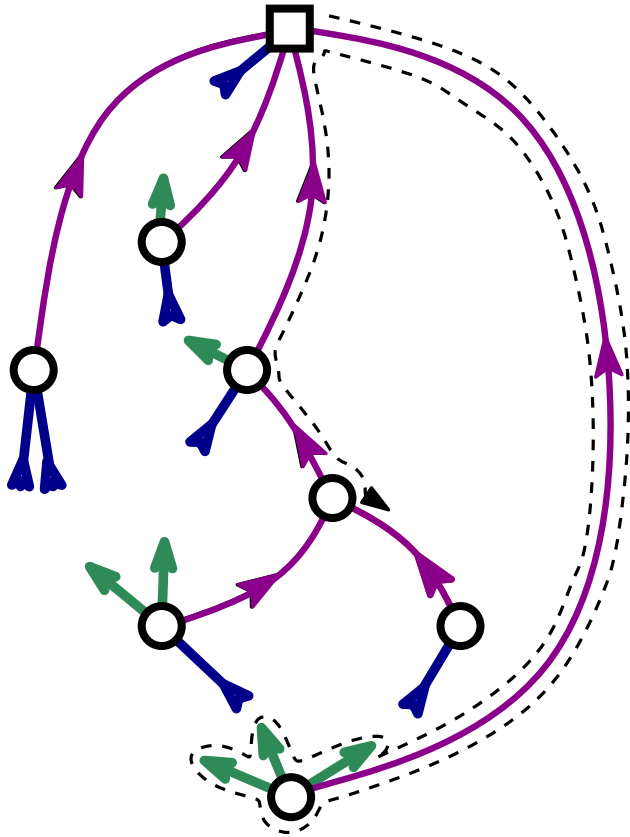
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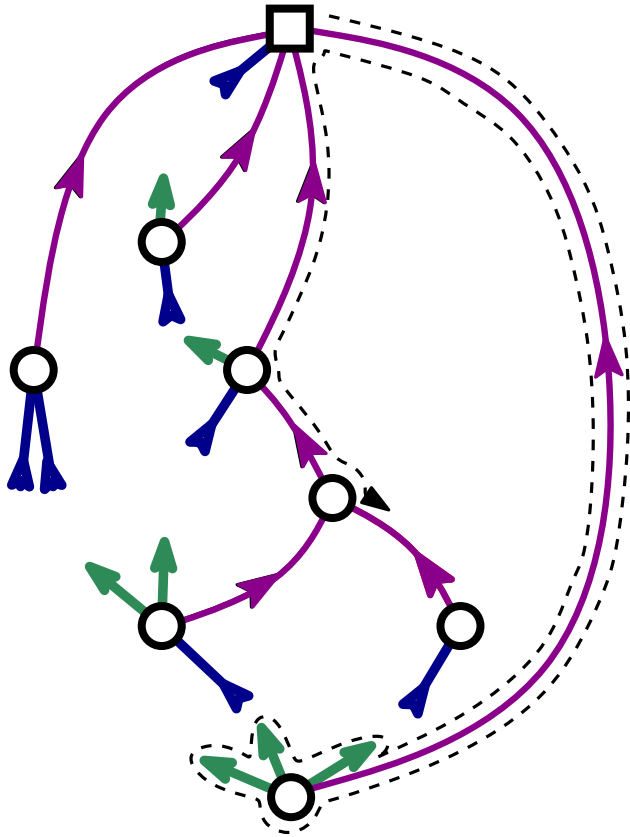
Encoding of the blossoming tree = **contour word** = word on $\{e, \bar{e}, b, \bar{b}\}$ s.t.:

e, \bar{e} : first time, second time we see an edge

b, \bar{b} : opening stem, closing stem.

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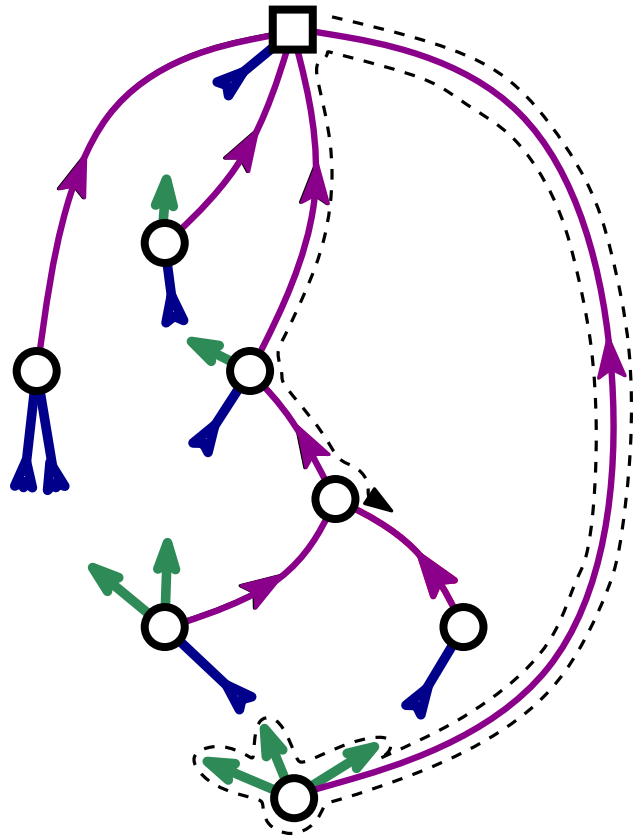
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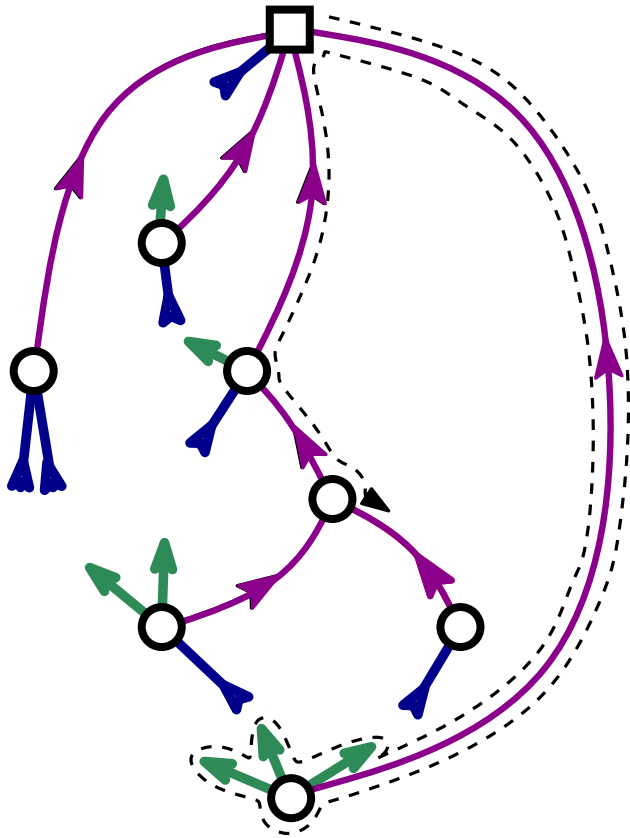
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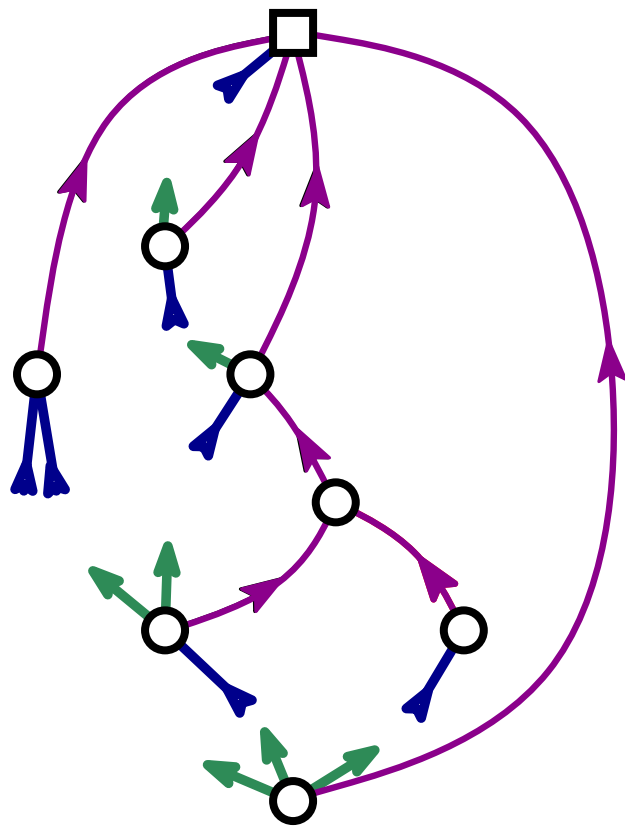
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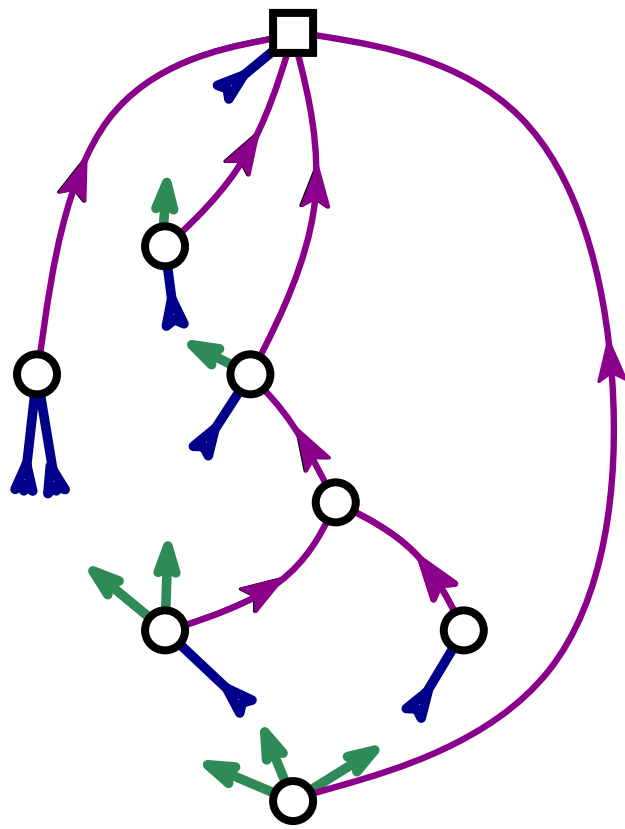
Trees of T_{bip} and triple of paths



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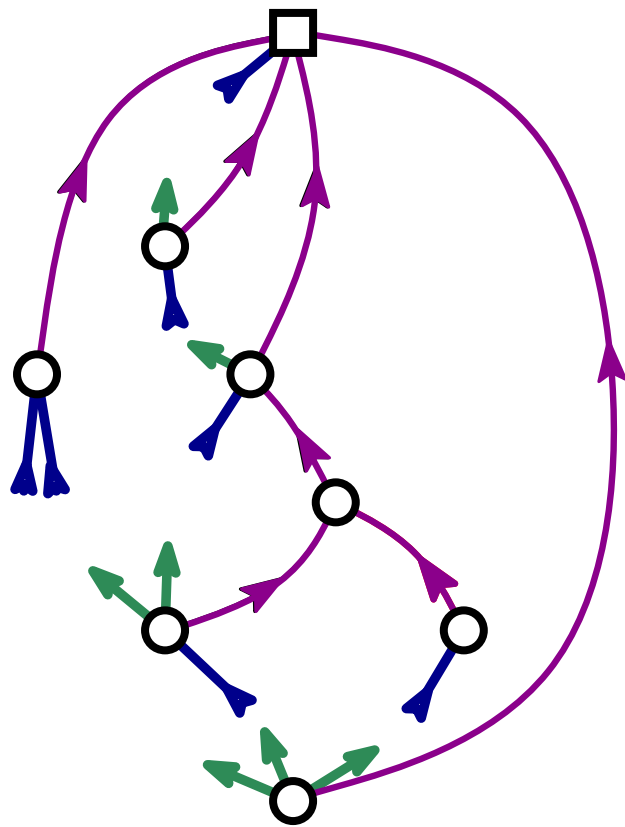


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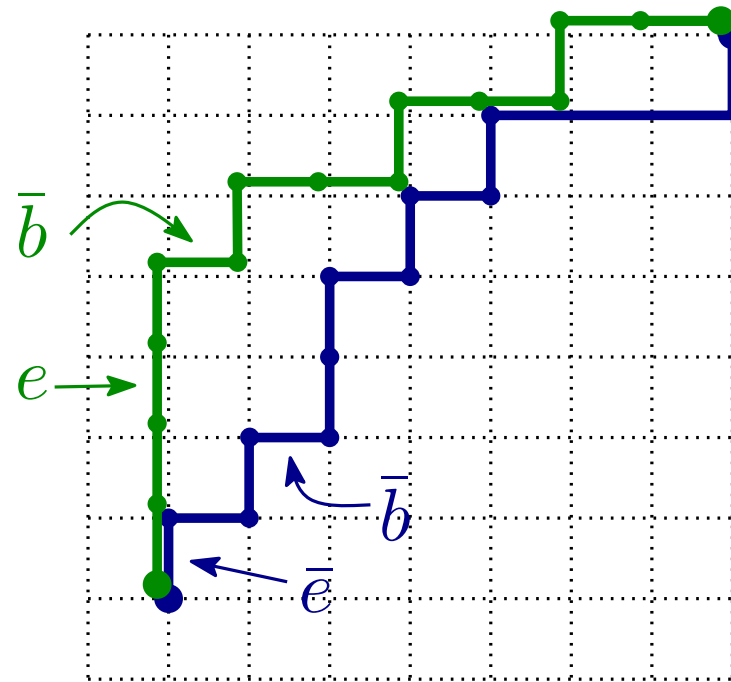


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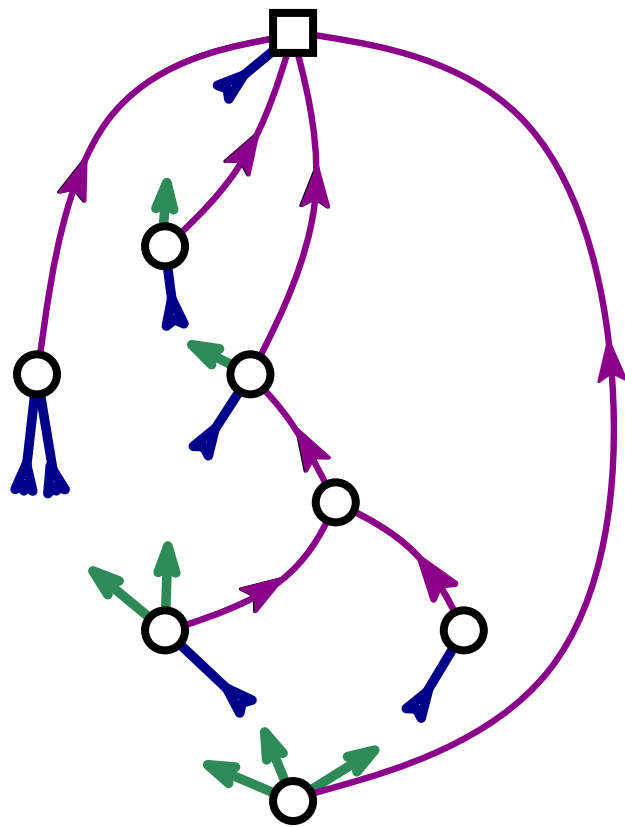
$$w_1 = e e e e \bar{b} e \bar{b} \bar{b} e \bar{b} \bar{b} e \bar{b} \bar{b}$$

$$w_2 = \bar{e} \bar{b} \bar{e} \bar{b} \bar{e} \bar{e} \bar{b} \bar{e} \bar{b} \bar{e} \bar{b} \bar{b} \bar{b} \bar{e}$$



$$w = e b b b \bar{e} e e e \bar{b} \bar{e} \bar{e} \bar{b} \bar{b} \bar{b} \bar{e} \bar{e} \bar{b} \bar{b} \bar{e} e \bar{b} \bar{b} \bar{e} \bar{b} e \bar{b} \bar{b} \bar{e}$$

Trees of T_{bip} and triple of paths

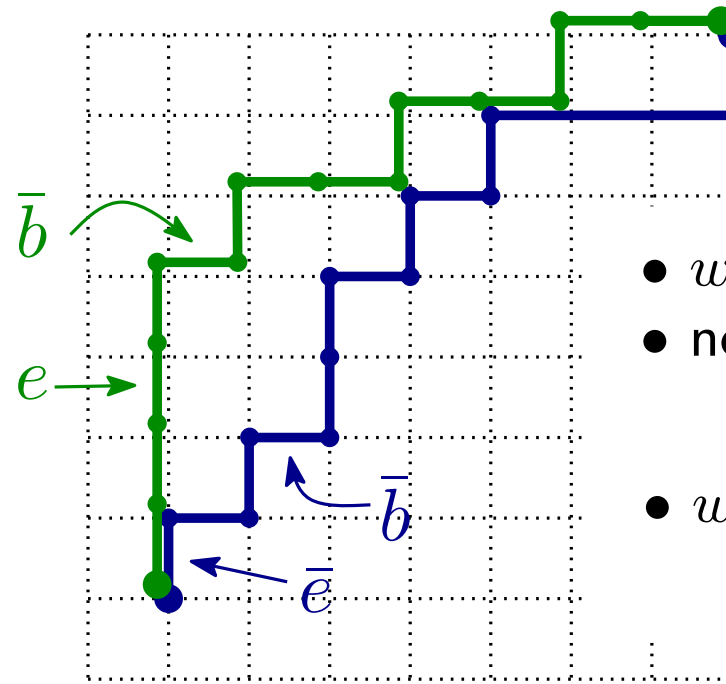


$w = eb$, $w[e, \bar{e}]$ and $w[b, \bar{b}] = \text{Dyck words}$,
 $w = \dots b \dots \bar{e} \dots \bar{b} \dots$, $w = e \dots e \dots \bar{b} \dots \bar{e} \dots$

$w_1 := w[e, \bar{b}]$, $w_2 := w[\bar{e}, \bar{b}]$

$w_1 = e e e e \bar{b} e \bar{b} \bar{b} e \bar{b} \bar{b} e \bar{b} \bar{b}$

$w_2 = \bar{e} \bar{b} \bar{e} \bar{b} \bar{e} \bar{e} \bar{b} \bar{e} \bar{b} \bar{e} \bar{b} \bar{b} \bar{b} \bar{e}$



- $w_2 =$ on the right/bottom of w_1
- no vertical edge in common.

+

- $w_1 = e \dots \bar{b}$ $w_2 = \bar{e} \bar{b}$

$w = e b b b \bar{e} e e e \bar{b} \bar{e} \bar{e} \bar{b} \bar{b} \bar{b} \bar{e} \bar{e} \bar{b} \bar{b} \bar{e} e \bar{b} \bar{b} \bar{e} \bar{b} e \bar{b} \bar{b} \bar{e}$

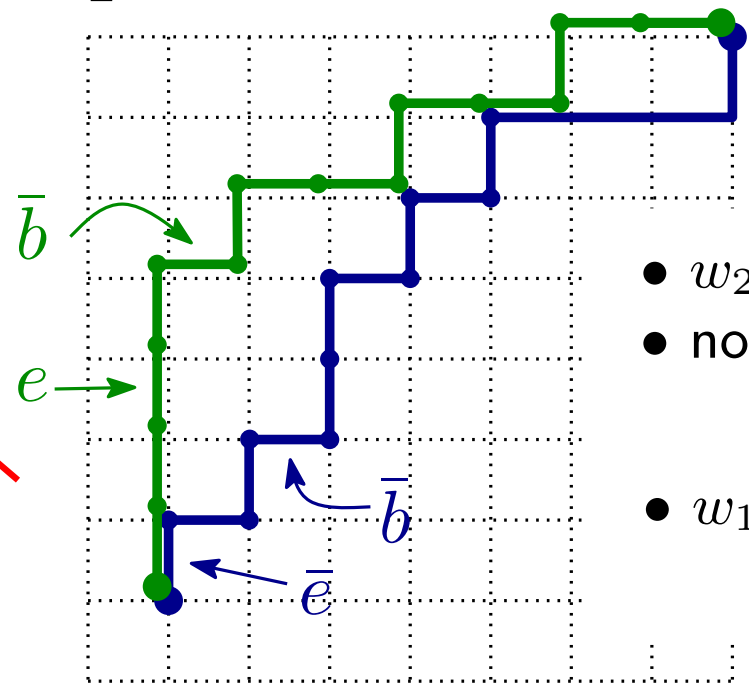
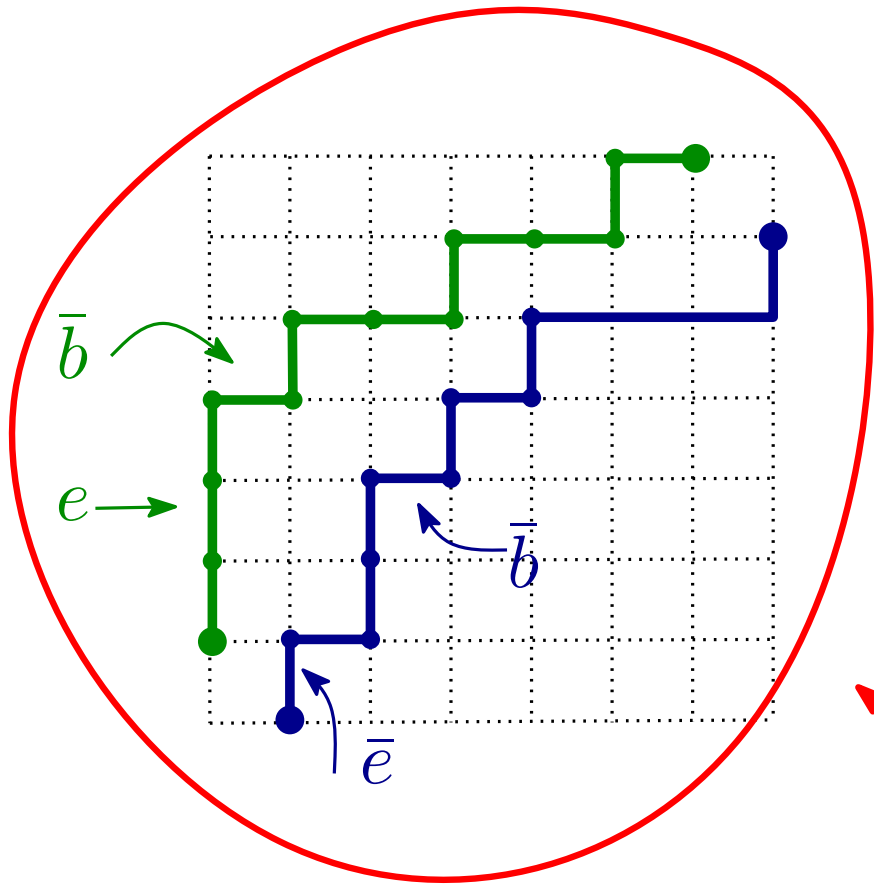
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$$w_1 := w[e, \bar{b}], \quad w_2 := w[\bar{e}, \bar{b}]$$

$$w_1 = e e e e \bar{e} \bar{e} \bar{b} \bar{b} \bar{e} \bar{b} \bar{b} e \bar{b} \bar{b}$$

$$w_2 = \bar{e} \bar{b} \bar{e} \bar{b} \bar{e} \bar{e} \bar{b} \bar{e} \bar{b} \bar{e} \bar{b} \bar{b} \bar{b} \bar{e}$$



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$$+ \\ \bullet w_1 = e \dots \bar{b} \quad w_2 = \bar{e} \bar{b}$$

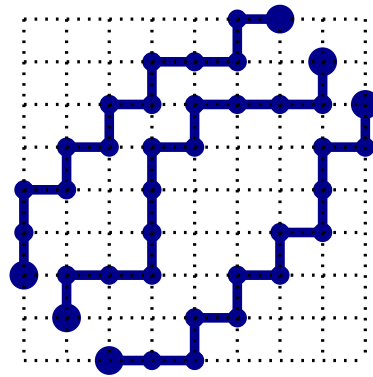
$$w = e b b b \bar{e} e e e \bar{e} \bar{e} \bar{b} \bar{b} \bar{e} \bar{b} \bar{b} \bar{e} \bar{b} \bar{b} \bar{e} \bar{b} \bar{e} \bar{b} \bar{b} \bar{e}$$



Summary

Proposition: [A., Poulalhon]

There exists a one-to-one correspondence between :

$T_{\text{bip}}(i, j)$ and



triplet of non-intersecting paths
with i  and j 
and fixed first and final points

Corollary : The number Θ_{ij} of bipolar orientations with $i + 2$ vertices and $j + 1$ faces is equal to:

$$\Theta_{ij} = \frac{2(i + j)!(i + j + 1)!(i + j + 2)!}{i!(i + 1)!(i + 2)!j!(j + 1)!(j + 2)!}$$

General framework ?

Theorem requires accessible orientation without ccw cycles :

Too much to ask ? **NO !**

Map M fixed + function $\alpha : V(M) \rightarrow \mathbb{N}$,

α -orientation = orientation of the edges such that
 $\forall v \in V(M), \text{out}(v) = \alpha(v)$.

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Proposition: [Felsner '04]

If a map M admits an α -orientation, then there exists a unique α -orientation without ccw cycles.

If there exists one accessible α -orientation, all of them are accessible.

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- Simple triangulations and quadrangulations
- Eulerian and general maps
- Non-separable maps
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} simpler proofs of known bijections

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- Simple triangulations and quadrangulations
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} **New bijections**

General framework ?

Proposition: [Felsner '04]

If a map M admits an α -orientation, then there exists a unique α -orientation without crossings.

If there exists one α -orientation, all of them are accessible.

Our framework **Thank you !**

- Simple triangulations
- Eulerian arrangements
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- Constellations
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