Higher coherence equations of semi-simplicial types as *n*-cubes of proofs

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A brief reminder

Semi-simplicial $(\Delta_+) =$ Simplicial $(\Delta) -$ Degeneracies

- Sets $X_0, X_1, X_2, ...$
- ► Face maps $d_i : X_n \to X_{n-1}$ for any $0 \le i \le n$
- Satisfying the semi-simplicial identity...

$$d_i d_j = d_{j-1} d_i$$
 when $i < j$

...And nothing more!

Straightforward to do in type theory by restricting ourselves to **h-Set**s

A brief reminder

Such a simple construction *should be* possible using **Type**s instead of **h-Set**s...right?

Problem

Define semi-simplicial types (SSTs) in HoTT

Raised during the Special Year on UF by Voevodsky and others

- **Types** in HoTT are *weak* ∞ -groupoids
- Constructions are no longer set-truncated
- Higher coherence issues!

The semi-simplicial identity $d_i d_j = d_{j-1} d_i$ induces "higher proof terms" that interfere with each other and need to be identified

Suppose we want to show that $d_i d_j d_k =_{X_n \to X_{n-3}} d_{k-2} d_{j-1} d_i$ in type theory

The semi-simplicial identity is the data of a term

$$\alpha_{i,j}: d_i d_j =_{X_n \to X_{n-2}} d_{j-1} d_i$$

There are two ways to compose the α_{i,j}'s together in order to inhabit the type

$$d_i d_j d_k =_{X_n \to X_{n-3}} d_{k-2} d_{j-1} d_i$$

This is given by

$$\pi \coloneqq \alpha_{j,k} \cdot \alpha_{i,k-1} \cdot \alpha_{i,j} \quad \text{and} \quad \pi' \coloneqq \alpha_{i,j} \cdot \alpha_{i,k} \cdot \alpha_{j-1,k-1}$$

• We now need the data of a term $\beta_{i,j,k}$: $\pi = \pi'$, and so on...

Possible approaches

▶ *n*-Truncate (e.g. **h-Set**s, **h-Grp**s, ...)

$$X: \Delta^{\operatorname{op}}_+ \to \mathbf{h}\operatorname{-Level}(n+2)$$

For any externally fixed n, stop the construction at stage n

$$\begin{array}{l} X_0 : \mathbf{Type} \\ X_1 : X_0 \to X_0 \to \mathbf{Type} \\ X_2 : \prod_{a \, b \, c \, : \, X_0} X_1(a, b) \to X_1(b, c) \to X_1(a, c) \to \mathbf{Type} \\ \vdots \\ X_n : \prod \cdots \to \mathbf{Type} \end{array}$$

 Add an "outer" equality which is strict (2LTT, Altenkirch, Capriotti, Kraus)

A general solution is believed to be impossible in "plain" HoTT

Ten years of investigations

2012-13 Special Year on UF

- 2012–13 Homotopy Type System (Voevodsky)
- 2014–15 Indexed + h-Sets (Herbelin)
 - 2015 Logic-enriched HoTT (Part, Luo)
- 2015–16 2-Level Type Theory (Altenkirch, Capriotti, Kraus)
 - 2022 MLTT + Type Streams (Kolomatskaia)
 - 2023 Indexed + h-Sets + Parametricity (Herbelin, Ramachandra)

Strict composition

TLCA 2015 - Warsaw, Poland

What if we write down all the coherence equations explicitly?

The explicit form taken by the higher-order coherence equations is well described in the case of (strict) ω -categories, i.e. when associativity and the exchange law hold on the nose for path composition

- Ross Street. The algebra of oriented simplexes. *Journal of Pure and Applied Algebra*, 49(3):283–335, 1987.
- Ian R. Aitchison. The geometry of oriented cubes. Macquarie University Research Report No: 86–0082, 1986.

n-Cubes of proofs

Our goal

Reformulate Aitchison's constructions in type theory, representing higher-order cubes as higher-order equality proofs

- Composition of *n* face maps corresponds to the "spine" of an *n*-cube
- ▶ (Higher) proof terms correspond to (composition of) *faces* of dimension ≥ 2
- The k-hemispheres are special cases of these compositions, as sources and targets of the equalities

How to give a recursive formulation of all the compositions involved?

3-Cube of proofs



 $\beta_{i,j,k}: \overbrace{\alpha_{j,k} \cdot \alpha_{i,k-1} \cdot \alpha_{i,j}}^{\mathcal{A}} =_{\left(d_i d_j d_k = (x_n \to x_{n-3})^d_{k-2} d_{j-1} d_i\right)} \underbrace{\alpha_{i,j} \cdot \alpha_{i,k} \cdot \alpha_{j-1,k-1}}_{\text{Frontmost}}$

Combinatorial structure

- ▶ When n ≥ 4 associativity and the exchange law are explicitly required
- There are nontrivial proof terms that are not equivalent to k-hemispheres...
- ...But are generated by hemispheres interfering at different levels!

Problem

How to describe the k-hemispheres?

- k-Hemispheres are made up of k-faces composed together in some order (not linear a priori!)
- We write $h_{k,n}^{\pm}$ for the *k*-hemispheres of the *n*-cube











Combinatorial structure

Definition

 D_k^n is the poset of *increasing* sequences $x_1 \dots x_n$ of length *n* with values $0 \le x_i \le k$ equipped with the "pointwise" order

Hasse diagram of
$$D_2^2$$
 $egin{array}{ccc} 00 \ \downarrow \ 01
ightarrow 11 \ \downarrow \ 02
ightarrow 12
ightarrow 22 \end{array}$

Theorem

The k-hemispheres of the n-cube are described by D_{n-k}^k in the sense that there is a bijection sending any $h_{k,n}^{\pm}$ onto a linear extension of D_{n-k}^k







Combinatorial structure

 D_k^n has a rich structure—so does $h_{k,n}^{\pm}$ as a result

There is (essentially) one canonical choice of linear extension of Dⁿ_k given by

$$x \preceq y \iff (x_n, x) \leq (y_n, y)$$

In our previous examples with D_2^2 , it chooses $11 \prec 02$

 \triangleright D_k^n can be recursively constructed with maps

$$d_*: D_{k-1}^n \to D_k^n$$
 and $R: D_k^{n-1} \to D_k^n$

Then $D_k^n = d_* D_{k-1}^n \amalg R D_k^{n-1}$

• This construction preserves \leq defined above!

Observation

$$h^\pm_{k,n}$$
 has a similar recursive formula

Extrusion of the 3-cube



Present and future work

Aitchison's report "only" contains formulae for the k-hemispheres

$$\psi_k[n] = \mu \psi_{k-1}[n-1] \cup \lambda \psi_k[n-1]$$
$$\omega_k[n] = \nu \omega_k[n-1] \cup \mu \omega_{k-1}[n-1]$$

- Give an explicit description of the "internal" compositions, which have a structure different from k-hemispheres (see Appendix)
- Make explicit the role of parenthesizing (associativity, exchange law, identity, ...)
- Combine these results to attempt a definition of semi-simplicial types with the coherence equations made explicit at every level

Questions?

Appendix



"Octagon of octagons" - Extracted from Aitchison's report

Appendix



Appendix

