# Measuring well quasi-orders and complexity of verification

(SLIDES FROM MY PHD DEFENSE)

Isa Vialard

November 20, 2024



### Some interesting sequences

 $\circ$  <  $\circ$  < V V  $\vee$   $\vee$   $\vee$   $\vee$ V V V V  $\vee$   $\vee$   $\vee$   $\vee$ V V V V  $\vee$ V V V V V V V V V V V /  $\vee$ V 0 < 0 < 0 < 0 < 0**ॉ**< o < o < o < o < o V V  $\mathbf{V}$ V V  $\vee$  $\vee$ V V V  $\vee$ V V.V V V (0,1) < 0 $\mathbf{i} < \mathbf{0} < \mathbf{0}$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$ 

- decreasing sequence

(5,5) > (4,4) > (4,3) > (2,3) > (1,1)

## Some interesting sequences



- decreasing sequence
- incomparable sequence (or antichain)
   i.e. pairwise incomparable

 $(1,9) \perp (3,8), (4,7), (7,5), \ldots$ 

## Some interesting sequences



#### **Definitions: Well Quasi-Order**



- decreasing sequence
- antichain i.e. pairwise incomparable
- bad sequence
  i.e. pairwise non increasing

#### 

No infinite antichain or decreasing seq

## Definitions: Well Quasi-Order

 $< \circ < < (\mathbb{N} \times \mathbb{N}, \leq_{\times})$ < < < >< 0 < 0Ο 0 V V V V  $\vee$ 0 V V  $< \circ < \circ < \circ < \circ < \circ < \circ < \circ$ 0 240 <V 0  $\vee$ 0 < 0 < 0 < 0 < $\mathbf{v} < \mathbf{o} < \mathbf{o} < \mathbf{o}$ V V V 0 < 0 < 0< 0 < 0 < 0 < 0 < 0 $\vee$ V Ο V (0,1) < 0V V V 

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## WQO \$

No infinite antichain or decreasing seq



WQO

 $\Diamond$ 

No infinite antichain or decreasing seq

 $\updownarrow$ 

Some see wqos as wells Blass & Gurevich (2008)



WQO

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WQO

↕

No infinite antichain or decreasing seq

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Some see wqos as wells Blass & Gurevich (2008)



- Reasons to study wqos
  - "It is fun" (Kříž & Thomas (1990))

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# Well-structured transition systems

Finkel (1994), Abdulla& Jonsson (1996)

- Set of configurations: WQO
- $\leq$  a simulation relation



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• Ex: Counter machines, Petri nets, VASS, Lossy channel systems . . .



Vector Addition Systems with States

## Well-structured transition systems

Finkel (1994), Abdulla& Jonsson (1996)

- Set of configurations: WQO
- $\leq$  is upward-compatible





• Ex: Counter machines, Petri nets, VASS, Lossy channel systems . . .



# Complexity and expressiveness

Schmitz& Schnoebelen(2011)

- Controlled bad sequences (even decreasing, or antichains)
- Can we bound the length of controlled sequences by measuring wqo?

# Measuring wqos

Natural notions of measure when finite



Finite subsets of  $\{1, 2, 3, 4\}$  ordered by  $\subseteq$ .

## Let's extend height and width to infinite wqos

#### Let's extend height and width to infinite wqos with ordinals

 $\circ$  <  $N \times N < \circ$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$ 

Width and height: at least  $\omega$ 

#### Let's extend height and width to infinite wqos with ordinals

 $0 < 0 < 0 < 0 < 0 < 0 < 0 < 0 < 0 > N \times N < 0$  $\bigvee$ VV V V V V V V V V V V V V V V  $\vee$ V V V V V < Q < 0 < 0 < 0 < 0 < 0 < 0 < 00 < 0 < 0 < 0 < 0< 0< 0 < 0(0. < 0< 0 < 0 < 0(0, 0) $< \delta < 0 < 0 < 0$ 

Width and height: at least  $\omega$ 

Counting elements: at least  $\omega$ 

#### Let's extend height and width to infinite wqos with ordinals

 $\circ$  <  $\wedge$   $\mathbb{N}$   $\times$   $\mathbb{N}$  <  $\circ$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$ V  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$ V V  $(0, 1) \leftarrow \bigcirc \leftarrow \bigcirc$  $\vee$   $\vee$   $\vee$   $\vee$   $\vee$   $\vee$  $(0, \theta) \leftarrow 0 \leftarrow 0 \leftarrow 0 \leftarrow 0$ 

Width and height: at least  $\omega$ 

Counting elements: at least  $\omega^2$ 

ω

ω

 $(\omega)$ 

## Definition (Maximal order type, Width and Height)

$$o(X)$$
  
w(X) = ordinal rank of the tree of 
$$\begin{cases} bad sequences \\ antichain sequences \\ decreasing sequences \end{cases}$$
 in X.

Definition from Kříž & Thomas(1990) (first definition of ordinal width)

First definition of maximal order type by De Jongh & Parikh(1977)



















# Measuring with games



• Game:  $\alpha$  vs w(X)

- Initial configuration:
  - Odile :  $\gamma = \alpha$ ,
  - Antoine :  $S = \emptyset$
- Player alternate:
  - Odile picks  $\gamma' < \gamma$
  - Antoine extends *S* into
    - S :: x an antichain,
- End: You lose if you cannot play anymore

# Measuring with games



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- End: You lose if you cannot play anymore

Theorem (Blass & Gurevich (2008))

- Antoine has winning strategy when Odile begins  $\Leftrightarrow \alpha \leq w(X)$
- Odile has winning strategy when Antoine begins  $\Leftrightarrow \alpha \ge w(X)$

# Example: Playing on the on disjoint sum



Disjoint sum  $A \sqcup B$ 

## **Theorem:** $o(A \sqcup B) = o(A) \oplus o(B)$ (De Jongh & Parikh(1977))

# Example: Playing on the on disjoint sum



**Theorem:**  $o(A \sqcup B) = o(A) \oplus o(B)$  (De Jongh & Parikh(1977))

Ex:  $(\omega^{\omega} + \omega^3) \oplus (\omega^5 + \omega + 1) = \omega^{\omega} + \omega^5 + \omega^3 + \omega + 1$ 

# Example: Playing on the on disjoint sum



Disjoint sum  $A \sqcup B$ 

**Theorem:**  $o(A \sqcup B) = o(A) \oplus o(B)$  (De Jongh & Parikh(1977))

This theorem is easy to prove with games!



 $o(A \sqcup B) \le o(A) \oplus o(B)$  if Odile wins when Antoine begins  $o(A \sqcup B) \ge o(A) \oplus o(B)$  if Antoine wins when Odile begins



 $o(A \sqcup B) \le o(A) \oplus o(B)$  if Odile wins when Antoine begins  $o(A \sqcup B) \ge o(A) \oplus o(B)$  if Antoine wins when Odile begins












Disjoint sum  $A \sqcup B$ 





Direct sum A + B





Direct sum A + B



Cartesian product  $A \times B$ 



А

А

	Space	M.O.T.	Height	Width
Disjoint sum	$A \sqcup B$	$o(A)\opluso(B)$	$\max(h(A), h(B))$	$w(A)\oplusw(B)$
Direct sum	A + B	o(A) + o(B)	h(A) + h(B)	$\max(\mathbf{w}(A),\mathbf{w}(B))$
Cartesian prod.	$A \times B$	$o(A)\otimeso(B)$	$h(A) \oplus h(B)$	?
Direct prod.	$A \cdot B$	?	$h(A) \cdot h(B)$	$w(A) \odot w(B)$

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Disjoint sum	$A \sqcup B$	$o(A)\opluso(B)$	$\max(h(A), h(B))$	$w(A)\oplusw(B)$
Direct sum	A + B	o(A) + o(B)	h(A) + h(B)	$\max(\mathbf{w}(A), \mathbf{w}(B))$
Cartesian prod.	$A \times B$	$o(A)\otimeso(B)$	$h(A) \oplus h(B)$	?
Direct prod.	$A \cdot B$	?	$h(A) \cdot h(B)$	$w(A) \odot w(B)$
Fin. words	<i>A</i> *	$\omega^{\omega^{(o(A)^{\pm})}}$	$h^*(A)$	$\omega^{\omega^{(\circ(A)^{\pm})}}$
	$M^\diamond(A)$	$\omega^{\widehat{\mathbf{o}(A)}}$	$h^*(A)$	?
Fin. multisets	$M^{o}(A)$	$\omega^{\circ(A)}$	?	?
Fin. Powerset	$P_{f}(A)$	?	?	?

Space	М.О.Т.	Height	Width
$A \sqcup B$	$o(A)\opluso(B)$	$\max(h(A), h(B))$	$w(A)\oplusw(B)$
A + B	o(A) + o(B)	h(A) + h(B)	$\max(w(A), w(B))$
A  imes B	$o(A)\otimeso(B)$	$h(A)  \hat{\oplus}  h(B)$	$\geq w(o(A) \times o(B))$
$A \cdot B$	$o(A) \cdot pred_k(o(B)) + o(A) \otimes k$	$h(A) \cdot h(B)$	$w(A)\odotw(B)$
A*	$\omega^{\omega^{(\mathrm{o}(A)^{\pm})}}$	$h^*(A)$	$\omega^{\omega^{(\mathrm{o}(A)^{\pm})}}$
$M^\diamond(A)$	$\omega^{\widehat{\mathbf{o}(A)}}$	$h^*(A)$	$\omega^{\widehat{\mathbf{o}(A)}-1}$
$M^{o}(A)$	$\omega^{o(A)}$	$\omega^{h(A)}$	$\omega^{o_{\perp}(A)}$
$P_{f}(A)$	$\leq 2^{o(A)}$	$\leq 2^{h(A)}$	$\geq 2^{w(A)}$

#### Back in time

Space	M.O.T.	Height	Width
$A \sqcup B$	$o(A)\opluso(B)$	$\max(h(A), h(B))$	$w(A)\oplusw(B)$
A + B	o(A) + o(B)	h(A) + h(B)	$\max(\mathbf{w}(A), \mathbf{w}(B))$
$A \times B$	$o(A)\otimeso(B)$	$h(A) \oplus h(B)$	?
$A \cdot B$	?	$h(A) \cdot h(B)$	$w(A) \odot w(B)$
A*	$\omega^{\omega^{(o(A)^{\pm})}}$	$h^*(A)$	$\omega^{\omega^{(o(A)^{\pm})}}$
$M^\diamond(A)$	$\widehat{\omega^{\circ(A)}}$	$h^*(A)$	?
$M^{o}(A)$	$\omega^{\circ(A)}$	?	?
$P_{f}(A)$	?	?	?

#### Quick look at the direct product



Lexicographic product  $A \cdot B$ 

#### • I was told that $o(A \cdot B) = o(A) \cdot o(B)$

... but only the lower bound is true:  $o(A \cdot B) \ge o(A) \cdot o(B)$  Mistake noticed by Harry Altman (March, 2024)





### Quick look at the direct product



 $(\omega + 1) \cdot 
abla$   $(\omega + 1) \cdot \Delta$ 

### Quick look at the direct product



$$o((\omega + 1) \cdot \nabla) = o((\omega + 1) \cdot \Delta) =$$

$$[(\omega + 1) \oplus (\omega + 1)] + +$$

$$(\omega + 1) = [(\omega + 1) \oplus (\omega + 1)]$$

$$= \omega \cdot 3 + 2 = \omega \cdot 3 + 1$$

$$= o(\omega + 1) \cdot o(\nabla)$$

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$A \sqcup B$	$o(A)\opluso(B)$	$\max(h(A), h(B))$	$w(A)\oplusw(B)$
A + B	o(A) + o(B)	h(A) + h(B)	$\max(w(A), w(B))$
A  imes B	$o(A)\otimeso(B)$	$h(A)  \hat{\oplus}  h(B)$	?
$A \cdot B$	Not functional	$h(A) \cdot h(B)$	$w(A)\odotw(B)$
A*	$\omega^{\omega^{(\mathrm{o}(A)^{\pm})}}$	$h^*(A)$	$\omega^{\omega^{(o(A)^{\pm})}}$
$M^\diamond(A)$	$\omega^{\widehat{\mathbf{o}(A)}}$	$h^*(A)$	?
$M^{o}(A)$	$\omega^{o(A)}$	?	?
$P_{f}(A)$	?	?	?

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$P_{f}(A)$	?	?	?

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$M^{o}(A)$	$\omega^{o(A)}$	$\omega^{h(A)}$	Not functional
$P_{f}(A)$	Not functional	Not functional	Not functional

#### Non functional example for P<sub>f</sub>



 $Y_1 = (\omega + \omega) \sqcup (\omega + \omega) \qquad \qquad Y_2 = (\omega \sqcup \omega) + (\omega \sqcup \omega)$ 

 $f(\mathsf{P}_{\mathsf{f}}(\underline{\mathsf{Y}_1})) \neq f(\mathsf{P}_{\mathsf{f}}(\underline{\mathsf{Y}_2})) \text{ for } f = \mathsf{o},\mathsf{h},\mathsf{w}$ 



# **Fixing non-functionality**

An underrated measure: maximal number of elements

# **Theorem (M.o.t. of the direct product)** $o(A \cdot B) = o(A) \cdot pred^{k}(o(B)) + o(A) \otimes k$ if $max\_elt(B) = k$



With Mirna Džamonja



# Theorem (M.o.t. of the direct product) $o(A \cdot B) = o(A) \cdot pred^{k}(o(B)) + o(A) \otimes k$ if $max\_elt(B) = k$



With Mirna Džamonja

If B has k maximal elements, then  $o(B) = \lambda + m$  with  $m \ge k$ 

Then  $o(A \cdot B) = o(A) \cdot (\lambda + (m - k)) + o(A) \otimes k$ 

 $(A, \leq_A)$  is an augmentation of  $(B, \leq_B)$  iff

- Same support: A = B
- A has more relations:  $\leq_B \subseteq \leq_A$



If  $A \geq_{\mathsf{aug}} B$  then  $\mathsf{o}(A) \leq \mathsf{o}(B)$  and  $\mathsf{w}(A) \leq \mathsf{w}(B)$ 

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If  $A \geq_{\mathsf{aug}} B$  then  $\mathsf{o}(A) \leq \mathsf{o}(B)$  and  $\mathsf{w}(A) \leq \mathsf{w}(B)$ 

• If  $A \ge_{aug} B$  then  $C(A) \ge_{aug} C(B)$ ,

for most wqo-constructors C

#### Proof: First separate infinite and finite part

If B has k maximal elements, then  $o(B) = \lambda + m$  with  $m \ge k$ 

We can partition B as:

- $B_{\lambda}$ , with  $o(B_{\lambda}) = \lambda$
- $B_m$ , with  $o(B_m) = m$



#### **Proof:** First separate infinite and finite part



## Focus on $A \cdot B_m$ : Lower bound



$$o(A \cdot B_m) \ge o(A \cdot ((m-k) + \Gamma_k)) = o(A) \cdot (m-k) + o(A) \otimes k$$

#### Lemma

If max\_elt(B) = 1 then B = B' + 1, then  $o(A \cdot B) = o(A \cdot B') + o(A) = o(A) \cdot o(B)$  by induction



 $\mathsf{o}(A \cdot B_m) \leq \mathsf{o}(A) \cdot m_1 \otimes \cdots \otimes \mathsf{o}(A) \cdot m_k = \mathsf{o}(A) \cdot (m-k) + \mathsf{o}(A) \otimes k$ 

# Fixing non-functionality ?

Fixing the multiset ordering : a fourth ordinal invariant!
#### Definition (Friendly order type)

 $o_{\perp}(X) =$  rank of the tree of *open-ended* bad sequences



#### Definition (Friendly order type)

#### $o_{\perp}(X) =$ rank of the tree of *open-ended* bad sequences



#### The fourth ordinal invariant

Theorem (Width of $M^\circ$ )						
$w(M^{o}(X)) = \omega^{o_{\perp}(X)}$						
	Space	o,h,w	o⊥			
			?			
			?			
			?			

- How to compute the fot?
  - Exists  $X' \subseteq X$  such that  $\mathsf{Bad}(X') \subseteq \mathsf{Bad}_{\perp}(X)$
  - $limit\_part(o(str(X))) \le o_{\perp}(X) \le o(str(X))$  with  $str(X) = \{ x \in X \mid \exists y \in X, y \perp x \}$
  - w(X)  $-1 \leq o_{\perp}(X)$

• if 
$$w(A) = o(A)$$
 limit, then  $o_{\perp}(X) = o(X)$ 

 $(A \sqcup B) = o(A) \oplus o(B)$ 

31/33

Space	М.О.Т.	Height	Width
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A + B	o(A) + o(B)	h(A) + h(B)	$\max(w(A), w(B))$
$A \times B$	$o(A)\otimeso(B)$	$h(A)  \hat{\oplus}  h(B)$	$\geq w(o(A) \times o(B))$
A · B	$o(A) \cdot \textit{pred}^k(o(B)) + o(A) \otimes k$	$h(A) \cdot h(B)$	$w(A) \odot w(B)$
	if $max_elt(B) = k$		
<i>A</i> *	$\omega^{\omega^{(\mathrm{o}(A)^{\pm})}}$	$h^*(A)$	$\omega^{\omega^{(o(A)^{\pm})}}$
$M^\diamond(A)$	$\omega^{\widehat{o(A)}}$	$h^*(A)$	$\omega^{\widehat{\mathbf{o}(A)}-1}$
$M^{o}(A)$	$\omega^{o(A)}$	$\omega^{h(A)}$	$\omega^{o_{\perp}(A)}$
$P_{f}(A)$	$\leq 2^{o(A)}$	$\leq 2^{h(A)}$	$\geq 2^{w(A)}$

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  - is fun!
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  - Other approach: Elementary family of wqos

 $E := \alpha \ge \omega^{\omega}$  mult. indec.  $|E_1 \sqcup E_2 | E_1 \times E_2 | M^{\diamond}(E) | M^{\circ}(E) | E^* | P_f(E)$ 

Wqos that appear in well-structured transition systems!

- Measuring well quasi-orders
  - is fun!
  - Often not functional but... maybe we are just missing some measures?
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Wqos that appear in well-structured transition systems!

#### Open questions

- New invariants:
  - Computing the fot
  - Is there an invariant that would make CP and  $\mathsf{P}_\mathsf{f}$  functional?
- New operations: Infinite words, variants of trees, graph minor, ...

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