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# A Functorial model of Differential Linear Logic

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Marie Kerjean (UP13), Valentin 'Richie' Maestracci (AMU), Morgan Rogers (UP13)

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# Let's review: Models of LL

What is a model of linear logic ?

# Let's review: Models of LL

What is a model of linear logic ? A possible categorical answer: Seely Categories

$$\begin{array}{ccc} & \mathcal{E}' & \\ & \xrightarrow{\quad} & \\ (\mathcal{C}, \times) & & (\mathcal{L}, \otimes) \\ & \xleftarrow{\quad} & \\ & U & \end{array}$$

A strong monoidal adjunction (!  $\stackrel{\text{def}}{=} \mathcal{E}' \circ U$ )

Between a monoidal category of linear morphisms

And a cartesian category of non linear morphisms

# Let's review: Models of LL

How do we interpret a proof?

A proof  $\pi$  of conclusion  $\Gamma \vdash A$  is interpreted as a morphism  $[[\Gamma]] \rightarrow [[A]]$  compositionally:

$$\frac{\frac{\pi}{A \vdash \Gamma}}{!A \vdash \Gamma}$$

Use a natural transformation  $d_A : !A \rightarrow A$

$$!A \xrightarrow{d} A \xrightarrow{\pi} \Gamma$$

# Let's review: Models of LL

This is a compact version of requiring to have these natural transformations

Operator	Type	Intuition
w	$!A \multimap 1$	Create constant function
c	$!A \multimap !A \otimes !A$	From 2 to 1 parameter
d	$!A \multimap A$	Forget linearity
p	$!A \multimap !!A$	Higher order

Plus some commutative diagrams (to respect cut elimination)

# What is Differential Linear Logic (DiLL) ?

DiLL is a variant of linear logic that was discovered by Ehrhard & Regnier in 2004 with a notion of differentiation:

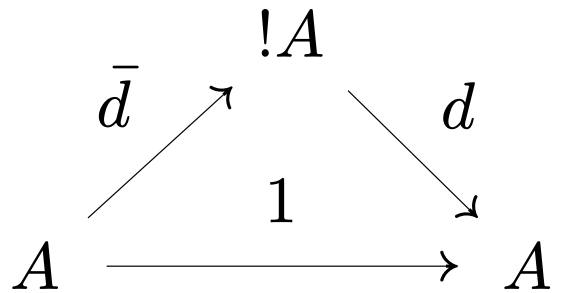

# What is Differential Linear Logic (DiLL) ?

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This amounts to adding these operator:

Operator	Type	Operator	Type	Intuition
w	$!A \multimap 1$	$\bar{w}$	$1 \multimap !A$	Evaluation at 0
c	$!A \multimap !A \otimes !A$	$\bar{c}$	$!A \otimes !A \multimap !A$	Convolution
d	$!A \multimap A$	$\bar{d}$	$A \multimap !A$	Differentiation
p	$!A \multimap !!A$	x	x	x

# Dbar diagrams





# Dbar diagrams 2

$$\begin{array}{ccccc}
 A & \xrightarrow{\bar{d}} & !A & \xrightarrow{p} & !!A \\
 \cong \downarrow & & & & \uparrow m \\
 A \otimes I & \xrightarrow{\bar{d} \otimes w} & !A \otimes !A & \xrightarrow{\bar{d}_! \otimes p} & !!A \otimes !!A
 \end{array}$$

$$\begin{array}{ccc}
 A \otimes !B & \xrightarrow{\bar{d} \otimes 1} & !A \otimes !B \xrightarrow{\varphi} !(A \otimes B) \\
 & \searrow 1 \otimes d & \nearrow \bar{d}_{A \otimes B} \\
 & & A \otimes B
 \end{array}$$

# What is Differential Linear Logic (DiLL) ?

There is a reformulation of such models with a biproduct  $\diamond$  which automatically gives most operators.

$$\begin{array}{ccc} & \mathcal{E}' \approx p & \\ & \xrightarrow{\quad} & \\ (\mathcal{C}, \times) & & (\mathcal{L}, \otimes, \diamond) \\ & \xleftarrow{\quad} & \\ & \mathcal{U} \approx d & \end{array}$$

Type	Operator
$!A \multimap A$	$\bar{d}$

But where would  $\bar{d}$  fit in such a setting ?

Our contribution: A model where  $\bar{d}$  is expressed as a functor.

(Purely functorial)

# Why would we want that ?

- Express a modular transformation of program  
→ (This is a key point of Differentiable Programming)
- Compactify definitions : Proving that something is a model becomes easier
- Makes link easier with Chiralities

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We want to capture the chain-rule in a functorial way:

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$\vec{D}$  should be a functor, with morally  $\vec{D}(f) = D_a(f)$

# The co-Slice Category

The category  $I \downarrow \mathcal{C}$ , the co-Slice of  $\mathcal{C}$  is defined as follows:

- Objects:  $(A, a)$  with  $a : I \rightarrow A$ , intuitively, an element of  $A$
- Arrows:  $f : (A, a) \rightarrow (B, b)$  such that  $f(a) = b$

$$\begin{array}{ccccc} I & \xrightarrow{a} & A & \xrightarrow{f} & B \\ & & & \searrow & \\ & & & & b \end{array}$$

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Now b:

$$D_a(g \circ f) = D(g) \circ D(f) = D_{f(a)}(g) \circ D_a(f)$$

# Second thing

The differential (ie, best linear approximation) of a linear is itself.

Hence  $\vec{D}$  should preserve linear morphisms:  $\forall a, l : D_a(l) = l$

$$\begin{array}{ccc} & & I \downarrow \mathcal{C} \\ & \swarrow \vec{D} & \downarrow \Pi \\ (\mathcal{L}, \diamond) & \xrightarrow{U} & (\mathcal{C}, \times) \end{array}$$

# But...

$$\begin{array}{ccc} & & I \downarrow \mathcal{C} \\ & \swarrow \vec{D} & \downarrow \Pi \\ (\mathcal{L}, \diamond) & \xrightarrow{U} & (\mathcal{C}, \times) \end{array}$$

We cannot go up from  $\mathcal{C}$  to  $I \downarrow \mathcal{C}$  !

→ Would require to choose a point

# The category of Generalized Elements

Given a functor  $U : \mathcal{L} \rightarrow \mathcal{C}$ , the category  $I \downarrow U$  of generalized elements over  $U$  is defined as:

- Objects:  $(A, a)$  with  $a : I \rightarrow U(A)$
- Arrows:  $l : A \rightarrow B$  such that  $U(l)(a) = b$

In a sense, the linear part of the co-Slice.

# Definition of Functorial DiLL Model

A pre-model of DiLL,

$$\begin{array}{ccc} I \downarrow U & \xrightarrow{\vec{U}} & I \downarrow \mathcal{C} \\ \Pi \downarrow & & \downarrow \Pi \\ (\mathcal{L}, \otimes, \diamond) & \xrightarrow[U]{} & (\mathcal{C}, \times) \end{array}$$

# Definition of Functorial DiLL Model

A pre-model of DiLL, plus a functor  $\vec{D}$

$$\begin{array}{ccc} & \vec{U} & \\ & \longrightarrow & \\ I \downarrow U & & I \downarrow \mathcal{C} \\ \Pi \downarrow & \swarrow \vec{D} & \downarrow \Pi \\ (\mathcal{L}, \diamond) & \xrightarrow{U} & (\mathcal{C}, \times) \end{array}$$



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(And well pointedness relative to  $I$  ...)

# Theorems

**Theorem:** Our functorial model is a model of DiLL.

**Theorem:** The converse is true for well pointed models.

# What if my model isn't well pointed?

A pre-model of DiLL, plus a functor  $\vec{D}$

$$\begin{array}{ccc} & \vec{U} & \\ & \longrightarrow & \\ U \downarrow U & & U \downarrow \mathcal{C} \\ \Pi \downarrow & \swarrow \vec{D} & \downarrow \Pi \\ (\mathcal{L}, \diamond) & \xrightarrow{U} & (\mathcal{C}, \times) \end{array}$$

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But we have no model/intuition of what is going on here...

# What are Chiralities ?

A categorical framework by Melliès which refines  $\ast$ -autonomous categories into models of Polarized MLL.

$$\begin{array}{ccc} & \xrightarrow{(-)_{\mathcal{P}}^{\perp}} & \\ (\mathcal{P}, \otimes, 1) & & (\mathcal{N}^{\text{op}}, \wp, \perp) \\ & \xleftarrow{(-)_{\mathcal{N}}^{\perp}} & \end{array} \quad \begin{array}{ccc} & \uparrow & \\ \mathcal{P} & \rightleftarrows & \mathcal{N} \\ & \downarrow & \end{array}$$

With the left being strong monoidal, and  $\downarrow \circ \uparrow = \text{Id}$

Plus some extra conditions on the adjunction

Appears a lot for “smooth” models of DiLL (in Functional Analysis)

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# A parallel to be made

# In Chiralities: Positive vs Negative | In DiLL: Linear vs Non-Linear

$  \begin{array}{ccc}  & (-)^{\perp_P} & \\  (\mathcal{P}, \otimes) & \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} & (\mathcal{N}^{\text{op}}, \wp) \\  & (-)^{\perp_N} &   \end{array}  $	$  \begin{array}{ccc}  & \vec{\uparrow} & \\  \text{El}'_0(\uparrow) & \longrightarrow & \mathcal{N} \downarrow 0 \\  \Pi \downarrow & \swarrow \vec{\downarrow} & \downarrow \Pi \\  (\mathcal{P}, \diamond) & \longrightarrow & (\mathcal{N}, \times) \\  & \uparrow &   \end{array}  $
$  \begin{array}{ccc}  & \mathcal{E} & \\  (\mathcal{C}, \times) & \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} & (\mathcal{L}^{\text{op}}, \otimes) \\  & U' &   \end{array}  $	$  \begin{array}{ccc}  & \vec{U} & \\  I \downarrow U & \longrightarrow & I \downarrow \mathcal{C} \\  \Pi \downarrow & \swarrow \vec{D} & \downarrow \Pi \\  (\mathcal{L}, \diamond) & \longrightarrow & (\mathcal{C}, \times) \\  & U &   \end{array}  $

# Last but not least: a funny remark

When  $\mathcal{L}$  is a calculus category (so with integration on top), we have an relative  $! \otimes$  Id-adjunction:

$$I \downarrow \mathcal{C} \begin{array}{c} \xrightarrow{\vec{D}} \\ \xleftarrow{\overline{U}} \end{array} (\mathcal{L}, \otimes)$$

$$(I \downarrow \mathcal{C})((a, A), (b, B)) \simeq \mathcal{L}(!A \otimes A, B)$$

With  $\overline{U}(l) = (U(l), U(u_A))$

This adjunction corresponds mathematically to the fundamental theorem of calculus !



# Future Work

- Pursue similarities and extending our framework to Polarized LL with Mellie's **Chiralities**
- New model of Kerjean & Pacaud-Lemay with co-promotion should have a symmetric behavior: a functorial **co-chain rule** ?
- Investigate the **dependent** flavour

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**Thank you for listening !**

**Any questions ?** 😊