Combining fixpoint and differentiation theory FICS 2024

Zeinab Galal, Jean-Simon Pacaud Lemay

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In math and computer science

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- various notions of recursion (fixpoint operators, trace, ...)

Goal: general account of the interactions between fixpoints, trace and differentiation

Cartesian differential categories

Blute, R. F. and Cockett, J. R. B. and Seely, R. A. G. "Cartesian Differential Categories." *Theory and Applications of Categories*, 22(23):622–672 (2009).

A Cartesian differential category is a Cartesian left additive category $\mathbb C$ equipped with an operator D

$$\mathbf{D}: \mathbb{C}(A, B) \longrightarrow \mathbb{C}(A \times A, B)$$
$$f: A \to B \longmapsto \mathbf{D}[f]: A \times A \to B$$

satisfying seven axioms.

There is an associated term calculus

$$\mathbf{D}[f](a,b) \coloneqq \frac{\mathsf{d}f(x)}{\mathsf{d}x}(a) \cdot b$$

Chain rule for Cartesian differential categories

Chain rule axiom: for $f : A \rightarrow B$ and $g : B \rightarrow C$,

$$\frac{\mathrm{d}g(f(x))}{\mathrm{d}x}(a) \cdot b = \frac{\mathrm{d}g(y)}{\mathrm{d}y}(f(a)) \cdot \left(\frac{\mathrm{d}f(x)}{\mathrm{d}x}(a) \cdot b\right)$$

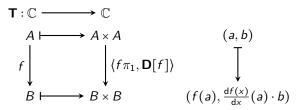
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$$A \times A \xrightarrow{\mathbf{D}[gf]} C = A \times A \xrightarrow{\mathbf{T}(f)} B \times B \xrightarrow{\mathbf{D}[g]} C$$

where \mathbf{T} is the tangent bundle



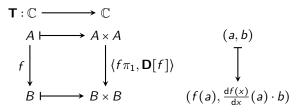
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$$A \times A \xrightarrow{\mathbf{D}[gt]} C = A \times A \xrightarrow{\mathbf{T}(t)} B \times B \xrightarrow{\mathbf{D}[g]} C$$

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T is a functor \Leftrightarrow the chain rule holds

Examples

Smooth functions between Euclidean spaces \mathbb{R}^n , for $f : \mathbb{R}^n \to \mathbb{R}^m$,

Derivative $\mathbf{D}[f]: \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}^m$ $\xrightarrow{\text{curry}}$ Jacobian $\mathbf{J}[f]: \mathbb{R}^n \to \text{Lin}(\mathbb{R}^n, \mathbb{R}^m)$

Relational and weighted relational models

- Reduced power series
- Convenient vector spaces
- Any model of the differential λ-calculus (Cartesian closed differential categories)
- And many more...

Fixpoint operators

Parametrized fixpoint operators are fixpoints for morphisms in context:

$$f: \Gamma \times X \longrightarrow X \qquad \rightsquigarrow \qquad \mu x. f(-, x): \Gamma \to X$$

with $\mu x. f(a, x) = f(a, \mu x. f(a, x))$ for all $a \in \Gamma$

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• Example: in PCF, we have an operator **fix**

$$\frac{\Gamma, X \vdash M : X}{\Gamma \vdash \operatorname{fix} M : X} \operatorname{FIX}$$

with reduction fix $M \rightarrow M($ fix M)

Categorical fixpoint operators

- Bloom, S. L., Ésik Z. Iteration theories, 1993.
- Simpson, A. and Plotkin, G. Complete axioms for categorical fixed-point operators, 2000.

For a Cartesian category $\mathbb{C},$ a parametrized fixpoint operator on \mathbb{C} is family of functions

$$\mathbf{fix}_X^{\Gamma} : \mathbb{C}(\Gamma \times X, X) \longrightarrow \mathbb{C}(\Gamma, X)$$
$$f \longmapsto \mathbf{fix}_X^{\Gamma}(f)$$

indexed by pairs of objects in $\ensuremath{\mathbb{C}}$ verifying:

• fixpoint axiom: for all $f: \Gamma \times X \to X$,

$$f \circ \langle \mathrm{id}_{\Gamma}, \mathbf{fix}(f) \rangle = \mathbf{fix}(f)$$

• naturality axiom: for all morphisms $g : \Gamma \to \Delta$ and $f : \Delta \times X \to X$,

$$\mathbf{fix}(f) \circ g = \mathbf{fix}(f \circ (g \times \mathrm{id}_X)).$$

Examples

• for domains Γ, X and a Scott-continuous map $f : \Gamma \times X \to X$,

$$f(a,\bigvee_{n\in\omega}f^n(a,\bot))=\bigvee_{n\in\omega}f^n(a,\bot)$$

is a least parametrized fixpoint for f.

- co-Kleisli category (weighted) relations with the finite multisets comonad
- special case: power series over continuous semi-rings
- any cpo-enriched cartesian closed category
- complete metric spaces and contractive maps
- and many more ...

Additional axioms

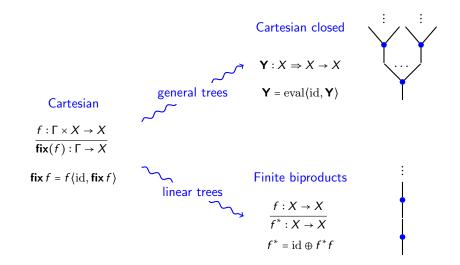
A fixpoint operator fix is a Conway operator if it satisfies

• Dinaturality axiom: for $f : \Gamma \times X \to Y$ and $g : Y \to X$

$$\mu x.g(f(a,x)) = g(\mu y.f(a,g(y)))$$

Bekić axiom: for f : Γ × X × Y → X and g : Γ × X × Y → Y
 μ(x,y). (f(a,x,y), g(a,x,y)) =
 (μx.f (a,x,μy.g(a,x,y)), μy.g (a, μx.f (a,x,μy.g(a,x,y)), y))

Conway fixpoint operator \Leftrightarrow Trace operator with \times as tensor



The function $f : \mathbb{R} \times \mathbb{R} \to \mathbb{R}$ has a parametrized fixpoint $g : \mathbb{R} \to \mathbb{R}$ $(x, y) \mapsto 1 - x^2 y$ $x \mapsto \frac{1}{1 + x^2}$ g(x) = f(x, g(x)) (1)

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By the chain rule,

$$g'(x) = \frac{\partial f}{\partial x}(x,g(x)) + \frac{\partial f}{\partial y}(x,g(x)) \cdot g'(x) = \frac{-2x}{(1+x^2)^2}$$
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Another way of computing g'(x):

• Compute the tangent of f, $\mathbf{T}(f) : \mathbb{R} \times \mathbb{R} \times \mathbb{R} \times \mathbb{R} \to \mathbb{R} \times \mathbb{R}$:

$$\mathbf{T}(f)(x, y, a, b) = \left(f(x, y), \frac{\partial f}{\partial x}(x, y) \cdot a + \frac{\partial f}{\partial y}(x, y) \cdot b\right)$$

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Cartesian differential fixpoint categories

Computing the derivative of the fixpoint is equivalent to computing the fixpoint of the tangent.

Option 1 differentiating the fixpoint

$$\frac{\overbrace{\Gamma \times X \xrightarrow{f} X}{\Gamma \xrightarrow{\text{fix}(f)} X}}{\Gamma \times \Gamma \xrightarrow{\text{D[fix}(f)]} X}$$

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Option 2 fixpoint of the tangent

$$\frac{\Gamma \times X \xrightarrow{f} X}{\Gamma \times X \times \Gamma \times X \xrightarrow{\mathsf{T}(f)} X \times X} \\
\frac{\Gamma \times \Gamma \times X \times X \xrightarrow{c} \Gamma \times X \times \Gamma \times X \xrightarrow{\mathsf{T}(f)} X \times X}{\Gamma \times \Gamma \times X \xrightarrow{c} \Gamma \times X \times \Gamma \times X \xrightarrow{\mathsf{T}(f)} X \times X}$$

where $c := \operatorname{id}_{\Gamma} \times \langle \pi_2, \pi_1 \rangle \times \operatorname{id}_X : \Gamma \times \Gamma \times X \times X \xrightarrow{\cong} \Gamma \times X \times \Gamma \times X$

Cartesian differential fixpoint categories

For a Cartesian differential category $\mathbb{C},$ a parametrized fixpoint operator fix satisfies the

• differential-fixpoint axiom if for every $f : \Gamma \times X \to X$

$$\mathbf{D}[\mathbf{fix}(f)] = \pi_2 \mathbf{fix}(\mathbf{T}(f \circ c))$$

• **tangent-fixpoint axiom** if for every $f : \Gamma \times X \to X$

$$\mathbf{T}[\mathbf{fix}(f)] = \mathbf{fix}(\mathbf{T}(f \circ c))$$

Lemma

If the fixpoint operator fix is Conway, the following are equivalent:

- **fix** satisfies the differential-fixpoint rule;
- **fix** satisfies the tangent-fixpoint rule.

Examples

- Differential fixpoint axiom: any category with finite biproducts where D[f] = f ο π₂ provided that the fixpoint is uniform wrt projection maps
- Tangent fixpoint axiom: weighted relations
- Special case: formal power series over continuous semi-rings
- General case: quantale enriched profunctors with the free exponential
- More generally: any Cartesian differential category where the fixpoint operator is obtained from bifree algebras and is uniform wrt linear maps.

Cartesian closed differential fixpoint categories

Ehrhard, T. "A coherent differential PCF." *Logical Methods in Computer Science*, 19 (2023).

For a Cartesian closed differential category \mathbb{C} , a fixpoint combinator $\mathbf{Y}: X \Rightarrow X \to X$ satisfies the

differential-fixpoint axiom if

$$\mathbf{D}[\mathbf{Y}] = \pi_2 \mathbf{Y} \lambda(\mathbf{T}(\text{eval}) \circ \boldsymbol{c})$$

tangent-fixpoint axiom if

$$\mathbf{T}[\mathbf{Y}] = \mathbf{Y}\lambda(\mathbf{T}(\text{eval}) \circ c)$$

Examples cpo-enriched cartesian closed categories, categories with fixpoint objects

In a Cartesian differential category \mathbb{C} , a morphism $f: X \to Y$ is linear if

$$\mathbf{D}[f] = X \times X \xrightarrow{\pi_2} X \xrightarrow{f} Y$$

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Lemma

If \mathbb{C} is a Cartesian differential fixpoint category, then the fixpoint of a linear morphism is linear.

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Lemma

If $\mathbb C$ is a Cartesian differential fixpoint category, then the fixpoint of a linear morphism is linear.

- Differential fixpoint axiom \Rightarrow Lin(\mathbb{C}) has a fixpoint operator.
- Differential fixpoint axiom \Rightarrow Lin(\mathbb{C}) has a repetition operator. + Conway fixpoint

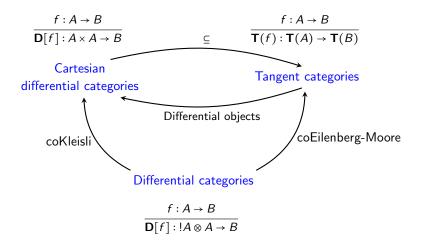
Inspired by:

- Combinatorics: Newton's method for combinatorial species (Decoste, Labelle, Leroux, Pivoteau, Salvy, Soria, ...)
- Formal languages and dataflow analysis: Newton's method for power series over ω-continuous semirings (Esparza, Kiefer, Luttenberger, Schlund,...)

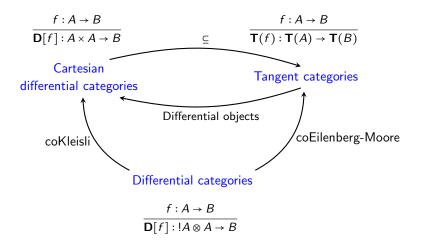
Main idea: assuming the differential-fixpoint axiom

compute the non-linear using the derivative and the linear repetition operator

What about differential categories?



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A natural notion of fixpoint for differential categories a trace for the monoidal product \otimes , how should it interact with differentiation?

- More general iterative schemes than Newton-Raphson
- Approximation of solutions of differential equations
- Reverse differential categories and fixpoints
- Restriction differential category for local implicit function theorems
- Guarded recursion and differentiation for fixpoints in metric models
- Syntactic counterpart (Taylor expansion)
- Differentiation in the coinductive world (e.g. stream calculus)

Thank you