

# Diagrammatic sets: weak higher categories for rewriting

Amar Hadzihasanovic

Formerly: RIMS, Kyoto University

Currently: IRIF, Université de Paris

TallCat Seminar

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There is a paper. But I'm reworking it heavily.  
Read at your own risk.

# Higher categories for all

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↪ **complete Segal spaces, complicial sets...** pick your favourite.

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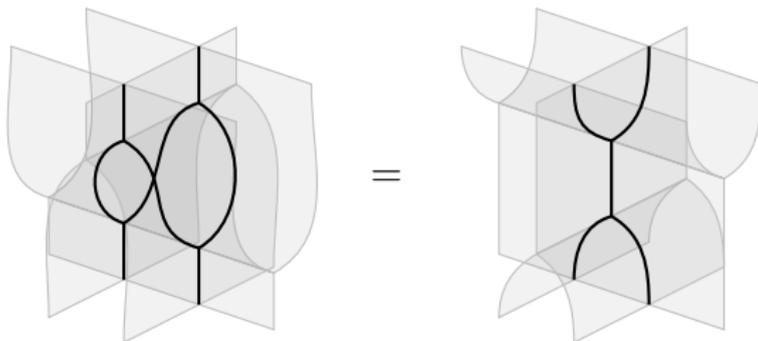
- We love diagrams! We love *presented* monoidal categories.
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- (*Then higher dimensions appear*) \*panic\*

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*How do we interpret this?*



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There is a lack of pasting theorems  
for mainstream models of **weak** higher categories.

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# The golden age of strict $\omega$ -categories

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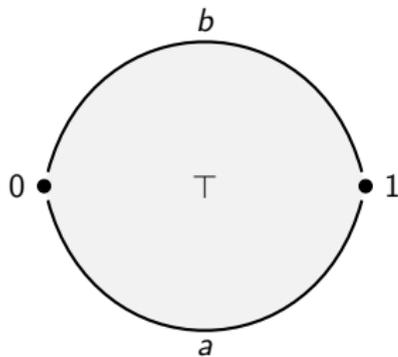
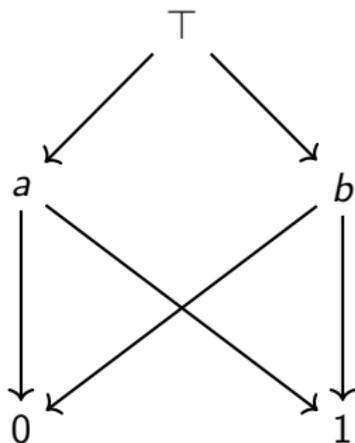
- **1987**: Ross Street's *The algebra of oriented simplexes* is out, sparking an interest in the combinatorics of higher-dimensional categorical diagrams.

Then several works on the combinatorics of *pasting diagrams* and their *pasting theorems* in strict  $n$ -categories:

- **1988**: John Power
- **1989**: Michael Johnson
- **1991**: Ross Street, John Power
- **1993**: Richard Steiner

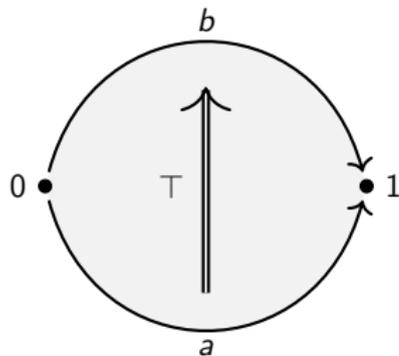
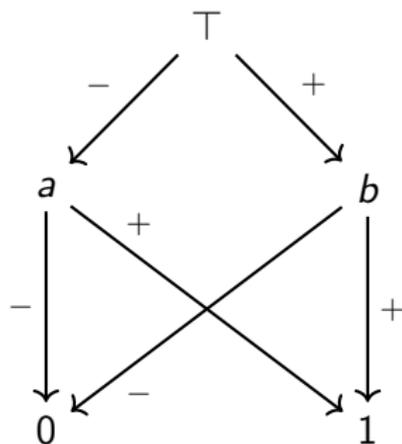
# Steiner's directed complexes

We can associate to a cell complex its **face poset**...



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and to a pasting diagram its **oriented face poset**.

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Steiner 1993, *The algebra of directed complexes*, gives sufficient conditions for

- an oriented poset to be the oriented face poset of a pasting diagram, and
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- an oriented poset to be the oriented face poset of a pasting diagram, and
- the pasting diagram to be **reconstructed** from its oriented face poset.

Many oriented posets present  $\omega$ -categories —  
fewer present **polygraphs**,  
that is,  
 $\omega$ -categories that are **freely generated** by some of their cells.

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If  $P$  has  $\dim n$  and  $Q$  has  $\dim k$ ,  $P \boxtimes Q$  has  $\dim n + k$ .

A variant of this was used to define  
the lax Gray product of  $\omega$ -categories  
(Steiner 2004, Ara-Maltsiniotis 2017)

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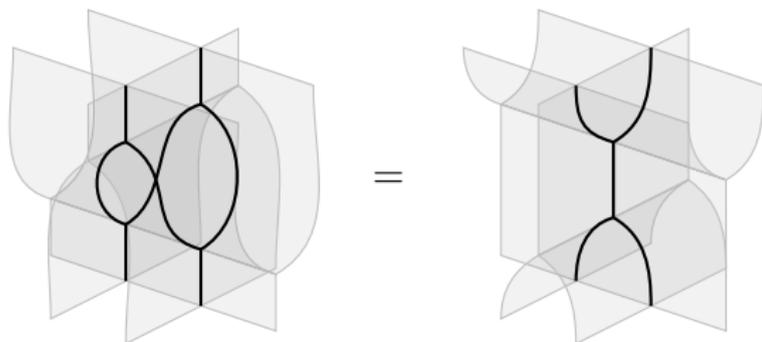
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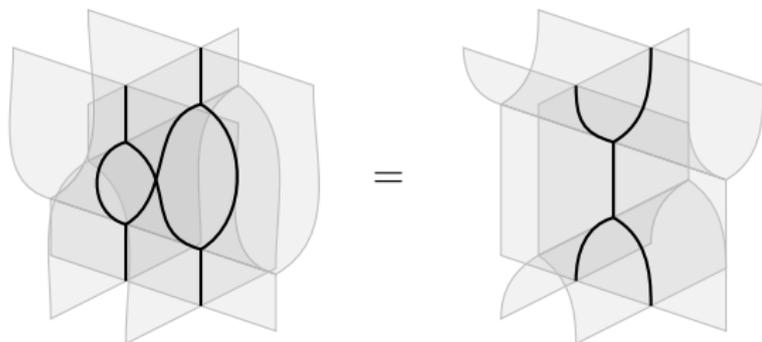
# Lax Gray products and diagrammatic algebra



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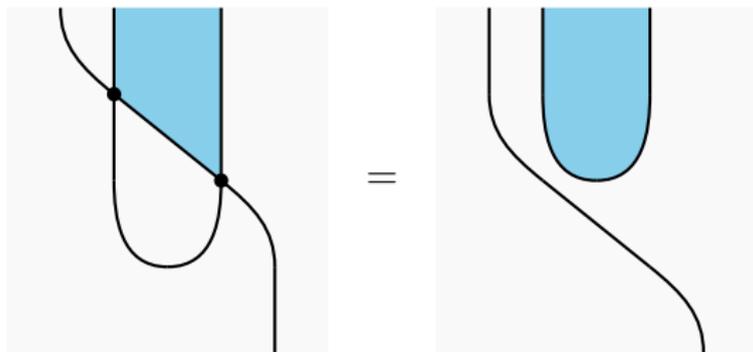
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(I'm not the only one)

# Lax Gray products and diagrammatic algebra

## Example: Biunitary equations

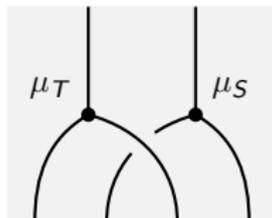
Used by Jamie Vicary and Mike Stay to unify quantum and encrypted communication protocols. They are models of a lax Gray product of 2-categories.



# Lax Gray products and diagrammatic algebra

## Example: Distributive laws of monads

They are models in **Cat** of a lax Gray product of 2-categories.



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Morally this should be a *braided monoidal category*.

But in strict  $\omega$ -categories, it is a *commutative monoidal category*.

**This breaks everything.**

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The core of the argument relies on the fact that “doubly monoidal” degenerates to “commutative” in strict 3-categories (**strict Eckmann-Hilton**).

...still contained some good ideas

Good takeaway #1 from Kapranov-Voevodsky:

*homotopy types may have **semi**strict algebraic models  
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- **2018**: Henry proves the homotopy hypothesis for “regular  $\omega$ -groupoids”.

# Diagrams with spherical boundary

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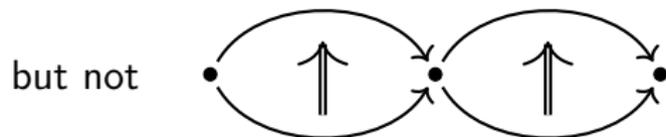
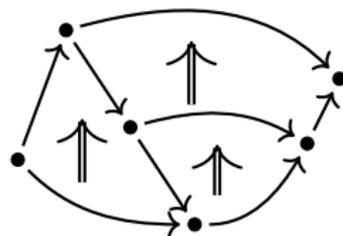
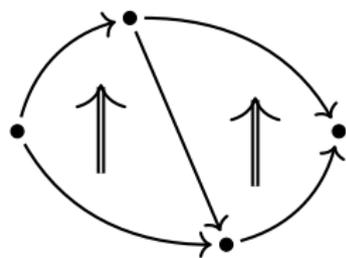
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~ “are homeomorphic to  $n$ -balls”

# Diagrams with spherical boundary



...and more good ideas

Good takeaway #2 from Kapranov-Voevodsky:

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Good takeaway #2 from Kapranov-Voevodsky:

**Diagrammatic sets**

Good takeaway #2 from Kapranov-Voevodsky:

## Diagrammatic sets

Kapranov-Voevodsky pass from spaces to  $\omega$ -categories through an intermediate notion of “spaces locally modelled on combinatorial pasting diagrams”, they call diagrammatic sets.

# Diagrammatic sets

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(Work in progress: axiomatic approach relative to "nice classes of diagrams")

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This is a model where we can interpret *every regular diagram* and compose *every diagram with spherical boundary*.

Just “stuff” any diagram with units and it will become regular!

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This started a slowly rising French school of rewriting with polygraphs (Yves Lafont, Philippe Malbos, Yves Guiraud, Samuel Mimram...)

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which is why I am in Paris now

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Many of the core ideas in polygraphic rewriting  
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*polygraphs* and *CW complexes*,  
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This analogy is limited by  
the fact that strict  $\omega$ -categories do not model all spaces.

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Everything that can be done with polygraphs  
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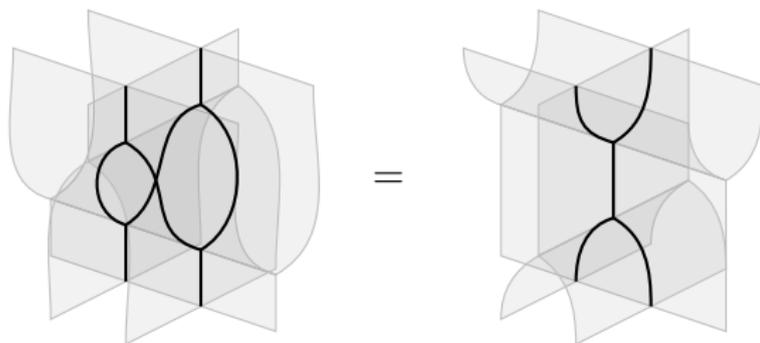
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- 4 Diagrams can be interpreted in models of all homotopy types, for rewriting homotopies
- 5 Lax Gray products, *joins* are easily defined and computed

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The smash product of diagrammatic sets produces this equation, the way it should.

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My hope is that diagrammatic sets can make the link between rewriting and homotopy theory tighter, on our way to figuring out what the right notions are.

*Work in progress:*

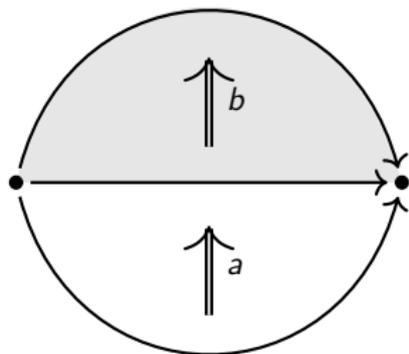
a model of computation in diagrammatic sets  
based on a “directed homotopy extension property”.

⋮  
2014  
2015  
2016  
2017  
2018  
2019  
**2020**

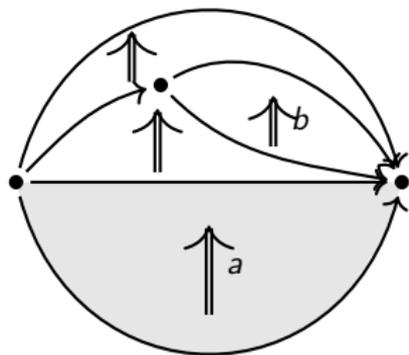
Thanks for listening!



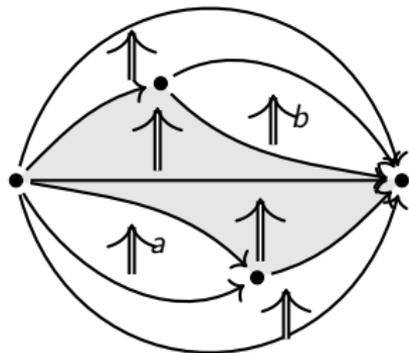
# Eckmann-Hilton in diagrammatic sets



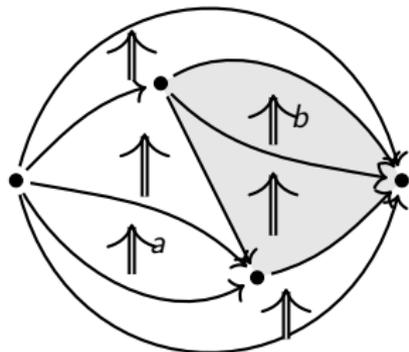
# Eckmann-Hilton in diagrammatic sets



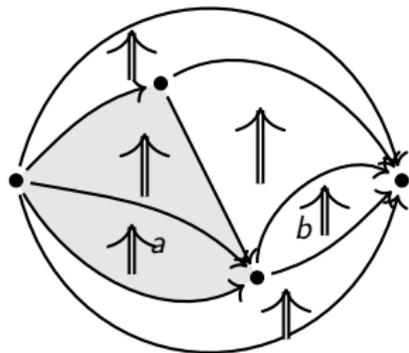
# Eckmann-Hilton in diagrammatic sets



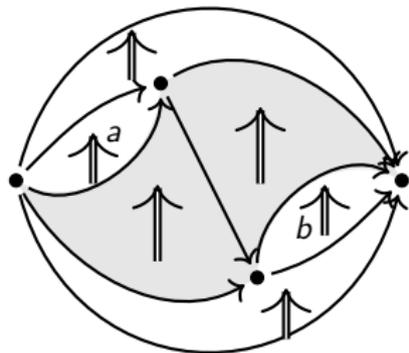
# Eckmann-Hilton in diagrammatic sets



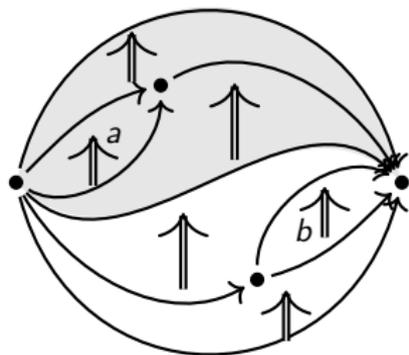
# Eckmann-Hilton in diagrammatic sets



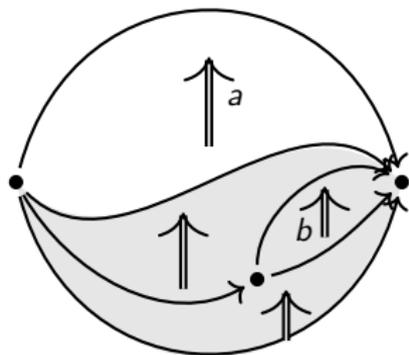
# Eckmann-Hilton in diagrammatic sets



# Eckmann-Hilton in diagrammatic sets



# Eckmann-Hilton in diagrammatic sets



# Eckmann-Hilton in diagrammatic sets

