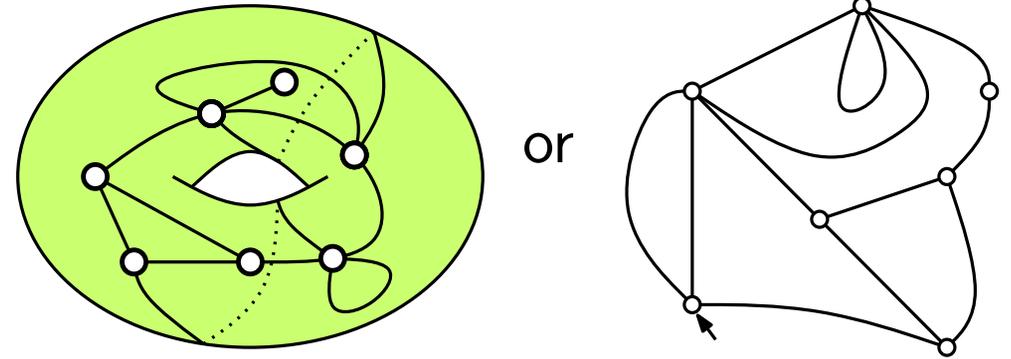


An introduction to map enumeration

Guillaume Chapuy, LIAFA, CNRS & Université Paris Diderot

What this course is about (I)

A **map** is a graph embedded in a surface:



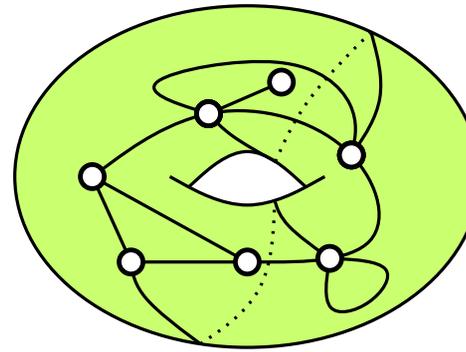
Maps appear (almost?) everywhere in mathematics.

Map enumeration alone is an enormous area.

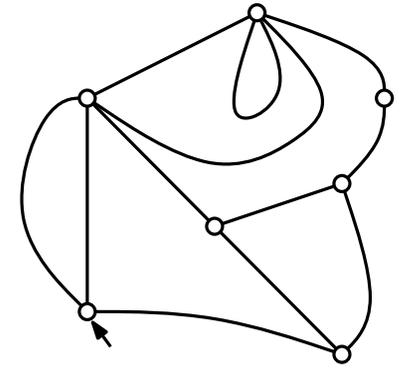
At least **three reasons** to be interested in it:

- because you would like to start **working seriously on the subject**.
- because map enumeration contains **powerful tools** that can be useful to other parts of combinatorics (functional equations, bijective tricks, algebraic tools...).
- because it is likely that **your favorite subject** is linked to map enumeration in at least some special case: looking at where this problem appears in the world of maps is a **very good source of new questions**.

What this course is about (II)

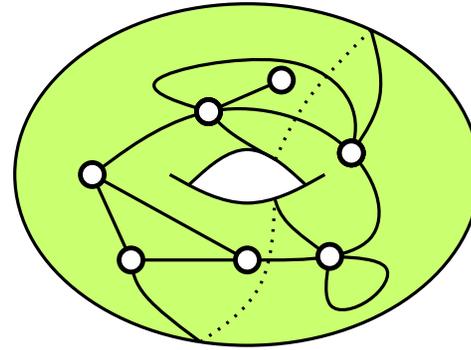


or

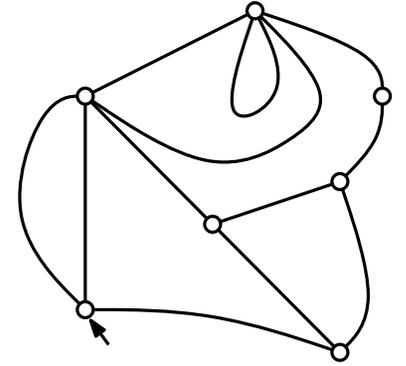


Topics **not covered**: - link with algebraic geometry
matrix integrals
string theory

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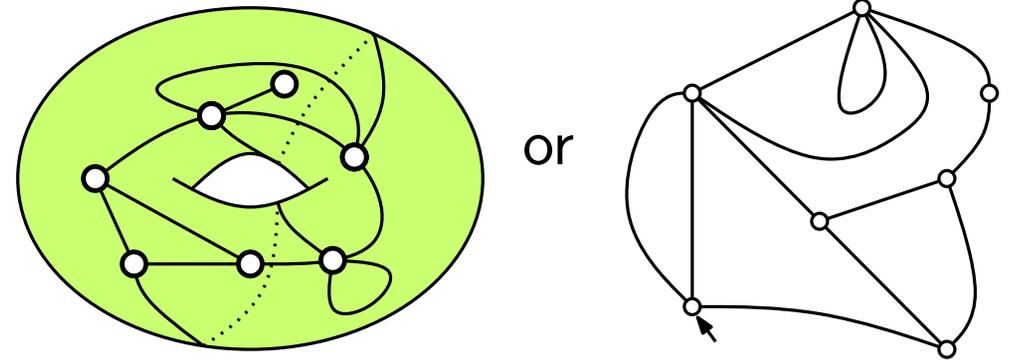


or



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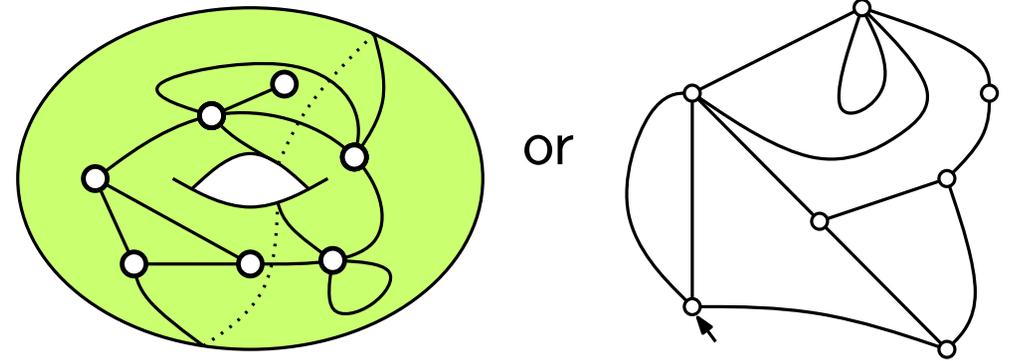


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SEE THE BOOK BY
LANDO-ZVONKIN

SEE E.G. SURVEY BY
MIERMONT

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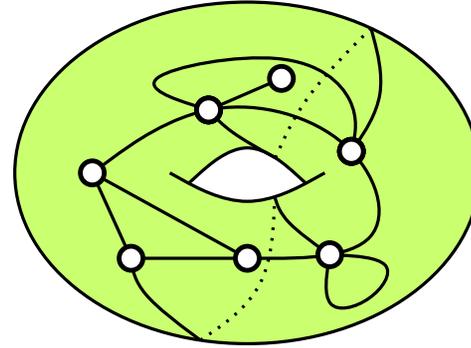
Topics **covered**: I – Maps
II – Tutte equation, counting planar maps
III – Tutte equation, counting maps on general surfaces
IV – Bijective counting of maps

The exercises contain **entry points to other subjects**
(one-face maps, link with the symmetric group)

Lecture I – What is a map? (the oral tradition)

Rough definition

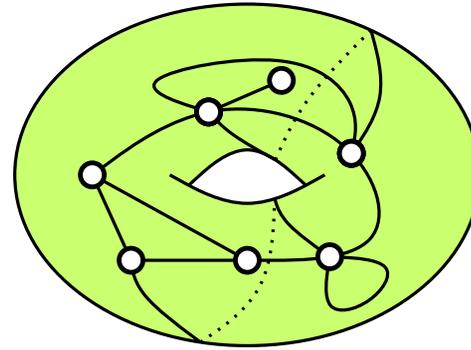
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Rough definition

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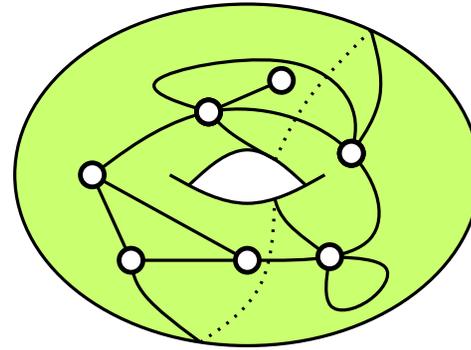


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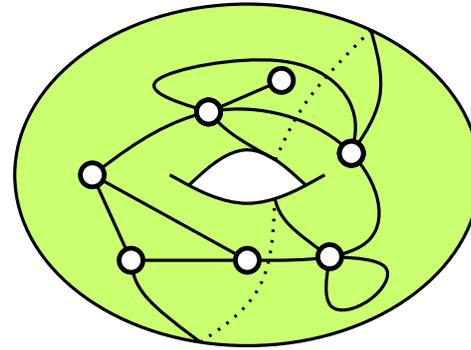
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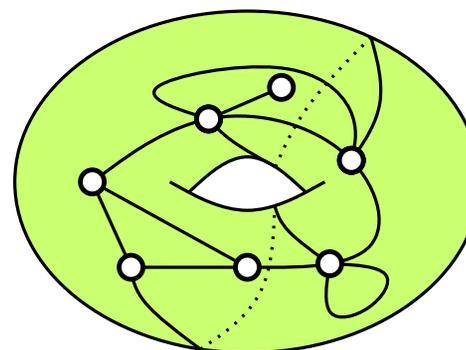
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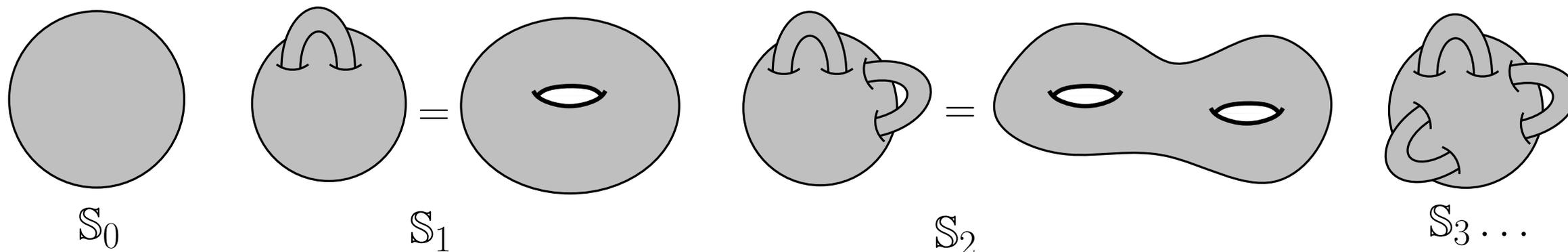
A **map** is a **graph drawn on a surface**.

? ? ?



A **surface** is a **connected, compact, oriented, 2-manifold** considered up to oriented homeomorphism.

Example: S_g := the g -torus = the sphere with g handles attached

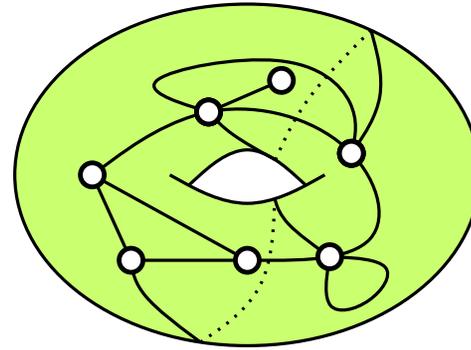


Theorem of classification: every surface is one of the S_g for some $g \geq 0$ called **the genus**.

Rough definition

A **map** is a **graph** drawn on a **surface**.

? ? ?

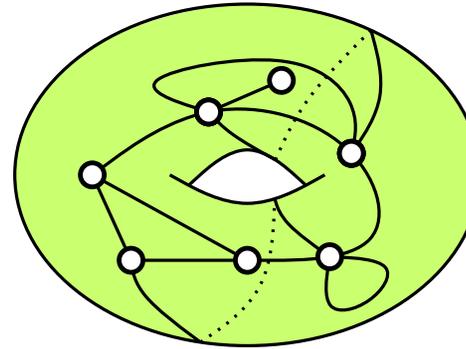


Our **graphs** are unlabelled, connected, and may have **loops** or **multiple edges**.

Rough definition

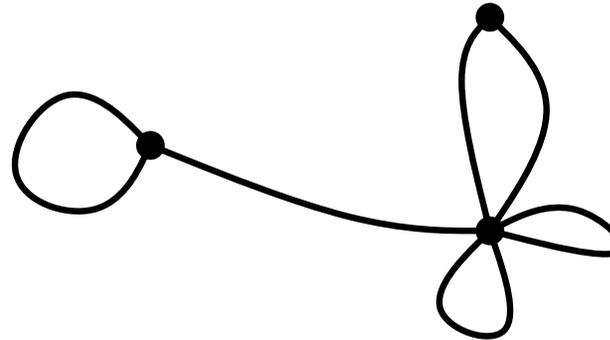
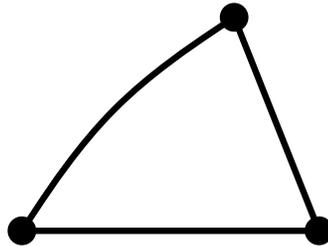
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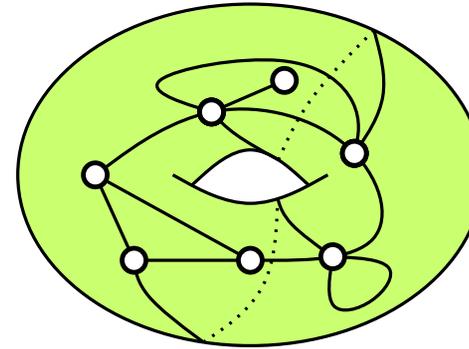
Examples:



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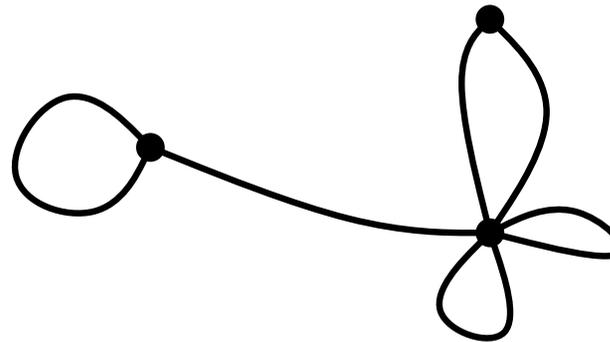
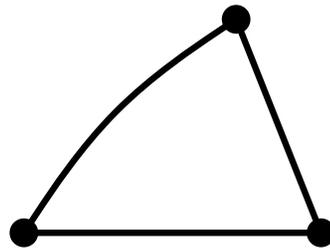
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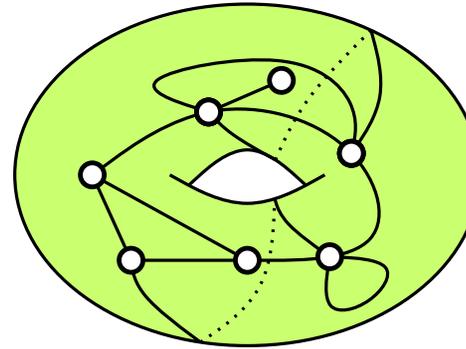


The **degree of a vertex** is the number of **half-edges** incident to it.

Rough definition

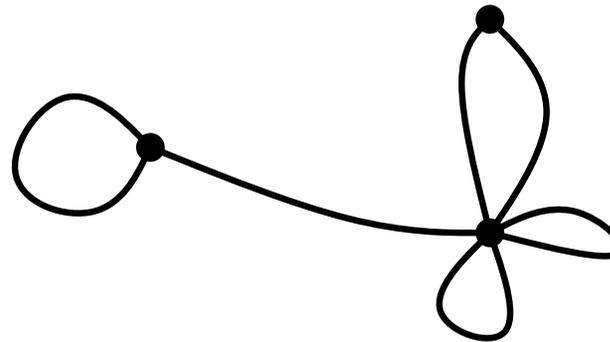
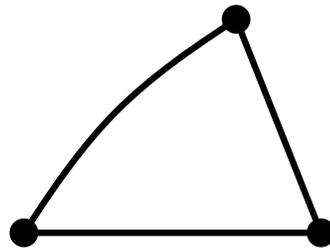
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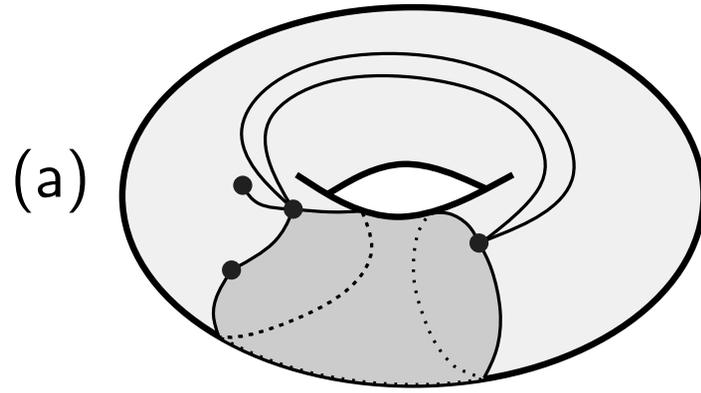


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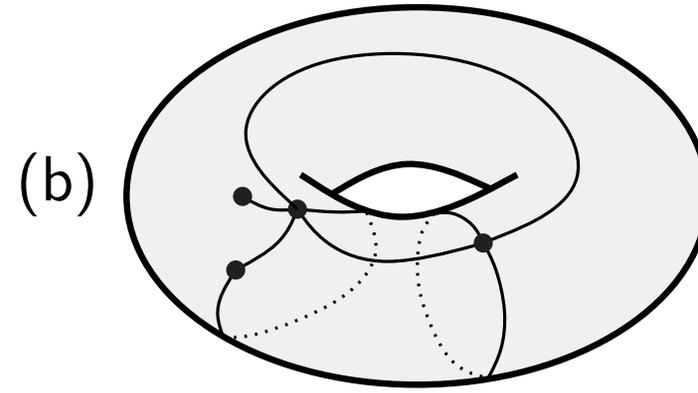
A **proper embedding** of a graph in a surface is a continuous drawing of the graph on the surface without edge-crossings.

Topological definition of a map

A **map** is a proper embedding of a graph G in a surface \mathbb{S} such the connected components of $G \setminus \mathbb{S}$ (called **faces**) are **topological disks**.



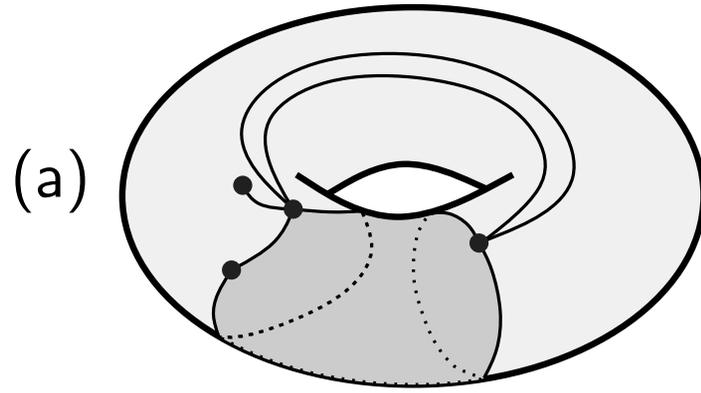
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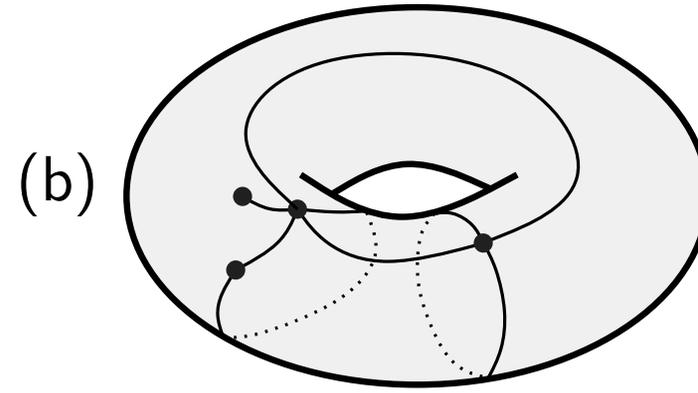
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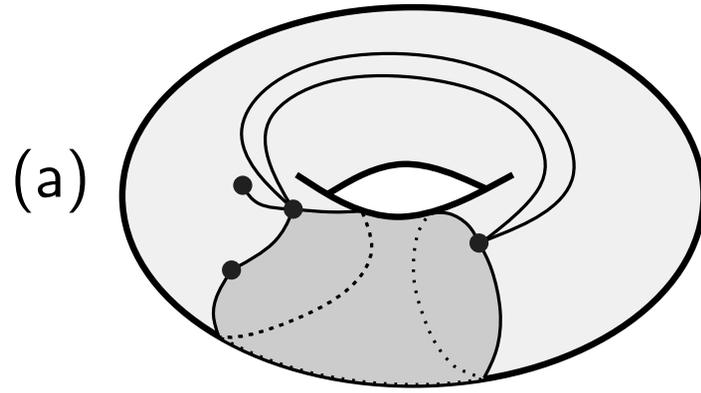


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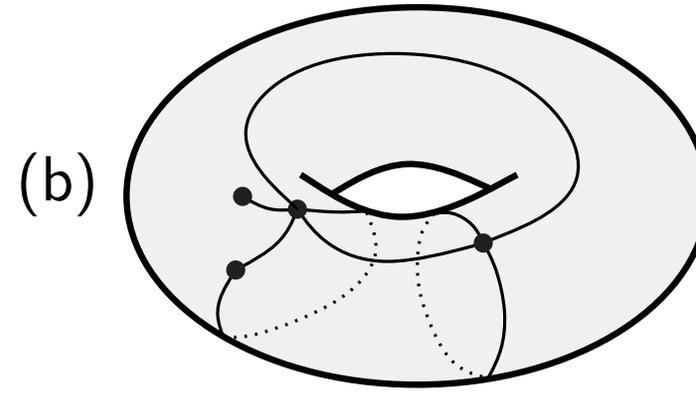
Maps are considered up to oriented homeomorphisms.

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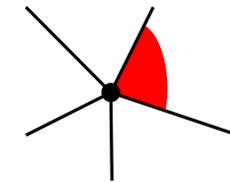
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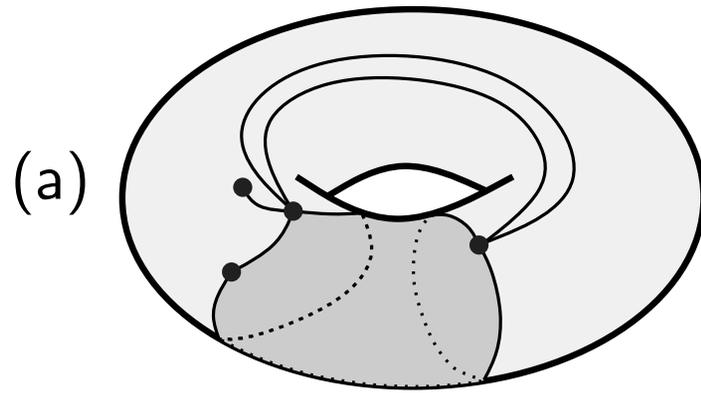
Important notion: a **corner** is an angular sector delimited by two consecutive half-edges in the neighborhood of a vertex.



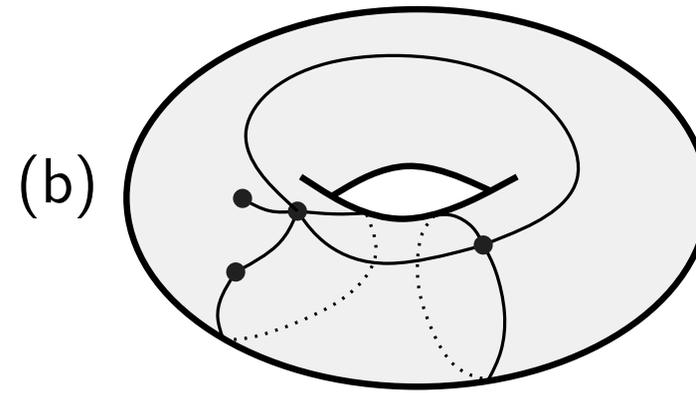
There is a canonical bijection between corners and half-edges.

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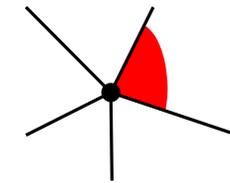
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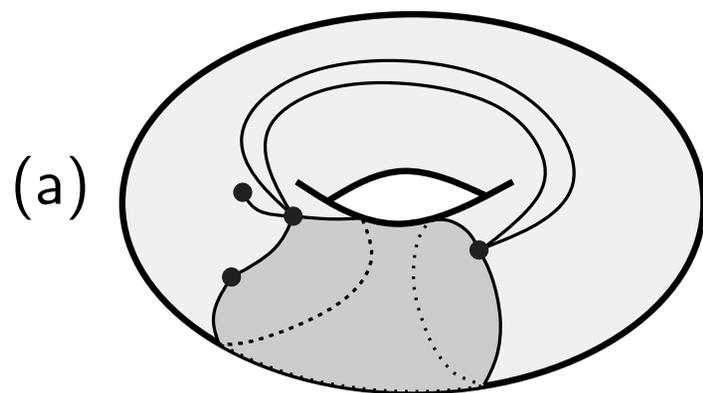
The **degree** of a vertex (or face) is the **number of corners** incident to it.

If a map has n edges then:

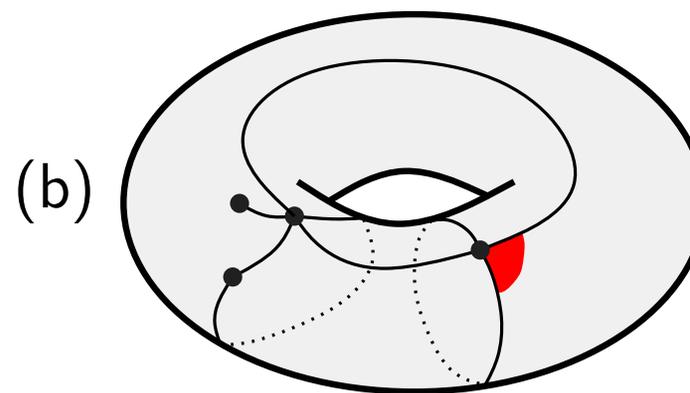
$$2n = \# \text{ corners} = \# \text{ half-edges} = \sum \text{ face degrees} = \sum \text{ vertex degrees}.$$

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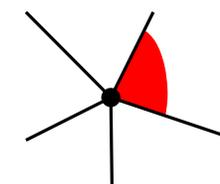
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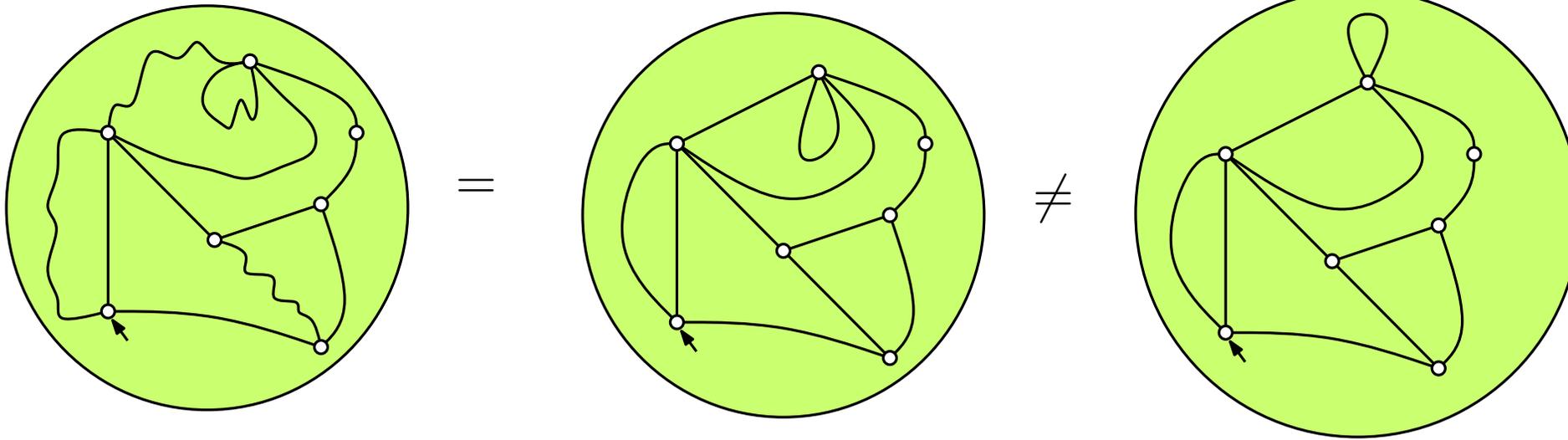
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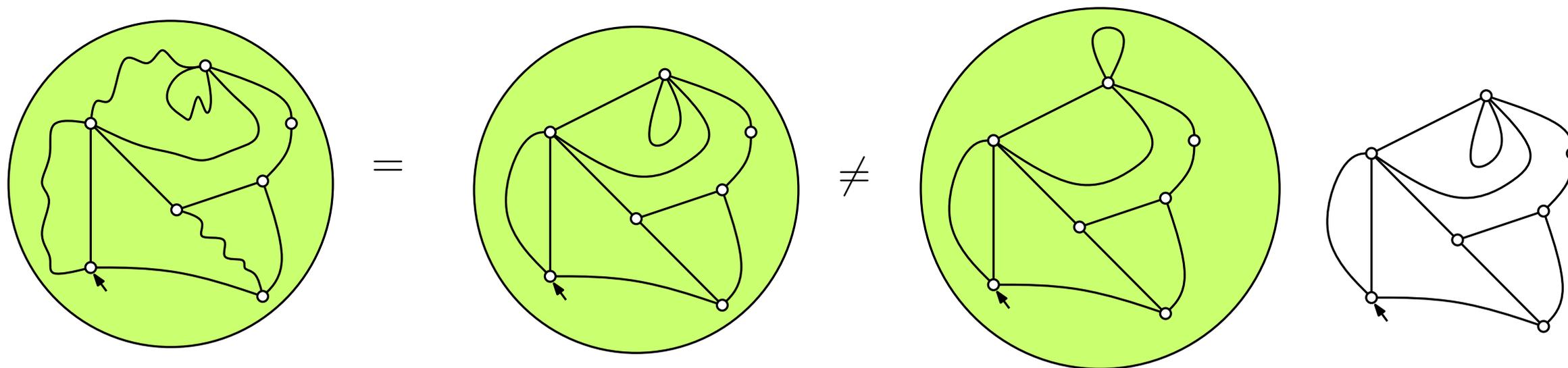
A **rooted map** is a map with a **distinguished corner** (or half-edge)

Example I: planar maps



Convention The infinite face is taken to be the root face.

Example I: planar maps



Maps of **genus 0** are called **planar** rather than spherical...

Convention The **infinite face** is taken to be the **root face**.

Advertisement for the next lecture:

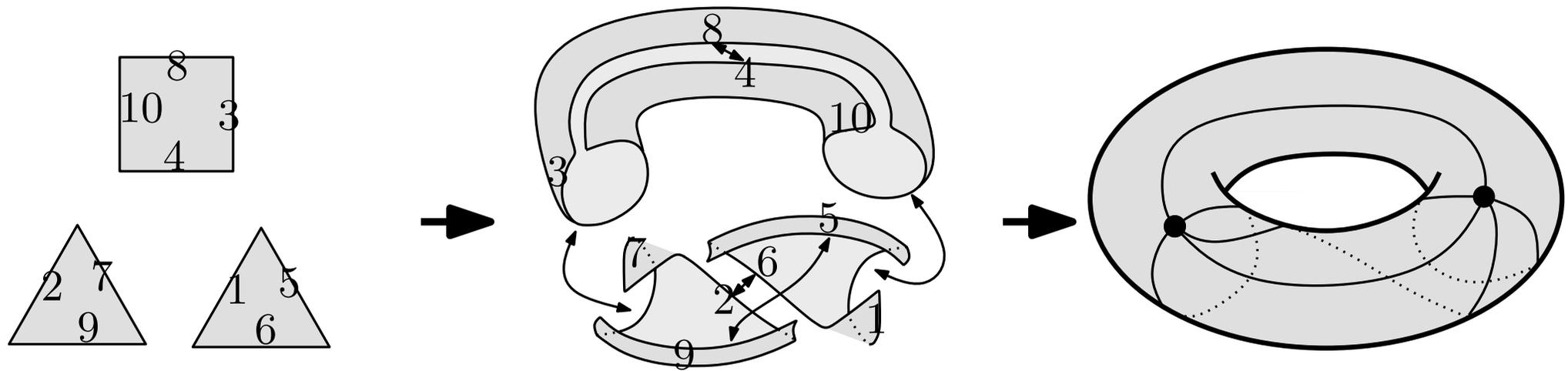
There are $\frac{2 \cdot 3^n}{n+2} \text{Cat}(n)$ rooted planar maps with n edges

(a nice number)

Maps as polygon gluings

An easy way to construct a map:

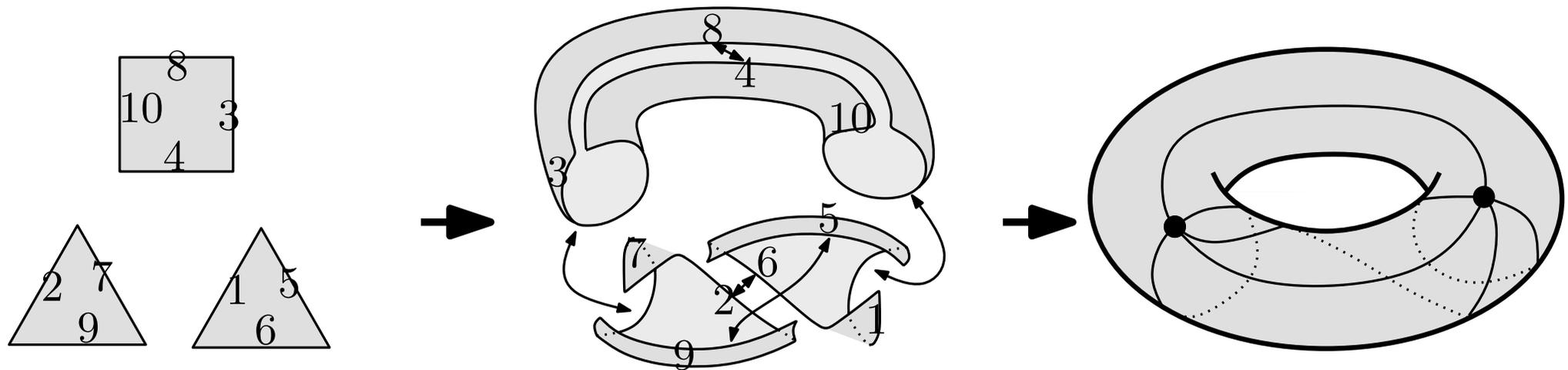
start with a family of polygons with $2n$ sides in total and glue them according to your favorite matching – just be careful to obtain something connected.



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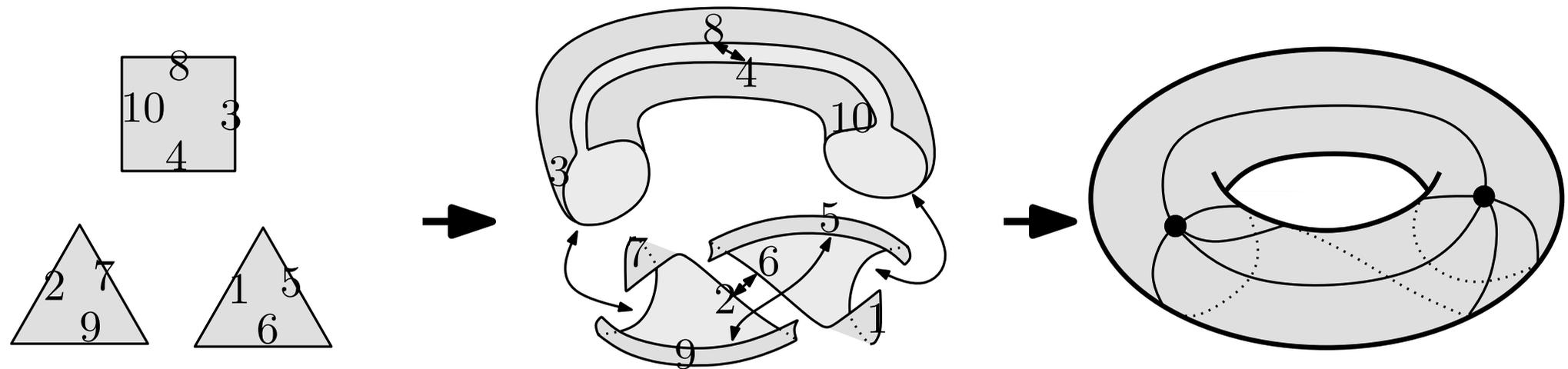


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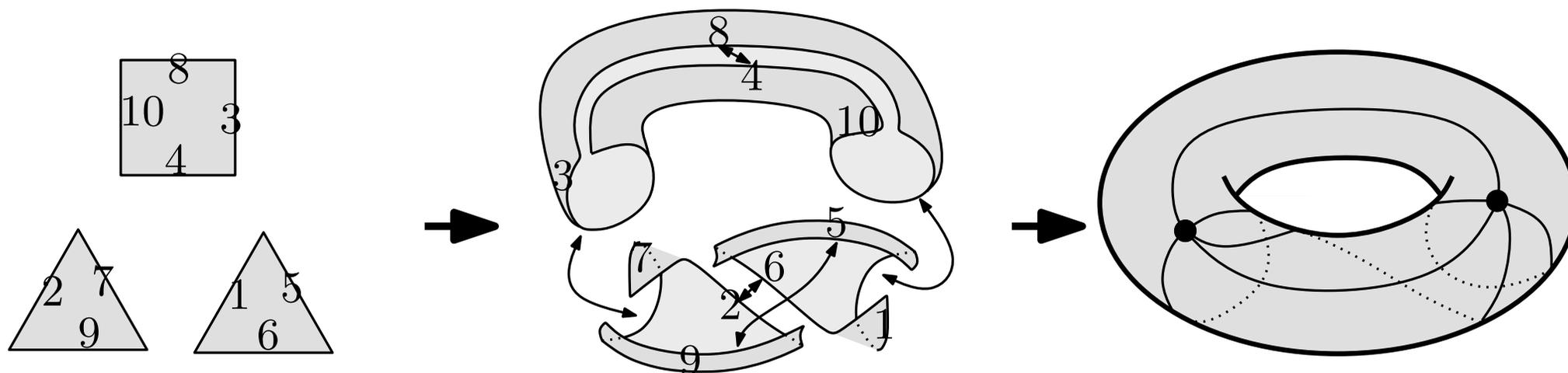
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Claim: provided it is connected, the object we construct is a [map](#).

Proof: We clearly build a surface with a graph on it, and by construction [the faces are our polygons](#) – hence topological disks.

Proposition: any map can be obtained in this way.

Heuristic proof: to go from right to left, just cut the surface along the edges of the graph.

Maps as rotation systems

A **rotation system** on a graph is the data of a **cyclic order** of the half-edges around each vertex.

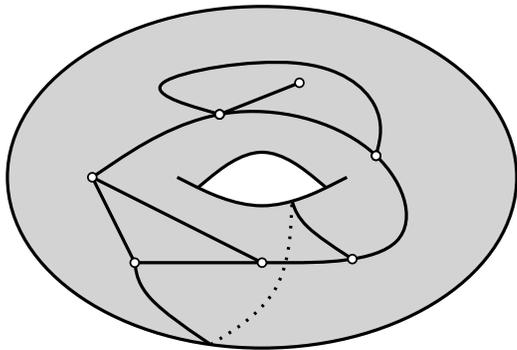
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Fact: There is a **natural mapping**:

Maps \longrightarrow Graphs equipped with a rotation system

...given by the local counterclockwise ordering!



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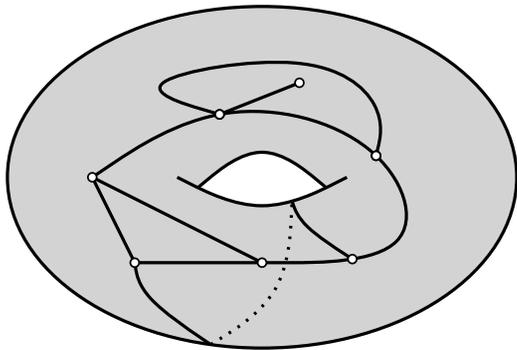
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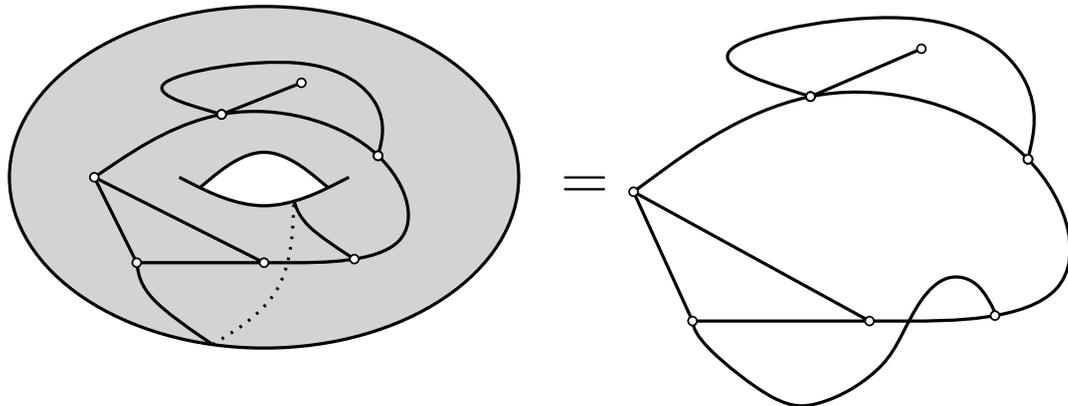
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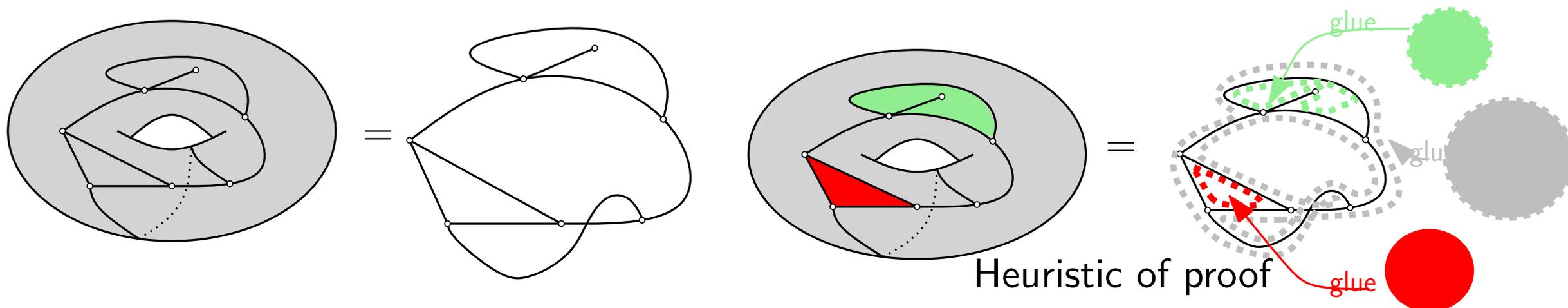
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Reminder: At this stage I should draw more examples on the board

Combinatorial definition with permutations

A **labelled map** of size n is a **triple of permutations** (σ, α, ϕ) in \mathfrak{S}_{2n} such that

- $\alpha\sigma = \phi$
- α has cycle type $(2, 2, \dots, 2)$.
- $\langle \sigma, \alpha, \phi \rangle$ acts **transitively** on $[1..2n]$.

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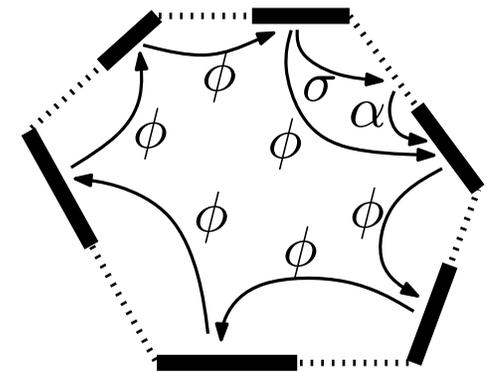
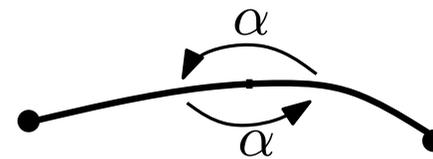
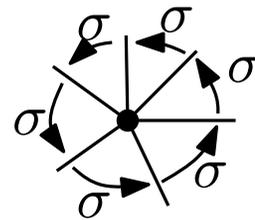
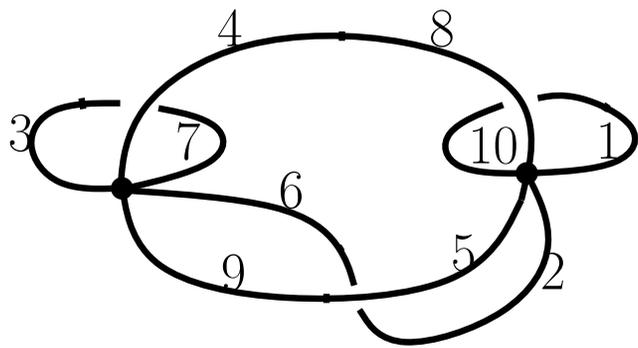
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Thm: There is a bijection between **labelled maps** of size n and **graphs with rotation systems** whose **half-edges** are labelled from 1 to $2n$.



$$\sigma = (1, 8, 10, 5, 2)(3, 9, 6, 7, 4)$$

$$\alpha = (1, 10)(2, 6)(3, 7)(4, 8)(5, 9)$$

$$\phi = \alpha\sigma = (1, 4, 7, 8)(2, 10, 9)(3, 5, 6)$$

Note:

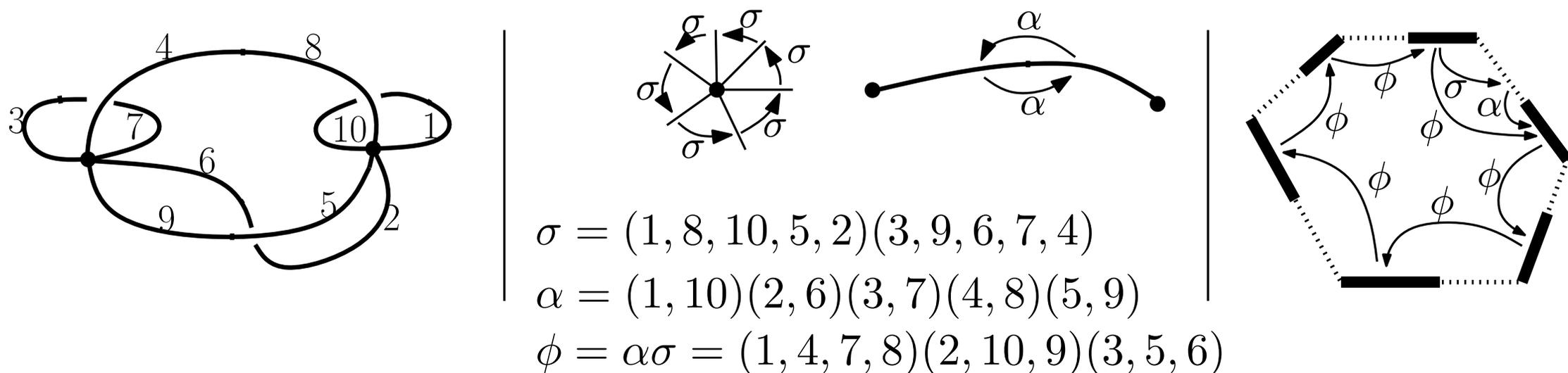
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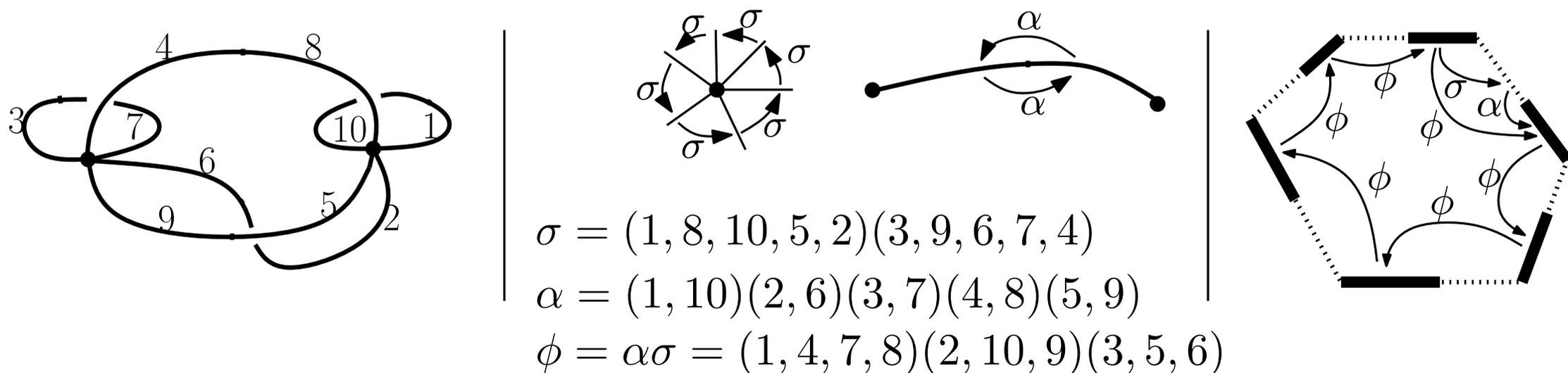
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A **rooted map** is an **equivalence class** of labelled maps under renumbering of $[2..2n]$.

labelled map " = " $(2n - 1)! \times$ **rooted map**

Duality

A **labelled map** of size n is a **triple of permutations** (σ, α, ϕ) in \mathfrak{S}_{2n} such that

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The mapping $(\sigma, \alpha, \phi) \rightarrow (\phi, \alpha, \sigma)$ is an **involution** on maps called duality. It exchanges **vertices** and **faces**.

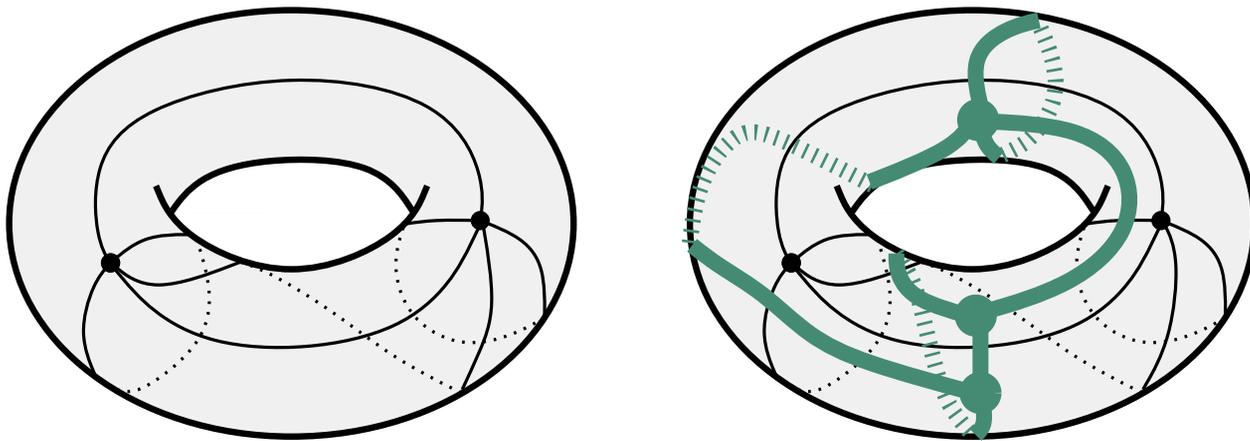
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There is also a **well-known** graphical version:



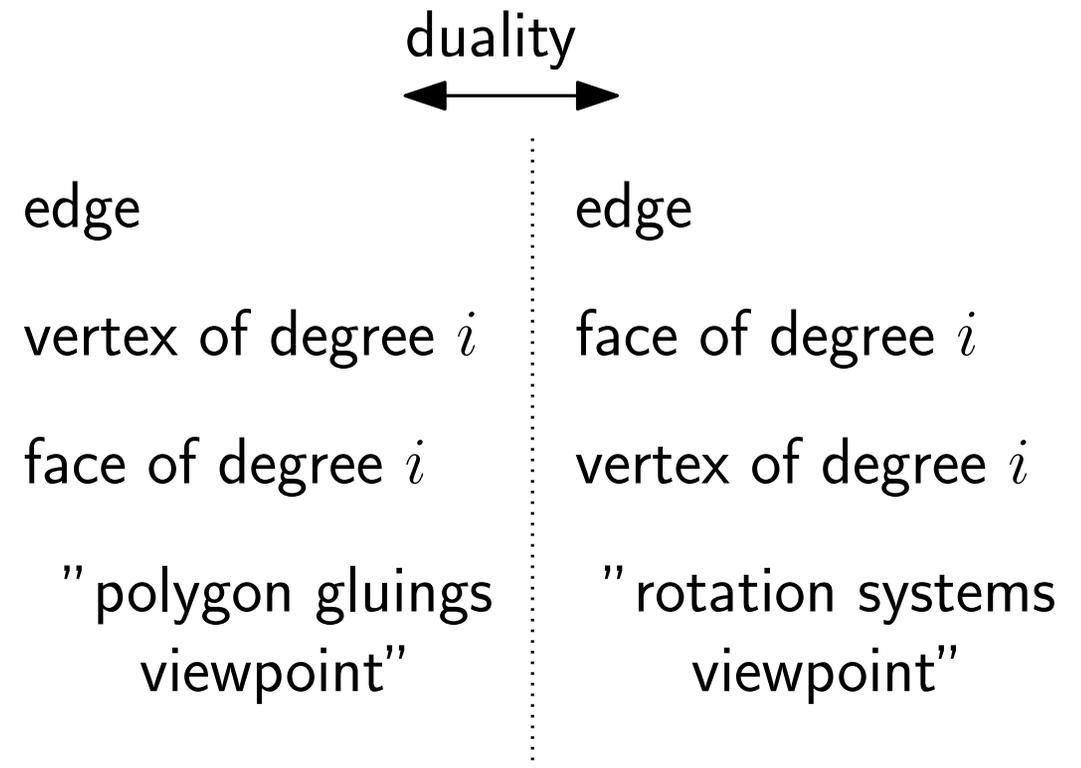
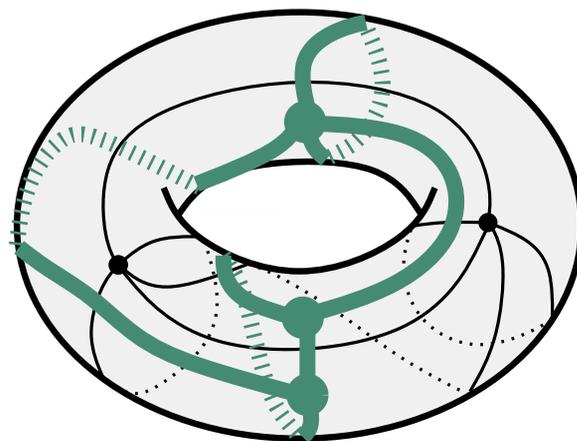
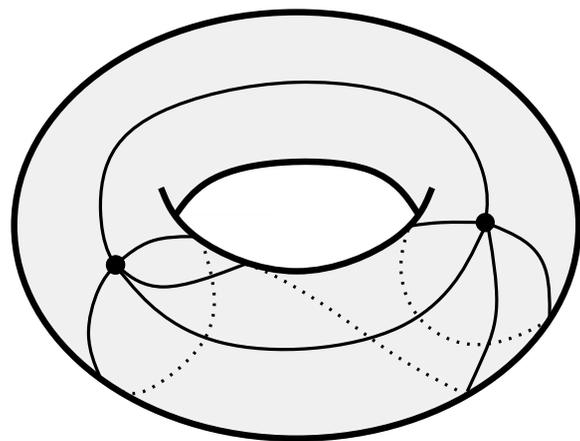
Duality

A **labelled map** of size n is a **triple of permutations** (σ, α, ϕ) in \mathfrak{S}_{2n} such that

- $\alpha\sigma = \phi$
- α has cycle type $(2, 2, \dots, 2)$.
- $\langle \sigma, \alpha, \phi \rangle$ acts **transitively** on $[1..2n]$.

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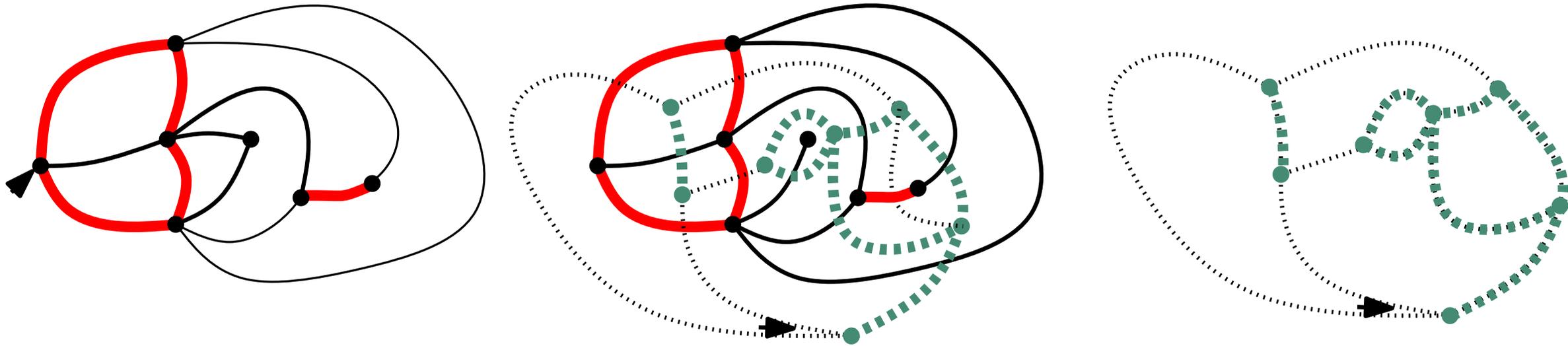
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Duality II – dual submap

If \mathfrak{m} is a map with underlying graph G then any subgraph $H \subset G$ induces a **submap** of G , with same vertex set, by **restricting the cyclic ordering to H** .

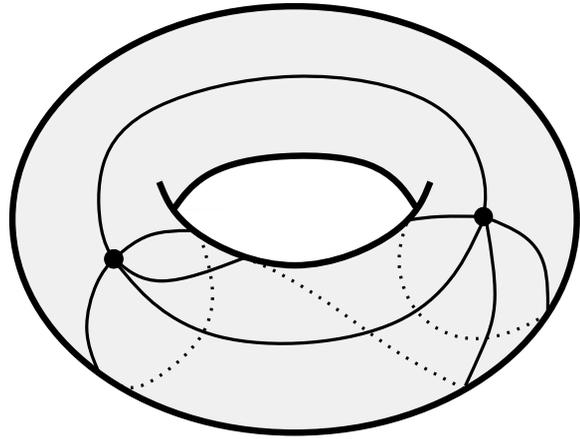
Note that the submap is **not necessarily connected** (and can have a different genus)



The **dual submap** is the submap of \mathfrak{m}^* formed by edges whose dual is not in H

Proposition: The total number of **faces** of a submap and its dual submap are equal.

Euler's formula



For a map of genus g with n edges, f faces, v vertices, we have:

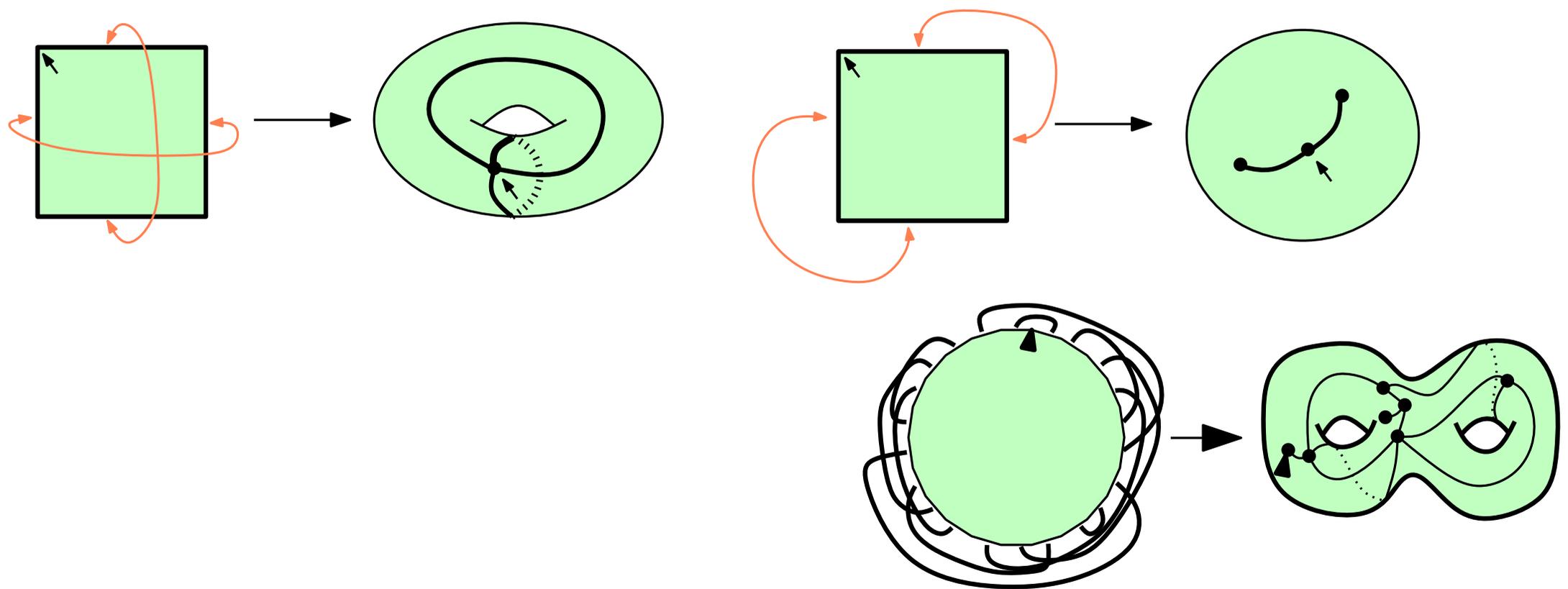
$$v + f = n + 2 - 2g$$

In particular we can recover the genus from the combinatorics
(we don't need to "see" the surface...)

Example II: one-face maps

What is a one-face map? Clear in the “polygon gluing viewpoint”.

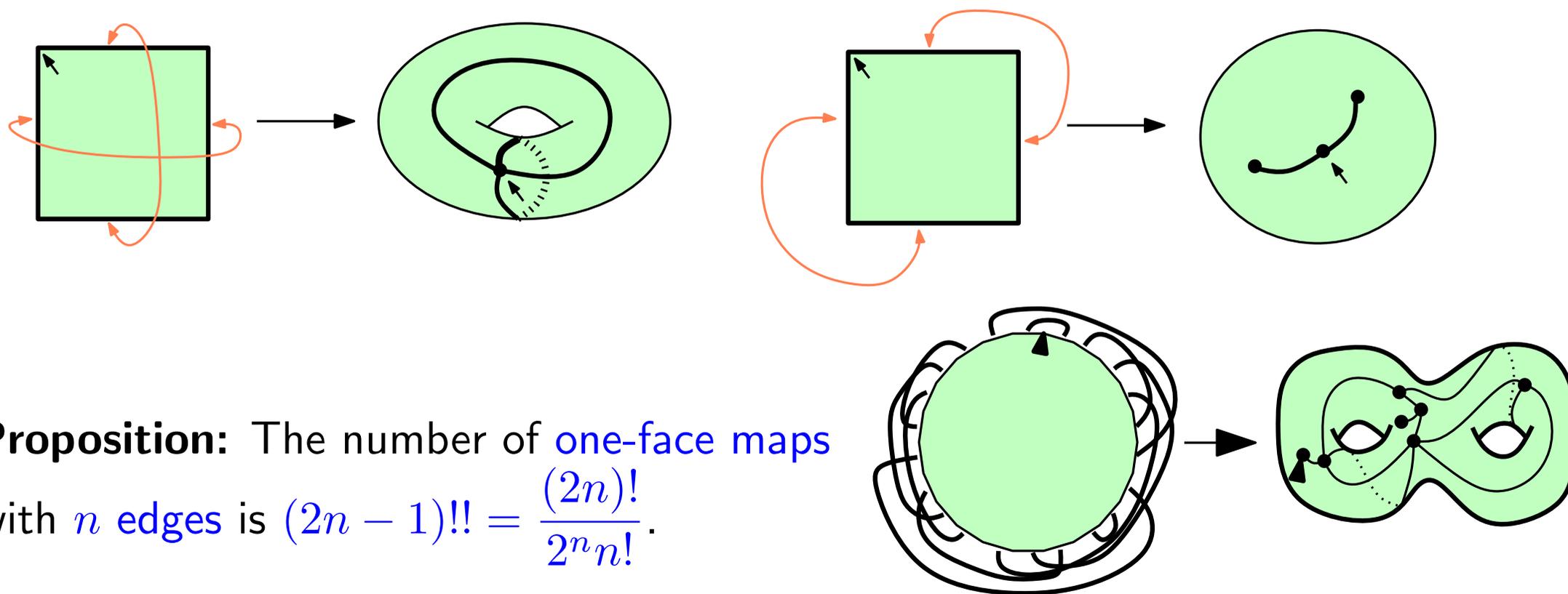
Start with a $2n$ -gon and glue the edges together according to some matching.



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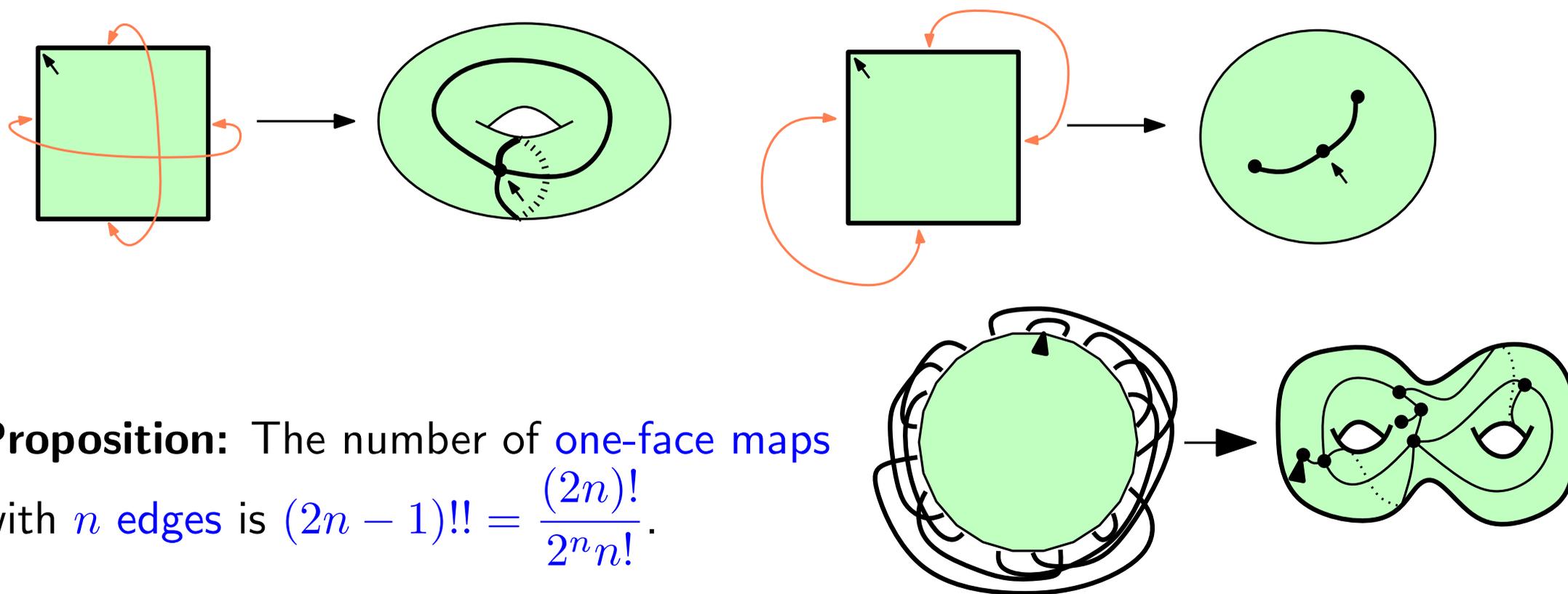


Proposition: The number of one-face maps with n edges is $(2n - 1)!! = \frac{(2n)!}{2^n n!}$.

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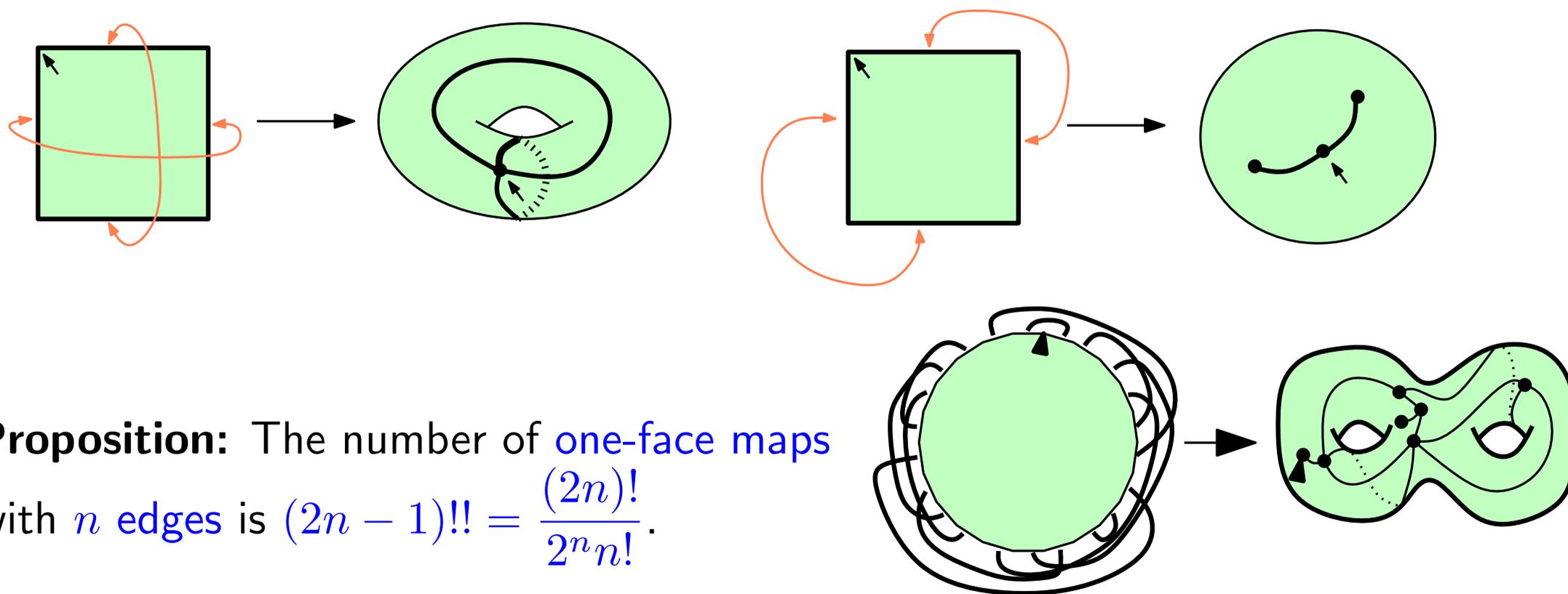
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Much harder: control the genus! (see the exercises)

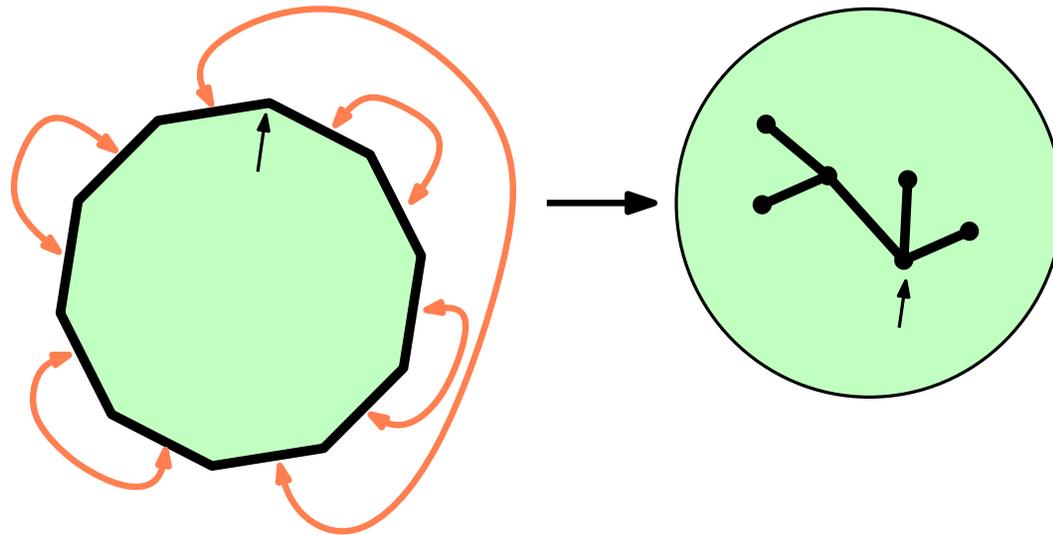
Example III: planar one-face maps (=plane trees, Catalan trees)

Euler formula: $v + f = n + 2 - 2g$

$f = 1, g = 0$ gives $v = n + 1$

this is a tree!

Tree+root corner+rotation system = plane tree (a.k.a. ordered tree)



Proposition: The number of rooted plane trees with n edges is $\text{Cat}(n) = \frac{1}{n+1} \binom{2n}{n}$.

Tomorrow: we start counting!

everything will be planar (no strange surface yet so don't be afraid)

if you don't know what to do tonight, try exercise 0 from the webpage.