

On recent results by Zuiddam and Bukh & Cox on the Shannon capacity

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with Sven Polak

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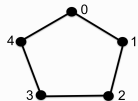
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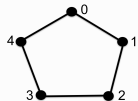
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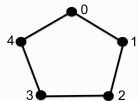
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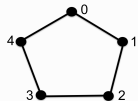
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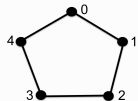
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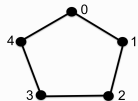
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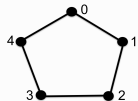
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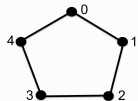
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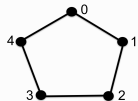
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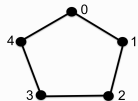
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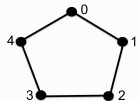
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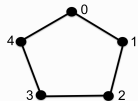
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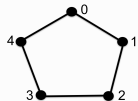
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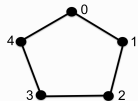
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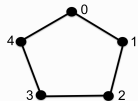
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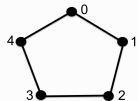
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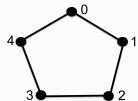
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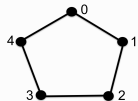
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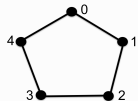
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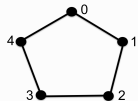
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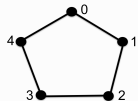
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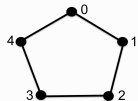
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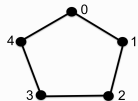
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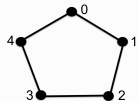
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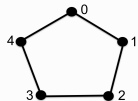
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It is easily seen that

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Proof. The "diagonal" in $G \cdot \bar{G}$ is independent, hence

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On the other hand, we have by Theorems 1, 6 and 7 that

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This proves the theorem. The proof also shows that in these cases $\Theta = \nu$.

Theorem 12. Let $n \geq 2r$ and let the graph $K(n, r)$ be defined as the graph whose vertices are the r -subsets of an n -element set S , two being adjacent iff they are disjoint. Then

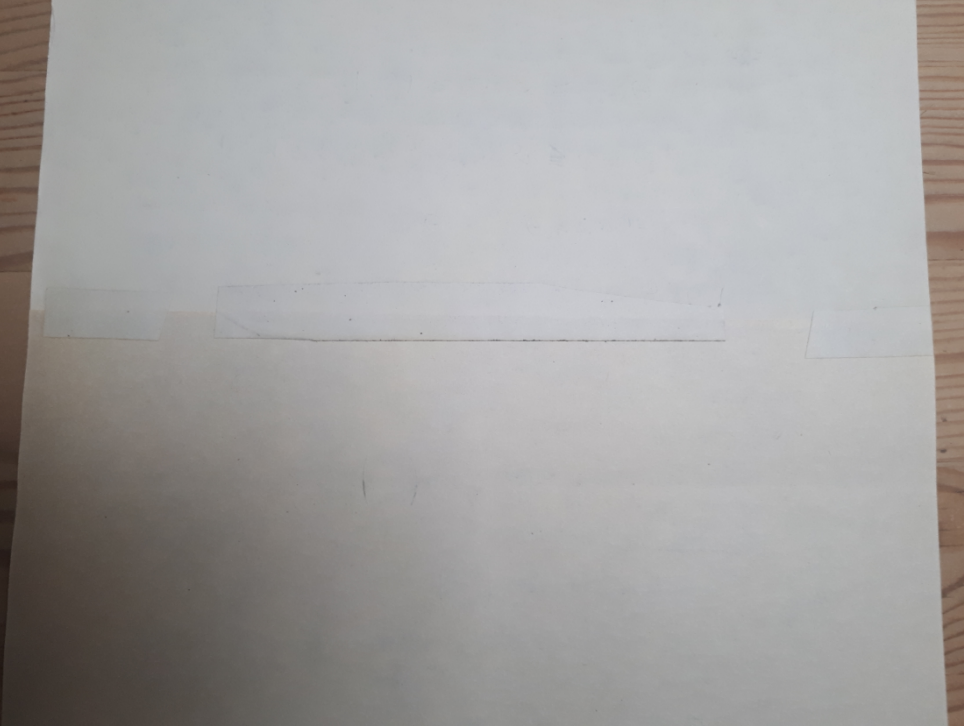
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Corollary (Erdős-Ko-Rado Theorem)

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Corollary (Erdős-Ko-Rado Theorem)

$$\alpha(K(n, r)) = \binom{n-1}{r-1}.$$

Proof. The r -subsets containing a specified element of S form an independent set of points in $K(n, r)$; hence

$$\alpha(K(n, r)) \geq \alpha(K(n, r)) \geq \binom{n-1}{r-1}$$

Proof. The "diagonal" in $G \cdot G$ is independent, hence

$$\Theta(G \cdot G) \geq \alpha(G \cdot G) \geq |V(G)|.$$

On the other hand, we have by Theorems 1, 6 and 7 that

$$\Theta(G \cdot G) \leq \vartheta(G \cdot G) = \vartheta(G)\vartheta(G) = |V(G)|.$$

If G is self-complementary then

$$\Theta(G \cdot G) = \alpha(G \cdot G) = \alpha(G)^2.$$

This proves the theorem. It also shows that in these cases $\Theta = \vartheta$.

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Willem Haemers 28/2/78.

Let G be a graph on n vertices. Let A be a symmetric matrix of size n such that $(A)_{ii} = 1$, $(A)_{ij} = 0$ iff i and j correspond to nonadjacent vertices of G .

Theorem $\Theta(G) \leq \text{rank}(A)$

Proof 1°. $\text{rank}(A^{*k}) = (\text{rank}(A))^k$

2°. A^{*k} has submatrix I of size $\alpha(A^{*k})$, hence $\text{rank}(A^{*k}) \geq \alpha(A^{*k})$. We now have

$$\Theta(G) = \sup \sqrt[k]{\alpha(A^{*k})} \leq \sup \sqrt[k]{\text{rank}(A^{*k})} = \text{rank}(A) \quad \square$$

Example G is the Schlegel graph on 27 vertices. The eigenvalues of its $(0,1)$ adjacency matrix B are 1, -5, 10 with multiplicities 20, 6 and 1 resp. Put $A = I - B$, then $\text{rank}(A) = 7$, so $\Theta(G) \leq 7$. G and $\bar{G}$ are edge transitive, so we have $\Theta(G) = 27 \frac{5}{15} = 9$, $\Theta(\bar{G}) = \frac{27}{3} = 9$, in fact $\Theta(\bar{G}) = 3$ because $\bar{G}$ admits a 3-coclique.

So in the above example we have seen:

$$\Theta(\bar{G}) \neq \Theta(G) \text{ and } \Theta(G) \cdot \Theta(\bar{G}) \leq 21 < 27.$$

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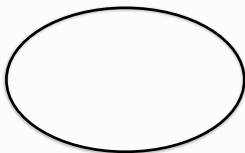
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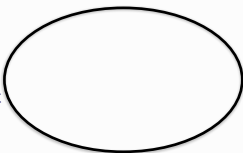
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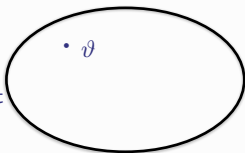
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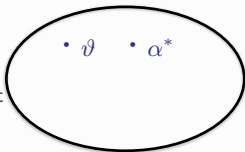
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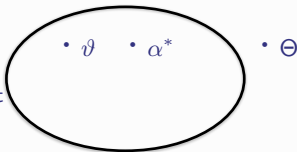
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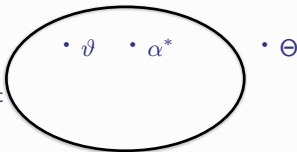
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Find $f \in \Delta$ with $f(C_7) = \Theta(C_7)$.

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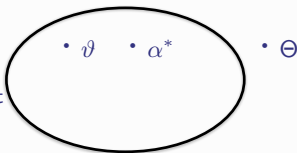
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$\varphi : V(G) \rightarrow V(H)$ with
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Δ
compact



👉 Find $f \in \Delta$ with $f(C_7) = \Theta(C_7)$.

👉 Is Δ connected ?

Jeroen Zuiddam's theorem (2018)

$$\Theta(G) = \min_{f \in \Delta} f(G),$$

where $\Delta := \{f : \{\text{graphs}\} \rightarrow \mathbb{R} \mid \text{for all graphs } G, H :$

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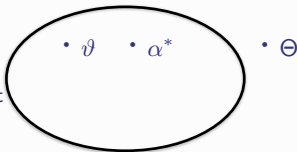
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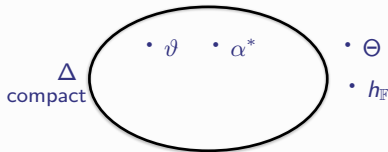
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What about the Haemers bound???

The fractional Haemers bound

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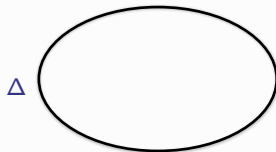
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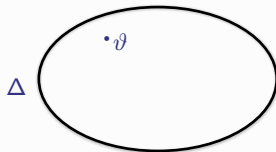
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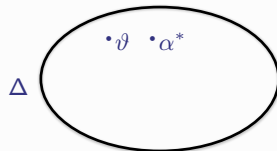
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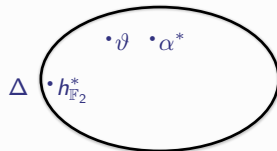
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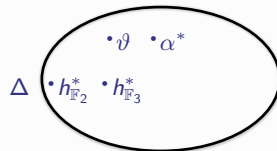
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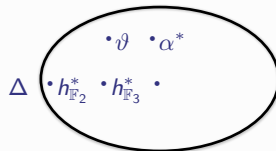
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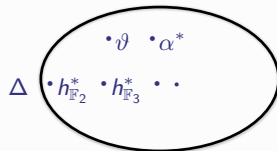
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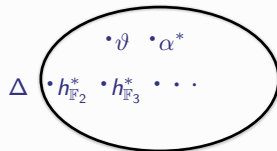
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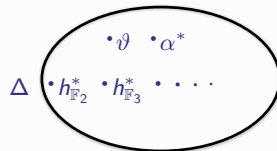
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Bukh & Cox 2018: **Theorem:** $h_{\mathbb{F}}^*$ is multiplicative
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Let $\mathbb{F}$ be a field and $G = (V, E)$ be a graph.

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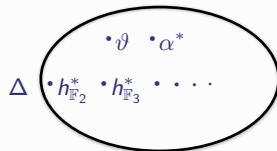
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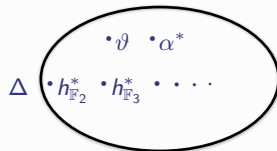
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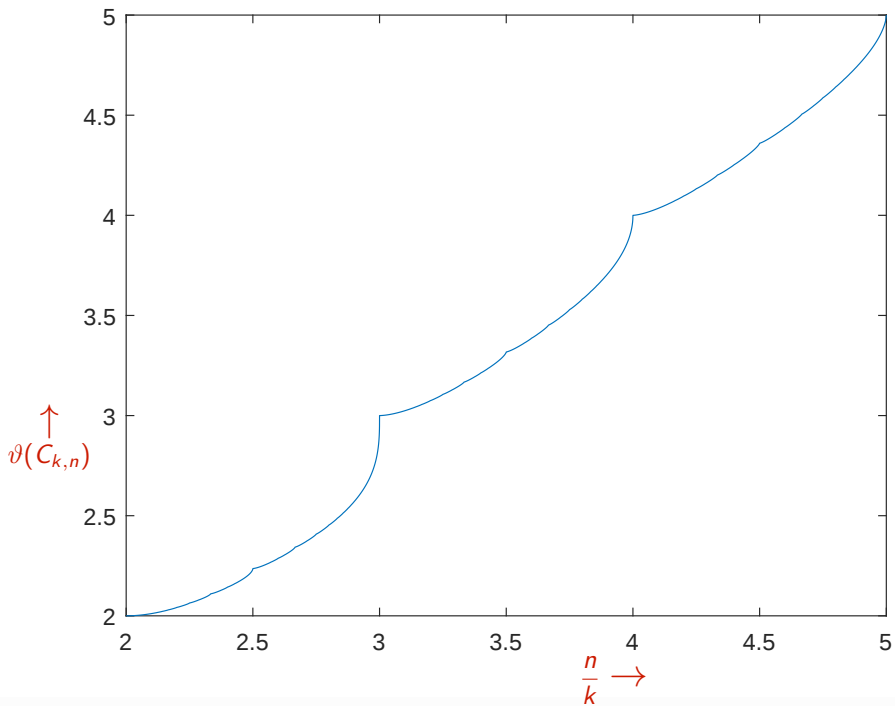
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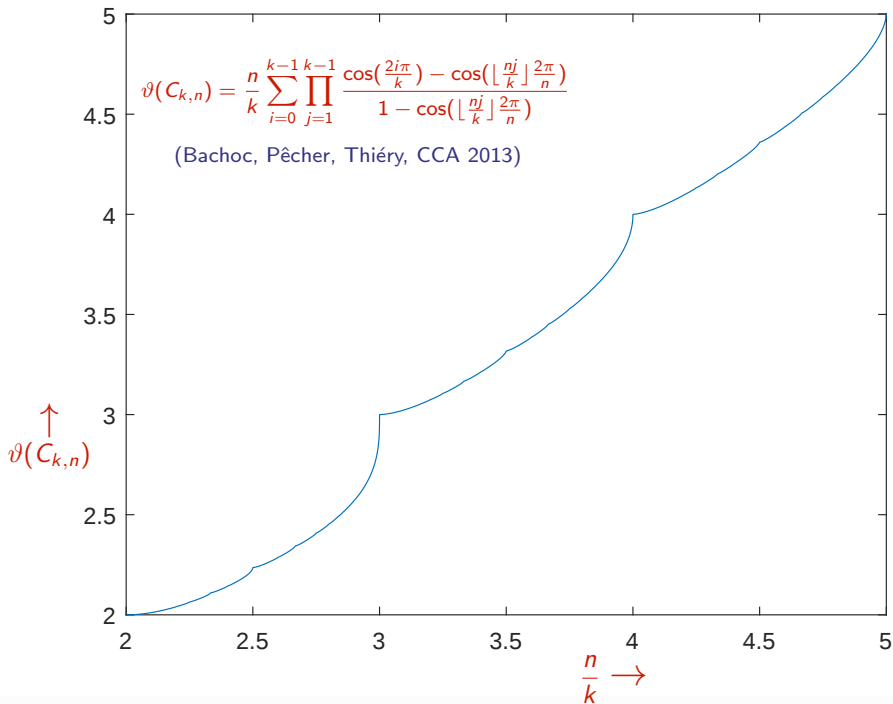
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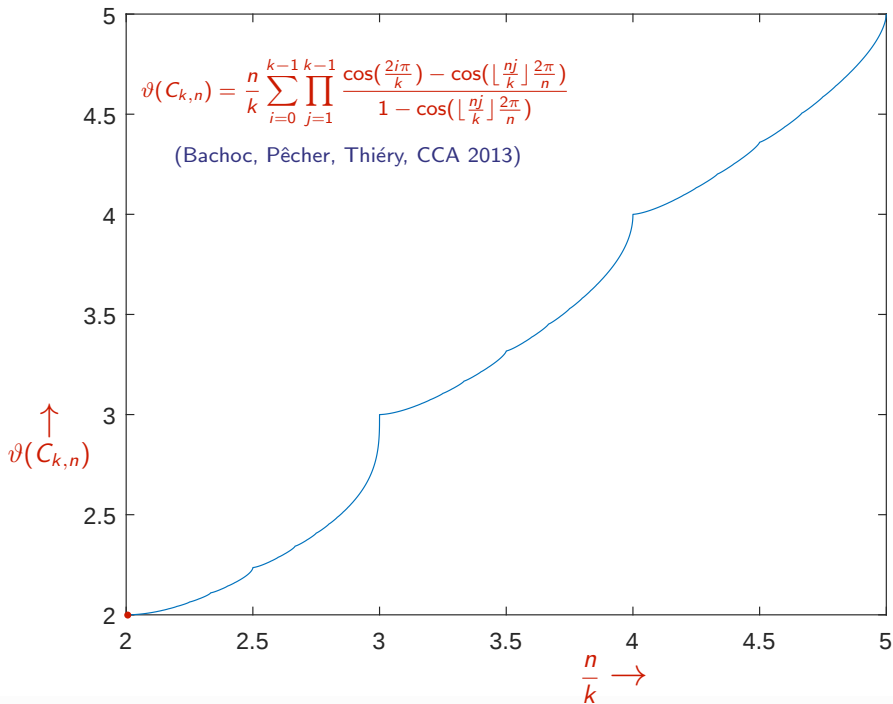
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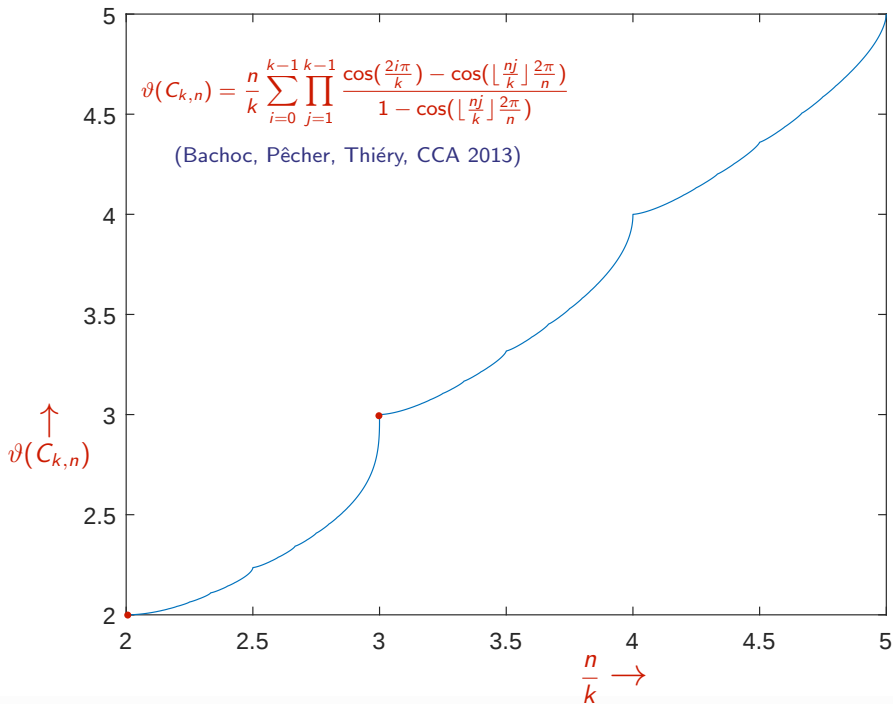
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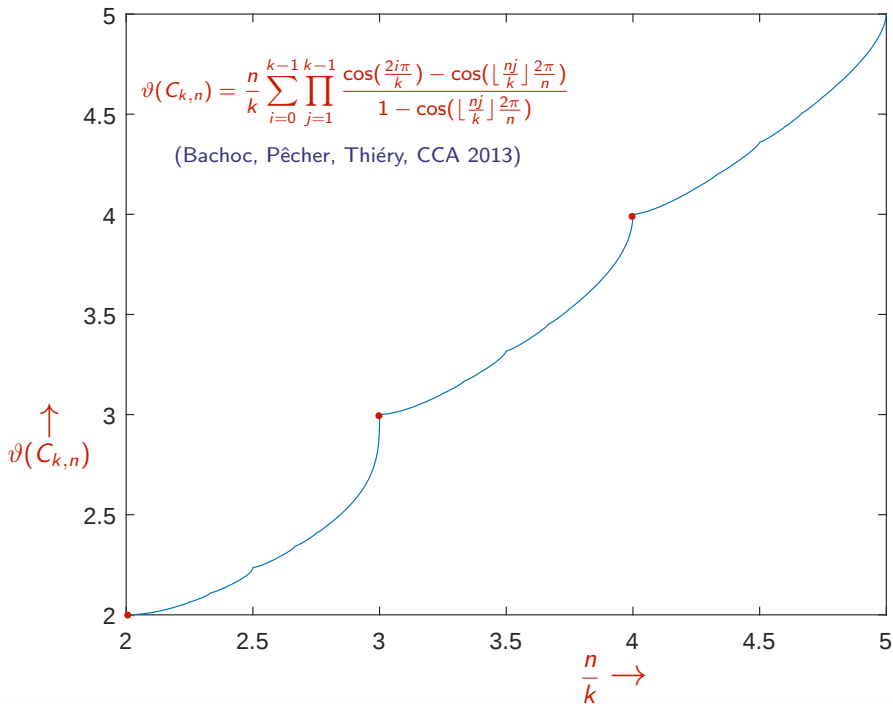
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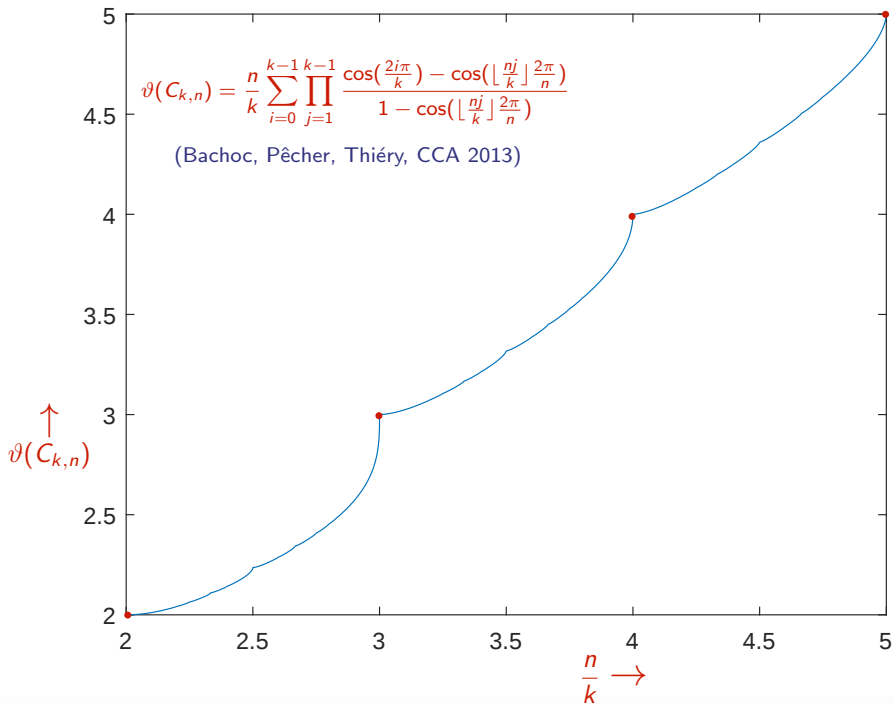


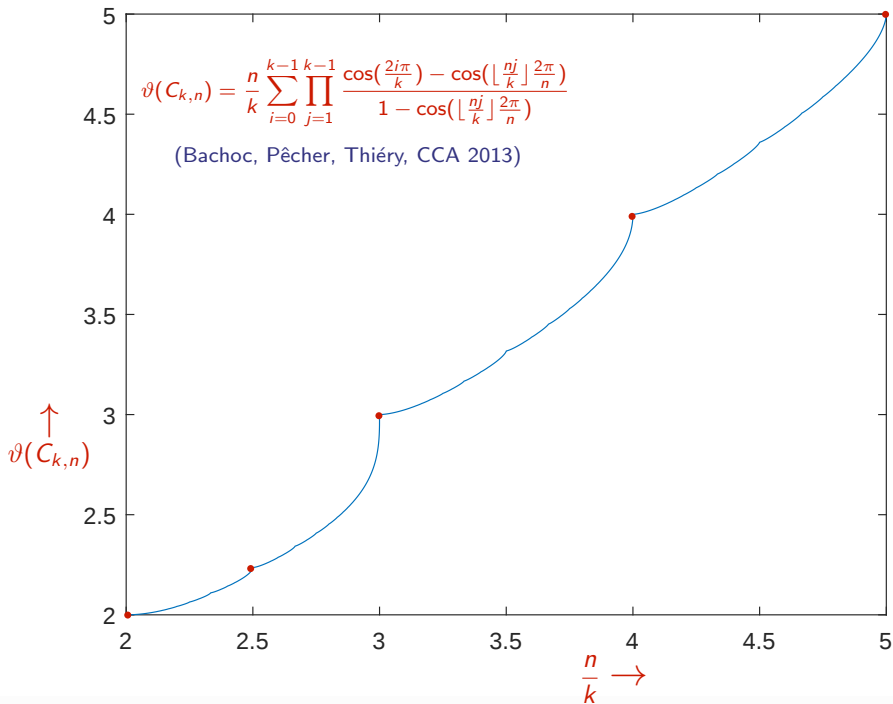


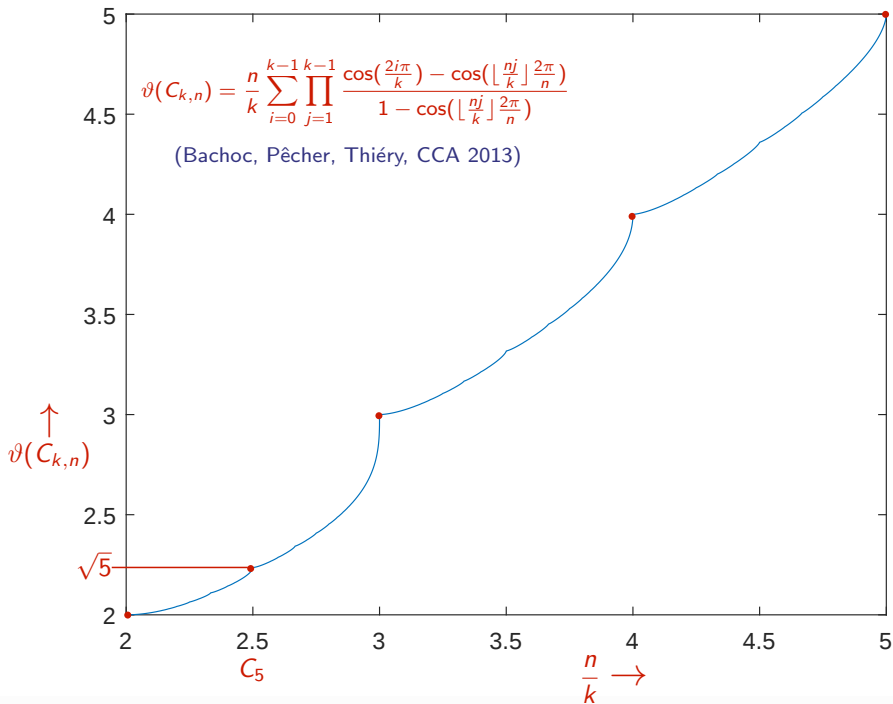


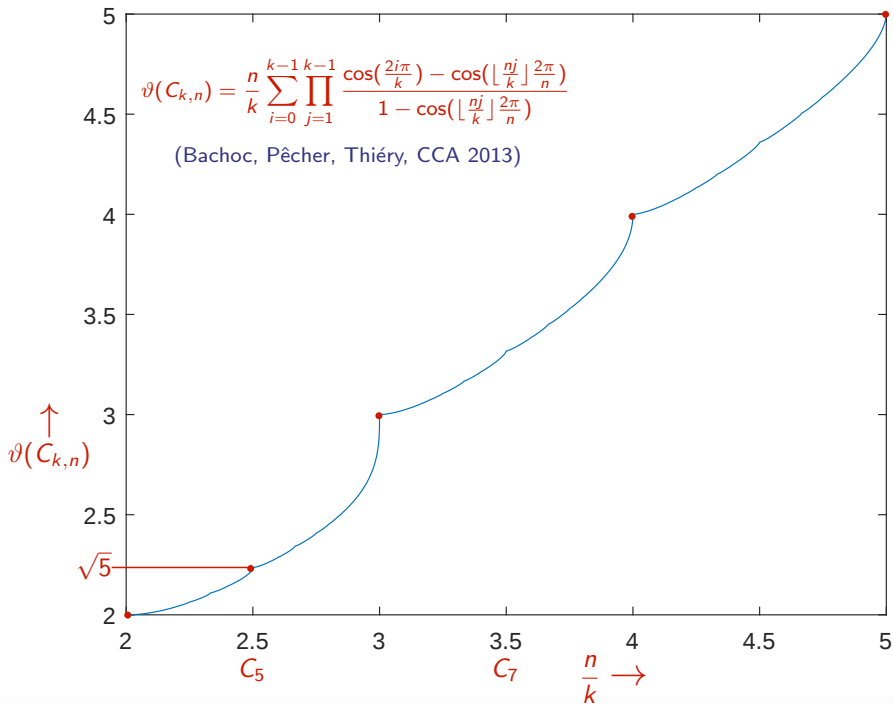


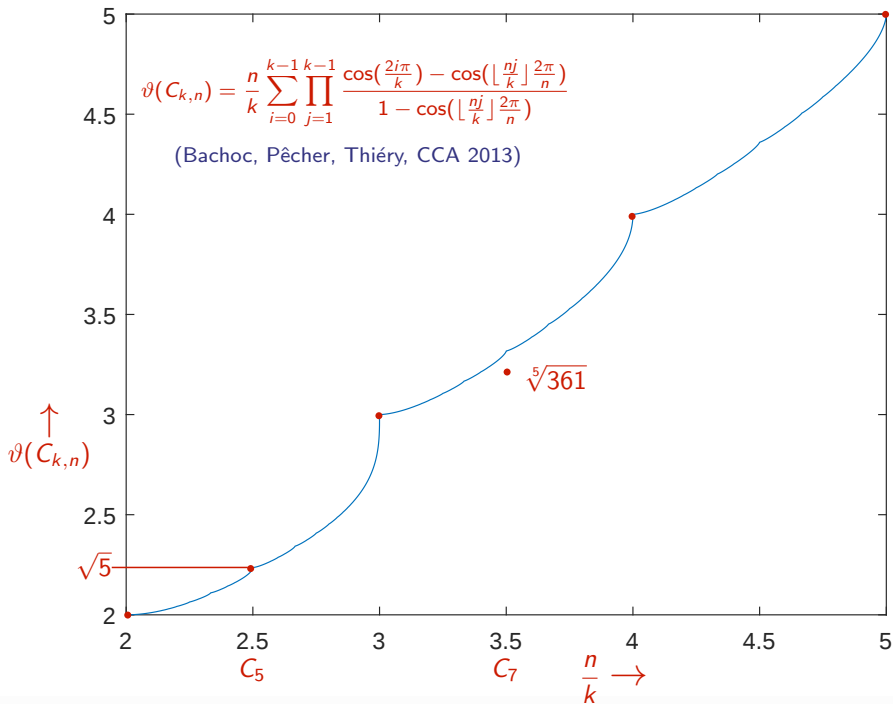


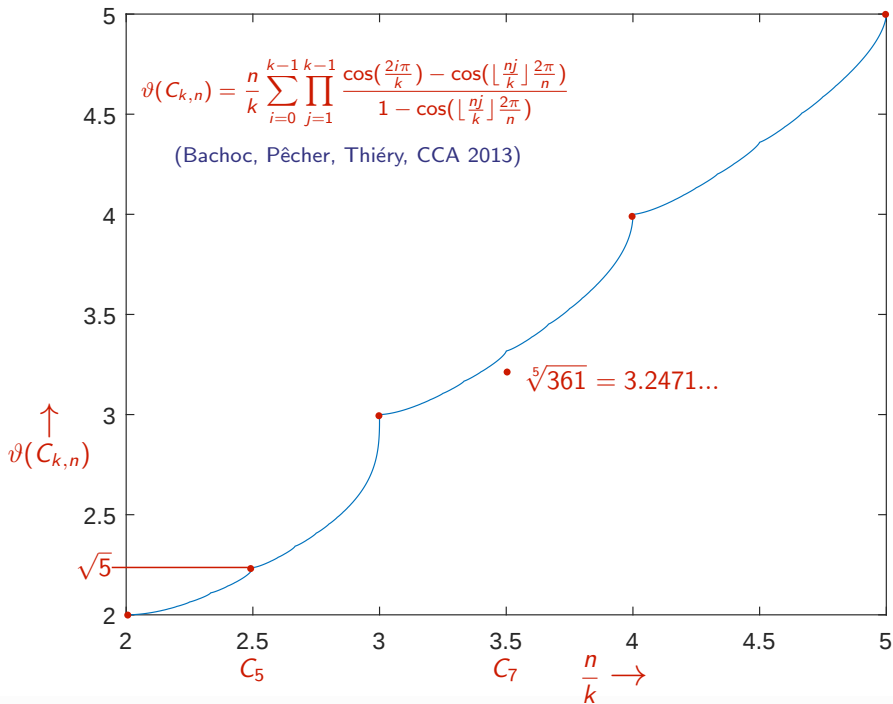


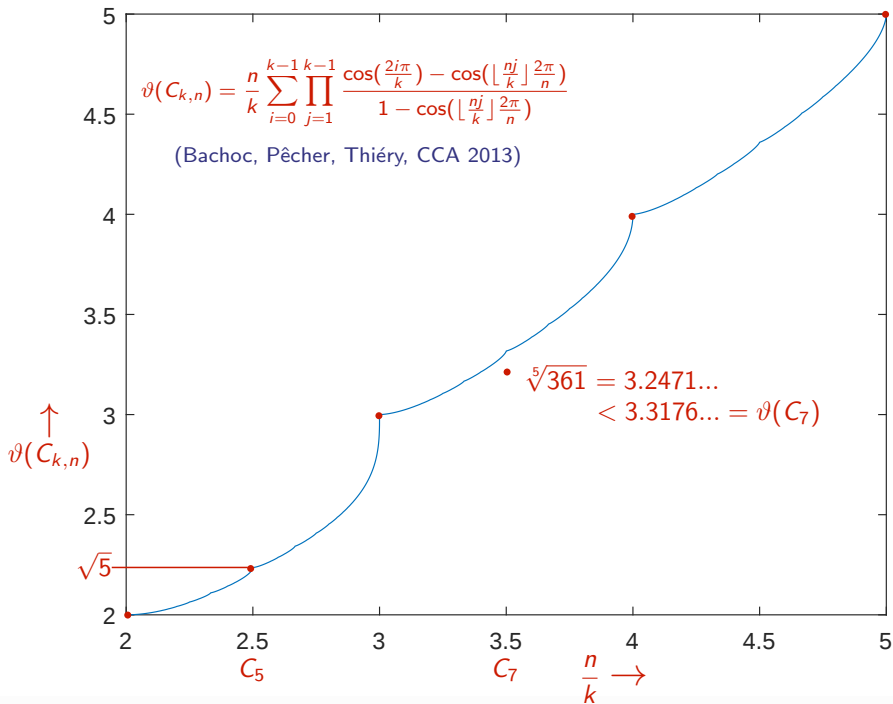


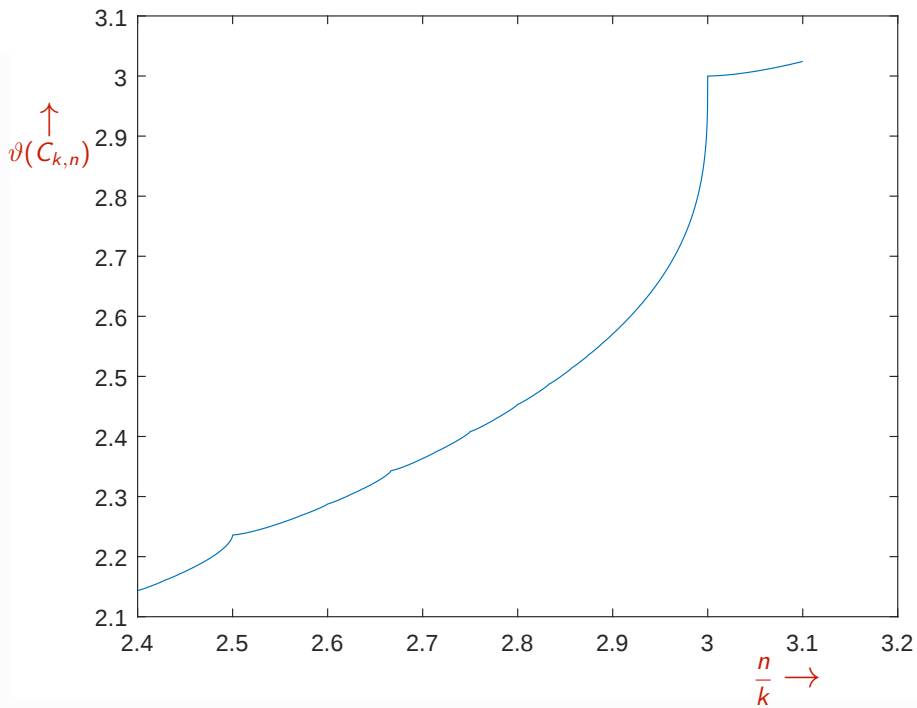


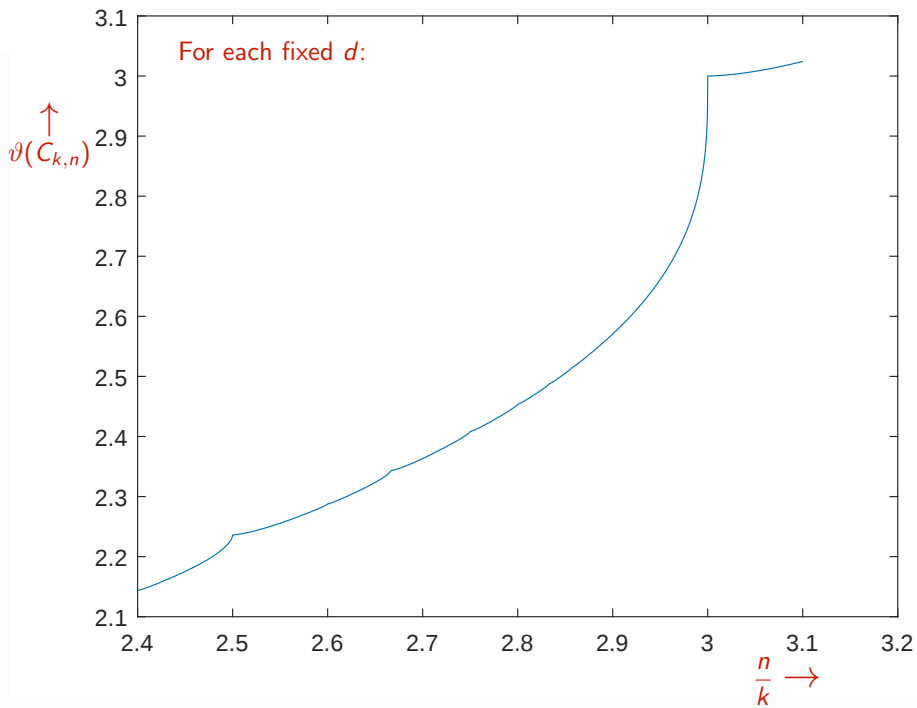


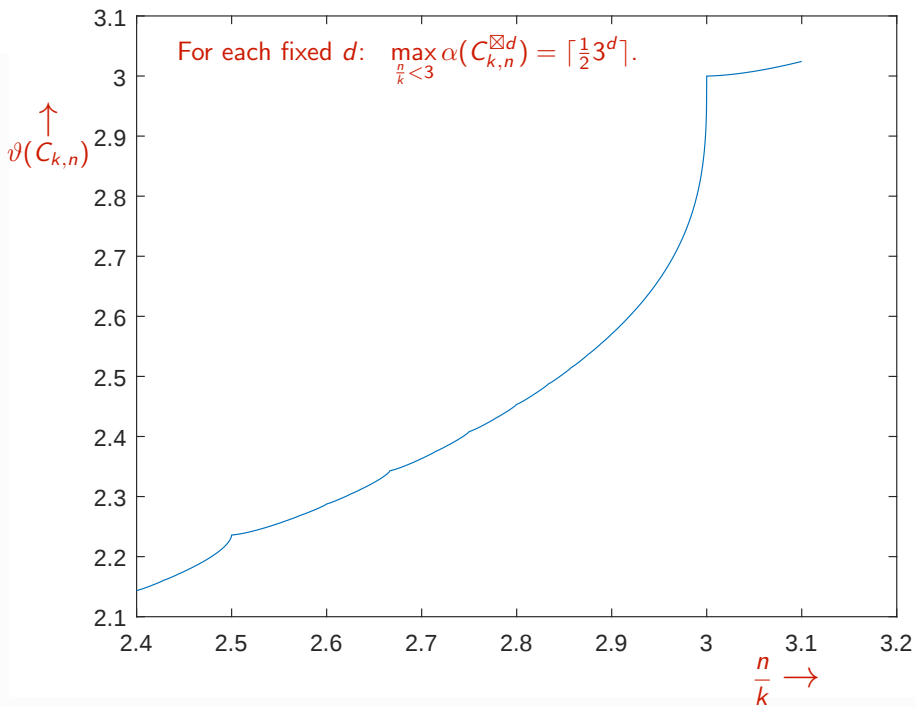






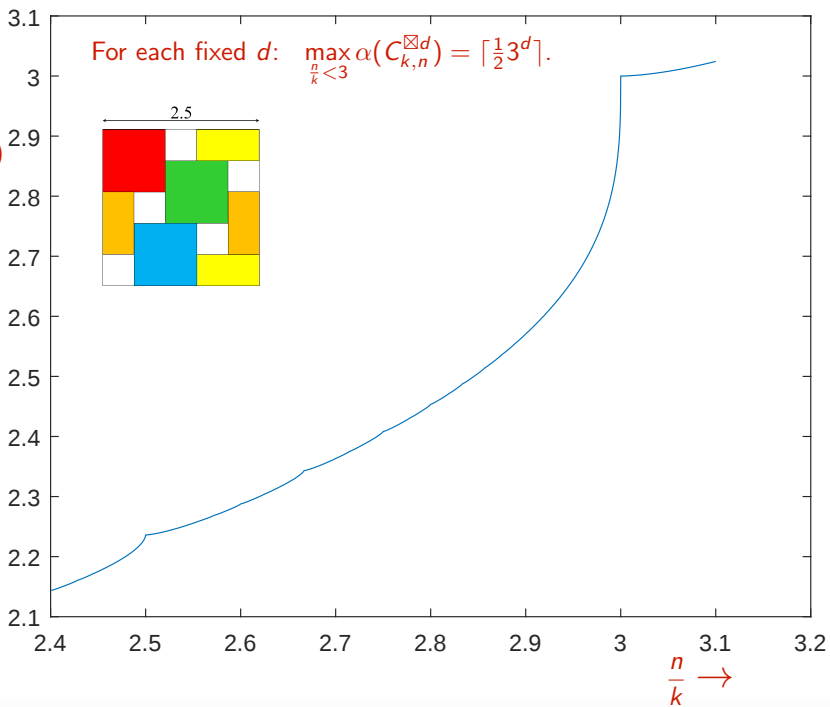
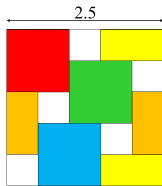


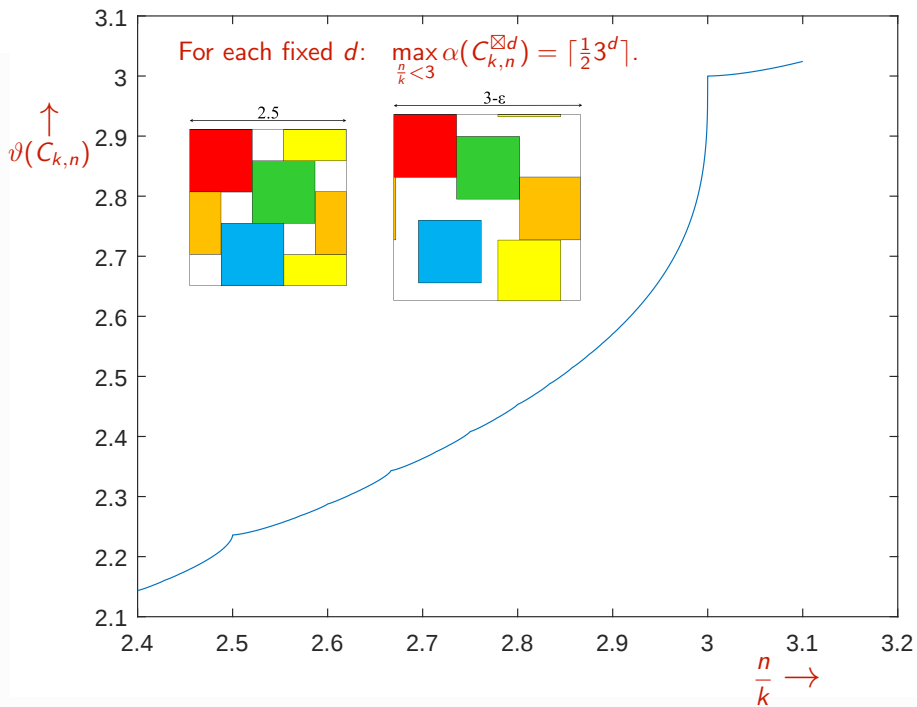


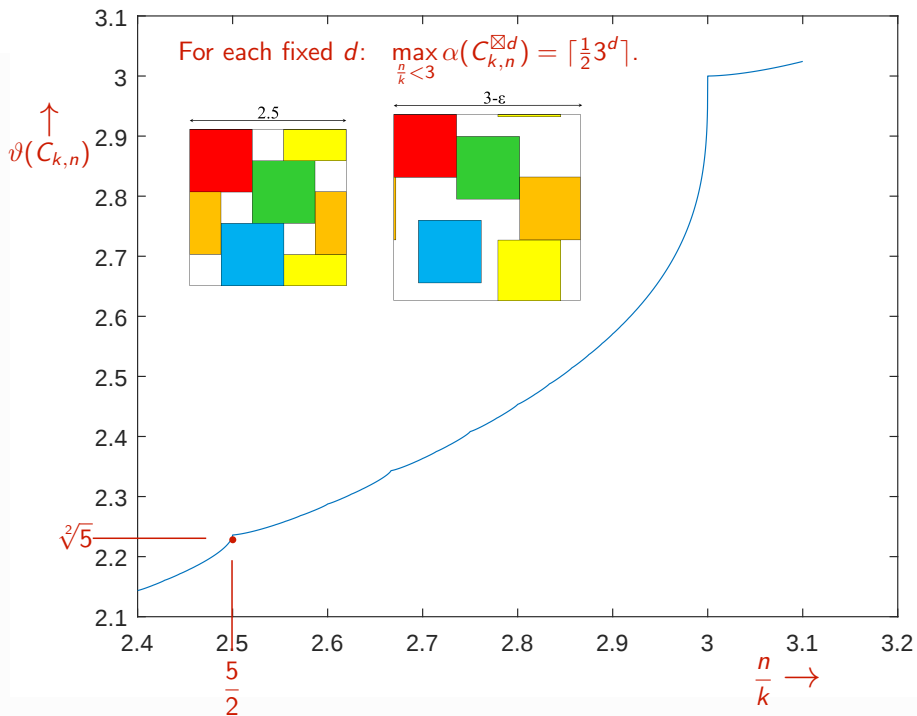


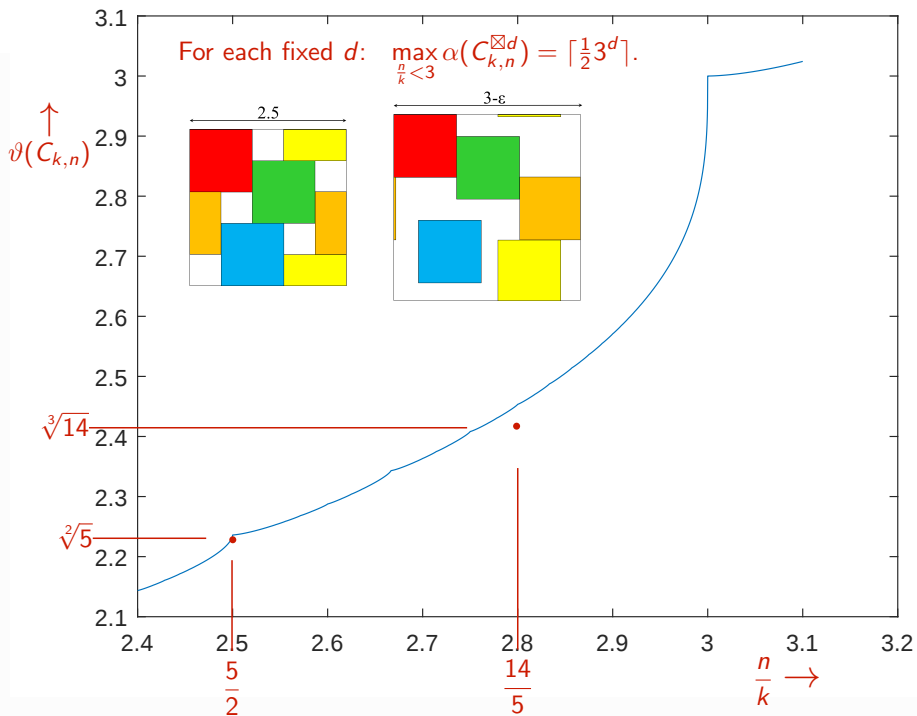
$\uparrow$
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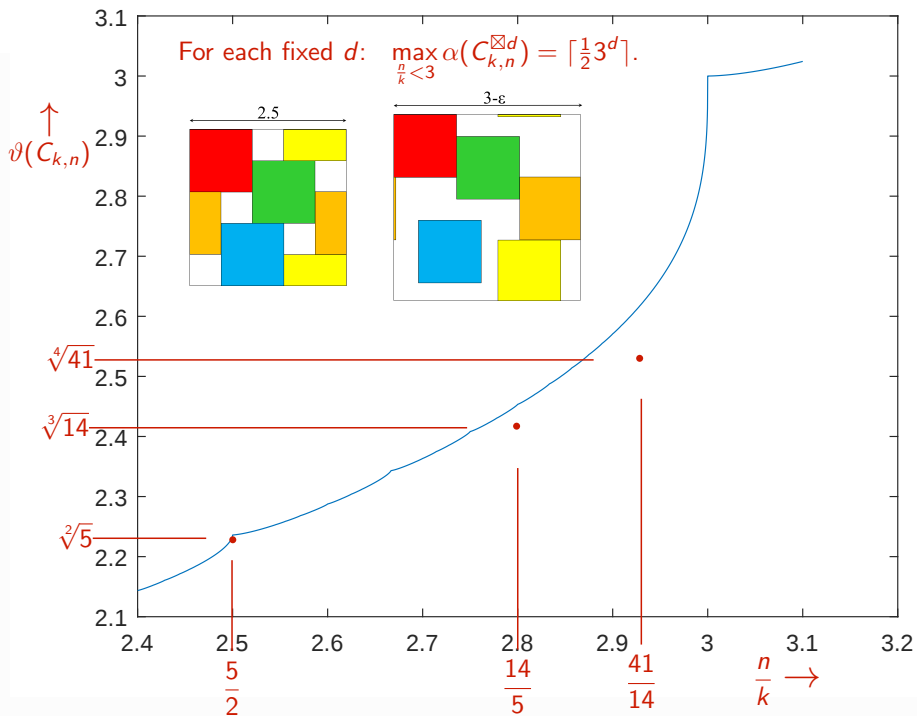
For each fixed d : $\max_{\frac{n}{k} < 3} \alpha(C_{k,n}^{\boxtimes d}) = \lceil \frac{1}{2} 3^d \rceil$.

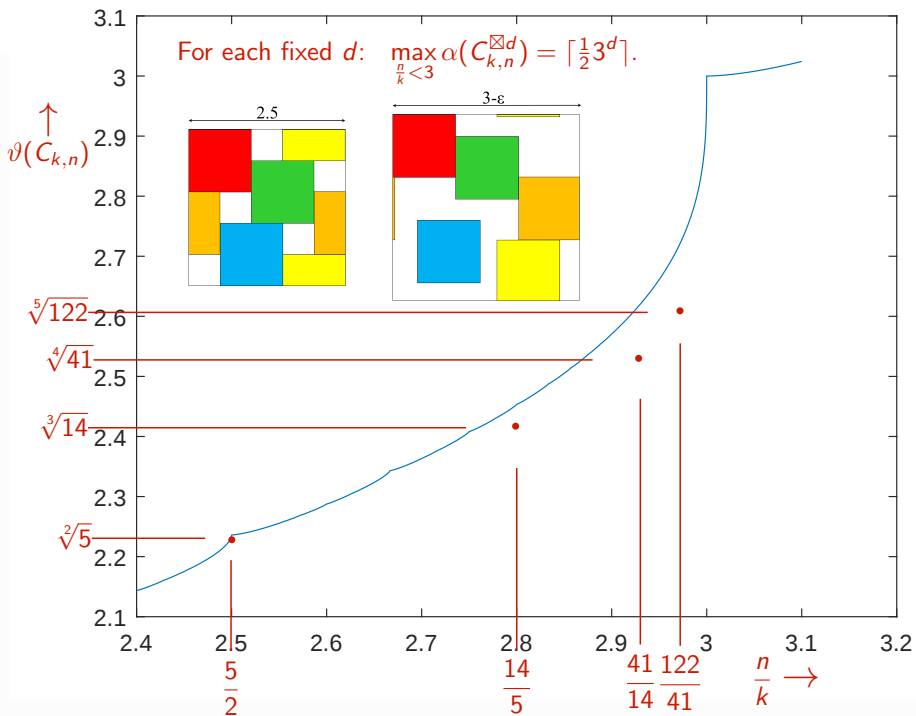


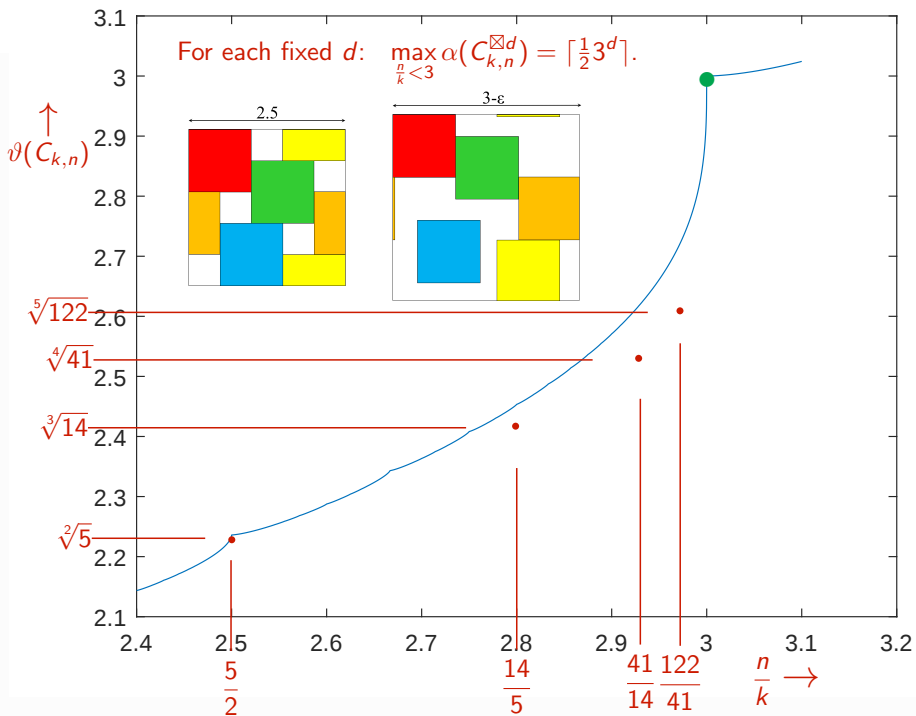












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Proved for $k \leq 11$.

And now for something completely different:

Conjecture: *Each simple, k -regular, properly k -edge-coloured graph contains a rainbow path of length $k - 1$.*

 all colours different

Proved for $k \leq 11$.

MERCI BIEN!