

The Limit of Recursion in State-based Systems

Giacomo Barlucchi

j.w.w. B. Afshari & G.E. Leigh

Göteborgs Universitet

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μ -calculus

Modal logic with fixed points

$$\mu x \Box \perp \vee (q \wedge \Diamond x) \quad \nu x p \wedge \Diamond x \wedge \mu y (\bar{p} \vee y)$$

Frame

A frame $\mathcal{S} = (S, R, \Lambda)$ consisting of: a set S of states; a binary relation $R \subseteq S \times S$; a labelling function Λ .

$$f_\varphi: \mathcal{P}(S) \rightarrow \mathcal{P}(S)$$

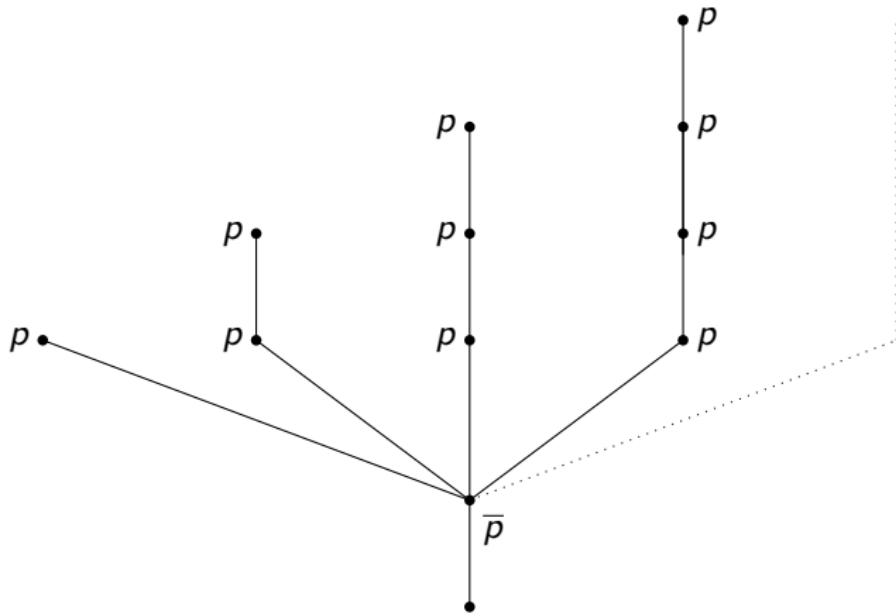
$$\|\mu x \varphi\|_{\mathcal{V}}^{\mathcal{S}} := \bigcap \{ U \subseteq S \mid \|\varphi\|_{\mathcal{V}[x \mapsto U]}^{\mathcal{S}} \subseteq U \}$$

$$\emptyset \subseteq f_\varphi(\emptyset) \subseteq f_\varphi(f_\varphi(\emptyset)) \subseteq \cdots \subseteq f_\varphi^\alpha(\emptyset) \subseteq f_\varphi^{\alpha+1}(\emptyset) \subseteq \cdots$$

$$\|\mu x \varphi\|_{\mathcal{V}}^{\mathcal{S}} := \bigcup_{\alpha} f_\varphi^\alpha(\emptyset)$$

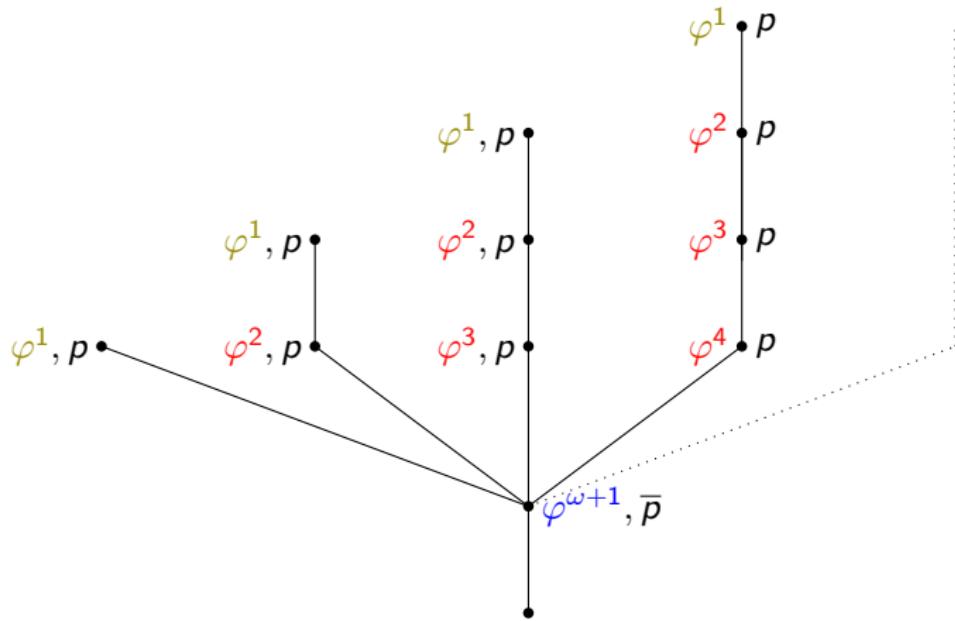
Approximation

$$\varphi \equiv \mu x. \square \perp \vee (p \wedge \square p \wedge \Diamond x) \vee (\bar{p} \wedge \square p \wedge \Box x)$$



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Closure Ordinal

Approximations

$$\varphi^\alpha := (\varphi, \alpha)$$

$$\|\mu x \varphi^\alpha\|_{\mathcal{V}}^{\mathcal{S}} := \bigcup_{\beta < \alpha} \|\varphi\|_{\mathcal{V}[x \mapsto \|\mu x \varphi^\beta\|_{\mathcal{V}}^{\mathcal{S}}]}^{\mathcal{S}}$$

$$\|(\psi_0 \vee \psi_1)^\alpha\| := \|\psi_0^\alpha\| \cup \|\psi_1^\alpha\|$$

$$\|(\psi_0 \wedge \psi_1)^\alpha\| := \|\psi_0^\alpha\| \cap \|\psi_1^\alpha\|$$

Closure

Given a frame \mathcal{S} and formula φ , $\text{CO}_{\mathcal{S}}(\varphi)$ is the least κ s.t.

$$\|\varphi^\kappa\|^{\mathcal{S}} = \|\varphi\|^{\mathcal{S}}$$

Closure Ordinal

$\text{CO}(\varphi)$ is the least ordinal κ such that for all countable frames \mathcal{S}

$$\text{CO}_{\mathcal{S}}(\varphi) \leq \kappa$$

Goal

Theorem

A countable ordinal α is the closure ordinal of a formula in the Σ -fragment iff $\alpha < \omega^2$

Σ -fragment

closure under logical connectives, modalities and μ -quantification of

- ▶ variables and literals,
- ▶ closed formulas in the language of μ -calculus.

$$\mu x. \Diamond x \wedge (\nu z. \Diamond(z \wedge p) \vee \Box \bar{p}) \quad \mu x. \Diamond x \wedge (\nu z. \Diamond(z \wedge p) \vee \Box x)$$

M.Czarnecki - 2010

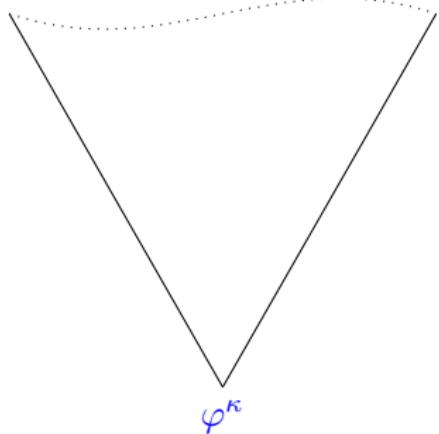
For every $\alpha < \omega^2$ there exists a Σ -formula φ such that $\text{CO}(\varphi) = \alpha$.

The argument

Theorem

If $\alpha < \omega_1$ is the closure ordinal of a Σ -formula φ then $\alpha < \omega^2$.

The argument for $\text{CO}(\varphi) > \kappa$



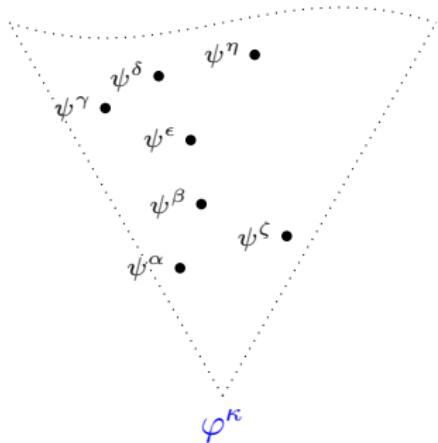
A assume \mathcal{S} maximal for φ^κ

The argument

Theorem

If $\alpha < \omega_1$ is the closure ordinal of a Σ -formula φ then $\alpha < \omega^2$.

The argument for $\text{CO}(\varphi) > \kappa$



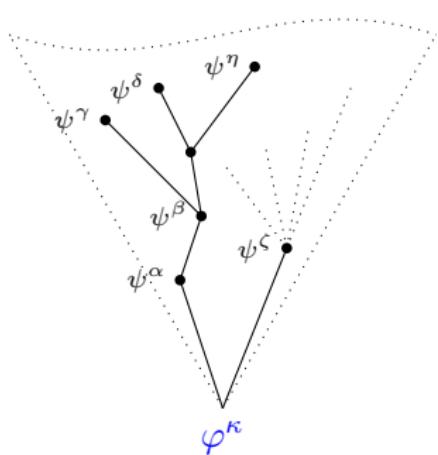
- A assume S maximal for φ^κ
- B annotate $\text{CL}(\varphi)$

The argument

Theorem

If $\alpha < \omega_1$ is the closure ordinal of a Σ -formula φ then $\alpha < \omega^2$.

The argument for $\text{CO}(\varphi) > \kappa$



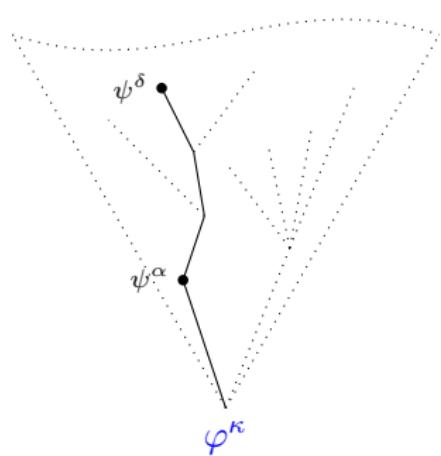
- A assume \mathcal{S} maximal for φ^κ
- B annotate $\text{CL}(\varphi)$
- C trace the ordinal progression

The argument

Theorem

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The argument for $\text{CO}(\varphi) > \kappa$



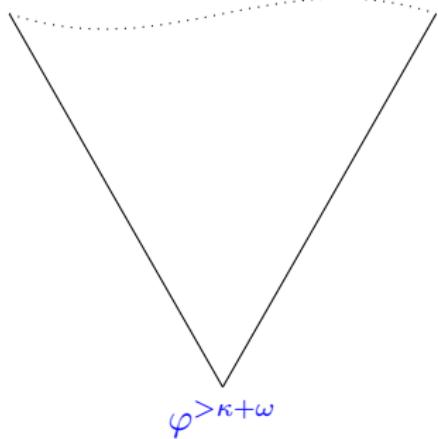
- A assume \mathcal{S} maximal for φ^κ
- B annotate $\text{CL}(\varphi)$
- C trace the ordinal progression
- D find repetitions

The argument

Theorem

If $\alpha < \omega_1$ is the closure ordinal of a Σ -formula φ then $\alpha < \omega^2$.

The argument for $\text{CO}(\varphi) > \kappa$



- A assume \mathcal{S} maximal for φ^κ
- B annotate $\text{CL}(\varphi)$
- C trace the ordinal progression
- D find repetitions
- E \mathcal{S} is not maximal

Syntax

$$\varphi := p \mid \bar{p} \mid x \mid \varphi \wedge \varphi \mid \varphi \vee \varphi \mid \nabla \Gamma \mid \mu x \varphi \mid \nu x \varphi$$

$$\nabla\{\varphi_0, \dots, \varphi_n\} := \bigvee_i^n \square \varphi_i \vee \Diamond \bigwedge_i^n \varphi_i$$

$$\square \varphi := \nabla\{\varphi, \perp\} \quad \Diamond \varphi := \nabla\{\varphi\} \wedge \nabla \emptyset$$

Conjunctive form φ

For every $\mu x \psi \in \text{CL}(\varphi)$, ψ is of the form

$$\bigwedge_{i < k} (\bigvee Z_i \vee \nabla \Gamma_i)$$

with Z_i closed formulas, and every formula in Γ_i is a variable, or is of the form $\mu y \psi'$.

Conservative Well-annotation

Well-annotation (Kozen -1988)

Given φ^κ , a *well-annotation* of \mathcal{S} is a function associating to each state $s \in S$ a set Θ_s of annotated formulas from $\text{CL}(\varphi)$ such that

1. if $\mu x^\alpha \psi \in \Theta_s$ then $\psi[\mu x \psi/x]^\beta \in \Theta_s$ for $\beta < \alpha$;
2. if $(\varphi \wedge \psi)^\alpha \in \Theta_s$ then $\varphi^\alpha, \psi^\alpha \in \Theta_s$;
3. ...

Conservative well-annotation

A well-annotation that is also minimal

- 1' if $\mu x^\alpha \psi \in \Theta_s$ then $\psi[\mu x \psi/x]^\beta \in \Theta_s$ for $\alpha = \beta + 1$;
- 2' ...

$O(\varphi, \Gamma)$

Given φ, Γ , $O(\varphi, \Gamma)$ is the supremum of $\kappa < \omega_1$ for which there exists a conservative Θ s.t. $\Gamma = \Theta_\rho$ and $\varphi^\kappa \in \Theta_\rho$.

Relevant part (1)

Relevant part

We call Φ a *relevant part* of Θ if for every $s \in S$,

1. $\Phi_s \subseteq \Theta_s$;
2. if $\mu x^{\alpha+1}.\psi \in \Phi_s$ then $\psi[\mu x.\psi/x]^{\alpha} \in \Phi_s$;
3. if $(\varphi_0 \wedge \varphi_1)^{\alpha} \in \Phi_s$ then exactly one $\varphi_i^{\alpha} \in \Phi_s$;
4. if $(\varphi_0 \vee \varphi_1)^{\alpha} \in \Phi_s$ then exactly one $\varphi_i^{\alpha} \in \Phi_s$;

$$\begin{aligned}\mu x. \bigwedge_i (\bigvee Z_i \vee \nabla \Gamma_i)^{\alpha+1} \\ \bigwedge_i (\bigvee Z_i \vee \nabla \Gamma_i)^{\alpha} \\ (\bigvee Z \vee \nabla \Gamma)^{\alpha} \\ \nabla \Gamma^{\alpha}\end{aligned}$$

Relevant part (2)

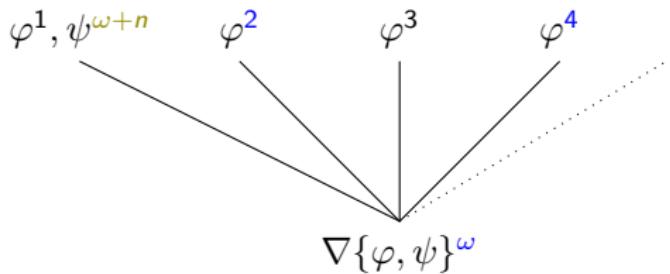
Relevant part

4 if $\nabla \Gamma^\alpha \in \Phi_s$ then:

- 4.1 one formula on every state where \Diamond
- 4.2 one/ ω formula/s for every \Box

Traces are:

- ▶ at most one ∇ per Φ_s
- ▶ possibly increasing
- ▶ exhaustive



$$\nabla\{\varphi, \psi\}^\omega = \Box\varphi^\omega \vee \Diamond(\varphi \wedge \psi)^{\omega+n}$$

Relevant part (2)

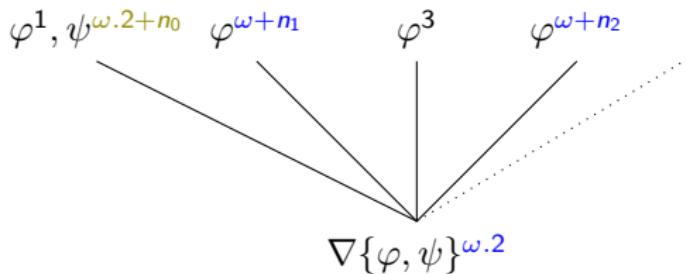
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$$\nabla\{\varphi, \psi\}^{\omega.2} = \Box\varphi^{\omega.2} \vee \Diamond(\varphi \wedge \psi)^{\omega.2+n_0}$$

Repetition and Optimality

Repetition pair

A pair of states (s, t) in \mathcal{S} is a *repetition pair* if:

1. there is a path $s = s_0 \dots s_k = t$,
2. $(\Theta_s^-, \Phi_s^-) = (\Theta_t^-, \Phi_t^-)$,
3. the unique relevant $\nabla\Gamma$ is s.t. $\nabla\Gamma^\alpha \in \Phi_s$ and $\nabla\Gamma^\beta \in \Phi_t$ for $\alpha > \beta$ limit ordinals.

Optimal state

Given a state s in a conservative well annotation Θ , Θ_s is *optimal* with respect to a formula $\varphi^\alpha \in \Theta_s$ if $O(\varphi, \Theta_s) < \alpha + \omega$.

Lemma

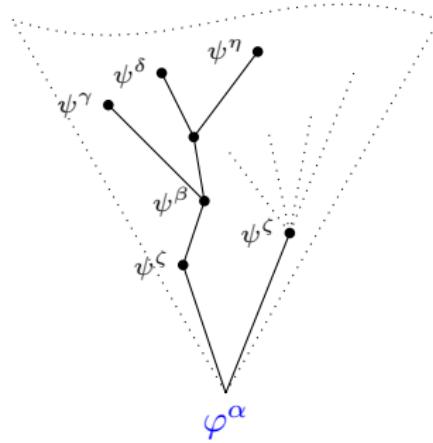
On every optimal path there are no repetition pairs.

Up to ω^2

First pumping lemma

For every φ and Γ there exists $N < \omega$ such that, if $O(\varphi, \Gamma) < \omega^2$ then $O(\varphi, \Gamma) \leq \omega \cdot N$.

- ▶ $N = 2^{2|\text{CL}(\varphi)|} + 1$ and $\omega^2 > \alpha > \omega \cdot N$
- ▶ Θ optimal for φ^α : $O(\varphi, \Theta_\rho) < \alpha + \omega$

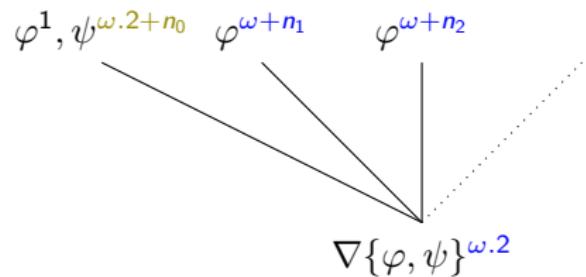


Up to ω^2

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- ▶ Θ optimal for φ^α : $O(\varphi, \Theta_\rho) < \alpha + \omega$
- ▶ an optimal relevant trace is strictly decreasing

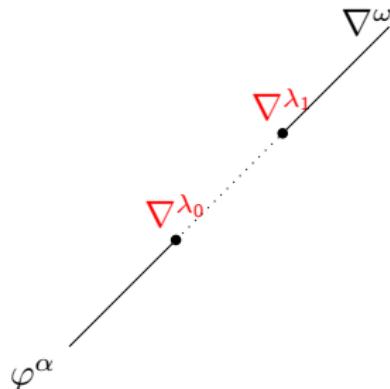


Up to ω^2

First pumping lemma

For every φ and Γ there exists $N < \omega$ such that, if $O(\varphi, \Gamma) < \omega^2$ then $O(\varphi, \Gamma) \leq \omega.N$.

- ▶ $N = 2^{2|\text{CL}(\varphi)|} + 1$ and $\omega^2 > \alpha > \omega.N$
- ▶ Θ optimal for φ^α :
 $O(\varphi, \Theta_\rho) < \alpha + \omega$
- ▶ an optimal relevant trace is strictly decreasing
- ▶ N limit points entail a repetition pair on the path
- ▶ contradiction by Lemma



Limits on CO

First pumping lemma

For every φ and Γ there exists $N < \omega$ such that, if $O(\varphi, \Gamma) < \omega^2$ then $O(\varphi, \Gamma) \leq \omega \cdot N$.

Second pumping lemma

For every φ and Γ , if $O(x, \Gamma) \leq \kappa$ then $O(x, \Gamma) < \omega^2$.

$$O(\varphi, \Gamma) = \kappa \implies CO(\varphi) \geq \kappa$$

$$CO(\varphi) = \kappa \implies O(\varphi, \Gamma') = \kappa$$

Theorem

A countable ordinal α is the closure ordinal of a formula in the Σ -fragment iff $\alpha < \omega^2$.

Thank you for your attention!

Definitions

Definition (Well-annotation)

An *annotation* of a frame \mathcal{S} for an equation system (X, E) is a function $\Theta: \mathcal{S} \rightarrow \mathcal{P}(\mathcal{L}_\mu \times \omega_1)$ associating to each state $s \in \mathcal{S}$ a set Θ_s of annotated formulas from $\text{CL}(X, E)$. A *well-annotation* of \mathcal{S} for (X, E) is an annotation Θ such that for all $s \in \mathcal{S}$, φ, Γ and $\alpha < \omega_1$:

1. if $\varphi^\alpha \in \Theta_s$ and $\varphi \in \mathcal{L}_\mu^-$ then $s \in \|\varphi\|^\mathcal{S}$;
2. if $x^\alpha \in \Theta_s$ for $x \in X$, then $E(x)^\beta \in \Theta_s$ for some $\beta < \alpha$;
3. if $\bigvee \Gamma^\alpha \in \Theta_s$ then $\Gamma^\beta \cap \Theta_s \neq \emptyset$ for some $\beta \leq \alpha$;
4. if $\bigwedge \Gamma^\alpha \in \Theta_s$ then $\Theta_s \preceq \Gamma^\alpha$;
5. if $\nabla \Gamma^\alpha \in \Theta_s$, then one of the following properties holds
 - 5.1 there exists $r \in R[s]$ such that $\Theta_r \preceq \Gamma^\alpha$,
 - 5.2 there exists $\varphi \in \Gamma$ such that $\Theta_r \preceq \{\varphi^\alpha\}$ for all $r \in R[s]$.

Definitions

Definition (Relevant part)

Let Θ be a conservative well-annotation of \mathcal{S} and Φ an annotation of the same frame. We call Φ a *relevant part* of Θ if for every $s \in \mathcal{S}$,

1. $\Phi_s \subseteq \Theta_s$;
2. if $x^\alpha \in \Phi_s$ then $E(x)^\beta \in \Phi_s$ where $\alpha = \beta + 1$;
3. if $\bigvee \Gamma^\alpha \in \Phi_s$ and $\alpha > 0$ then $\Gamma^\alpha \cap \Theta_s \subseteq \Phi_s$;
4. if $\bigwedge \Gamma^\alpha \in \Phi_s$ then $\chi^\alpha \in \Phi_s$ for exactly one $\chi \in \Gamma$;
5. if $\nabla \Gamma^\alpha \in \Phi_s$ and $\alpha > 0$ then:
 - 5.1 for all $\eta < \alpha$ and $\varphi \in \Gamma_\square^s$ there is a $r \in R[s]$ and $\beta > \eta$ s.t. $\varphi^\beta \in \Phi_r$,
 - 5.2 $\Gamma \cap \Phi_r^- \neq \emptyset$ for every $r \in \Gamma_\diamondsuit^s$, and
 - 5.3 for all $r \in R[s]$ if $\Phi_r \neq \emptyset$ then $y^\beta \in \Phi_r$ for some $y \in \Gamma$ and $\beta > p(\alpha)$.

Related works

1. Mathis Kretz (2006) - Proof-theoretic aspects of modal logic with fixed points.
2. G. Fontaine (2008) - Continuous Fragment of the mu-Calculus
3. M.J.Gouveia, L.Santocanale (2019) - \aleph_1 and the modal μ -calculus
4. G.C. Milanese (2018) - An exploration of closure ordinals in the modal μ -calculus.
5. M.Czarnecki (2010) - How fast can the fixpoints in modal μ -calculus be reached.
6. B.Afshari, G.E.Leigh (2013) - On closure ordinals for the modal mu-calculus