Reachability for multidimensional continued fractions and minimality of interval exchange transformations

Vincent Delecroix

CNRS / Université de Bordeaux
CODYS meeting Arles-Bordeaux-Caen-Paris, december 2020

## Reachability / orbit problem

Let $\mathcal{I}$ be a set of triples $(f, X, Y)$ where $f: D \rightarrow D$ is a function $X \subset D$ and $Y \subset D$.
input: a triple $(f, X, Y)$ in $\mathcal{I}$
$\operatorname{REACH}(\mathcal{I})$ : output: whether there exists $x \in X, y \in Y, n \geq 0$ such that $f^{n}(x)=y$

## Reachability / orbit problem

Let $\mathcal{I}$ be a set of triples $(f, X, Y)$ where $f: D \rightarrow D$ is a function $X \subset D$ and $Y \subset D$.
input: a triple $(f, X, Y)$ in $\mathcal{I}$
$\operatorname{REACH}(\mathcal{I})$ : output: whether there exists $x \in X, y \in Y, n \geq 0$ such that $f^{n}(x)=y$

Our functions $f: D \rightarrow D$ will be piecewise affine maps on $K^{d}$ where $K \subset \mathbb{R}$ and $X, Y$ will be polyhedron.

## Example 1: piecewise translations

$d \geq 1$
$\Lambda \subset \mathbb{R}$ stable under addition
$f:[0,1]^{d} \rightarrow[0,1]^{d}$ bijective piecewise translations with discontinuities contained in hyperplanes $\left\{x_{i}=\alpha\right\}$ with $i \in\{1, \ldots, d\}$ and $\alpha \in \Lambda$ and translations contained in $\Lambda^{d}$.

Example 1: piecewise translations $d=1$
interval exchange transformations with permutation $\pi=\left(\begin{array}{ccc}A & B & C \\ D \\ D & C & B\end{array}\right)$


## Example 1: piecewise translations $d=1$

interval exchange transformations with permutation $\pi=\left(\begin{array}{ccc}A & B & C \\ D \\ D & C & B\end{array}\right)$


Theorem (Keane condition)
$X=\left\{\lambda_{D}, \lambda_{D}+\lambda_{C}, \lambda_{D}+\lambda_{C}+\lambda_{B}\right\}$
$Y=\left\{\lambda_{A}, \lambda_{A}+\lambda_{B}, \lambda_{A}+\lambda_{B}+\lambda_{C}\right\}$
If there is no $n \geq 0, x \in X$ and $y \in Y$ such that $T_{\pi, \lambda}^{n}(x)=y$ then all infinite orbits of $T_{\pi, \lambda}$ are dense.

## Example 1: piecewise translations $d=1$

interval exchange transformations with permutation $\pi=\left(\begin{array}{ccc}A & B & C \\ D \\ D & C & B\end{array}\right)$


Theorem (Keane condition)
$X=\left\{\lambda_{D}, \lambda_{D}+\lambda_{C}, \lambda_{D}+\lambda_{C}+\lambda_{B}\right\}$
$Y=\left\{\lambda_{A}, \lambda_{A}+\lambda_{B}, \lambda_{A}+\lambda_{B}+\lambda_{C}\right\}$
If there is no $n \geq 0, x \in X$ and $y \in Y$ such that $T_{\pi, \lambda}^{n}(x)=y$ then all infinite orbits of $T_{\pi, \lambda}$ are dense.

Remark:

- $\lambda \in \mathbb{Q}^{4} \Rightarrow T_{\pi, \lambda}$ completely periodic
- $\lambda$ totally irrational $\Rightarrow T_{\pi, \lambda}$ satisfies Keane condition


## Example 1: piecewise translations $d=2$

P. Hooper rectangle exchange system $(d=2)$


## Example 2: simplicial systems

Piecewise affine functions with Markovian topological dynamics.

## Why do we care about reachability?

(1) Dynamical systems: periodic orbits, minimality of interval exchange transformations, $\mathrm{GL}(2, \mathbb{R})$-orbit closures of translation surfaces,
(2) Theoretical computer science
(3) Control theory

## Decidability

For any $d \geq 1, \operatorname{REACH}\left(\operatorname{Mat}_{d \times d}(K)\right)$ is decidable (Kannan-Lipton 1980 for $K=\mathbb{Q}$ ).

## Undecidability result

We define a restricted class $\mathcal{F}$ of 2-dimensional rational piecewise affine functions on $D=\left([0,1) \cap \mathbb{Z}\left[\frac{1}{2}\right]\right)^{2}$ with discontinuities contained in hyperplanes $x=p / 2^{n}$ or $y=p / 2^{n}$
Piecewise $(x, y) \mapsto(a x+b y+c, d x+e y+f)$ with $a, b, d, e \in\left\{1,2, \frac{1}{2}\right\}$ and $c, f \in \mathbb{Z}\left[\frac{1}{2}\right]$.

## Undecidability result

We define a restricted class $\mathcal{F}$ of 2-dimensional rational piecewise affine functions on $D=\left([0,1) \cap \mathbb{Z}\left[\frac{1}{2}\right]\right)^{2}$ with discontinuities contained in hyperplanes $x=p / 2^{n}$ or $y=p / 2^{n}$
Piecewise $(x, y) \mapsto(a x+b y+c, d x+e y+f)$ with $a, b, d, e \in\left\{1,2, \frac{1}{2}\right\}$ and $c, f \in \mathbb{Z}\left[\frac{1}{2}\right]$.

Theorem (folklore?)
The problem $\left.\operatorname{REACH}\left(\{(\mathcal{F},(1,1),(0,0))\}_{f \in \mathcal{F}}\right)\right)$ is undecidable.

## Undecidability result

We define a restricted class $\mathcal{F}$ of 2-dimensional rational piecewise affine functions on $D=\left([0,1) \cap \mathbb{Z}\left[\frac{1}{2}\right]\right)^{2}$ with discontinuities contained in hyperplanes $x=p / 2^{n}$ or $y=p / 2^{n}$
Piecewise $(x, y) \mapsto(a x+b y+c, d x+e y+f)$ with $a, b, d, e \in\left\{1,2, \frac{1}{2}\right\}$ and $c, f \in \mathbb{Z}\left[\frac{1}{2}\right]$.

Theorem (folklore?)
The problem $\left.\operatorname{REACH}\left(\{(\mathcal{F},(1,1),(0,0))\}_{f \in \mathcal{F}}\right)\right)$ is undecidable.

Theorem (folklore bis)
There exists an explicit function $f \in \mathcal{F}$ such that the problem $\operatorname{REACH}\left(\{(f, x,(0,0))\}_{x \in D}\right)$ is undecidable.

## Open problems

Whether the reachability problem is decidable or undecidable is an open question for the following classes of maps

- interval exchange transformations over number fields,
- 1-dimensional piecewise rational affine maps,
- d-dimensional piecewise translations over number fields,
- simplicial systems,


## This talk

(1) Partial result for reachability of interval exchange transformations.
(2) Link with a specific reachability problem for simplicial systems.

## Periodic orbits and relations

( $\pi, \lambda$ ): interval exchange transformation.
A periodic orbit gives rise to a non-negative relation on translations


$$
\tau_{A}+2 \tau_{D}+\tau_{C}=0
$$

## Periodic orbits and relations

( $\pi, \lambda$ ): interval exchange transformation.
A periodic orbit gives rise to a non-negative relation on translations


$$
\tau_{A}+2 \tau_{D}+\tau_{C}=0
$$

$$
\mathrm{R}(\tau):=\left\{r \in \mathbb{Z}^{A}: \sum_{\alpha \in A} r_{\alpha} \cdot \tau_{\alpha}=0\right\}
$$

## Periodic orbits and relations

$(\pi, \lambda)$ : interval exchange transformation.
A periodic orbit gives rise to a non-negative relation on translations


$$
\tau_{A}+2 \tau_{D}+\tau_{C}=0
$$

$$
\mathrm{R}(\tau):=\left\{r \in \mathbb{Z}^{A}: \sum_{\alpha \in A} r_{\alpha} \cdot \tau_{\alpha}=0\right\}
$$

## Lemma

Let $\pi$ be a permutation and $\lambda \in \mathbb{R}^{d}$. If $\mathbb{Z}_{\geq 0}^{d} \cap \mathbb{R}(\tau)=\{0\}$ then the iet $T_{\pi, \lambda}$ has no periodic orbit.

## Rauzy induction

Idea: dynamics on the space of interval exchanges $\mathcal{R}:(\pi, \lambda) \mapsto\left(\pi^{\prime}, \lambda^{\prime}\right)$
integral non-negative matrix cocycle $A_{n}(\pi, \lambda): \lambda=A_{n}(\pi, \lambda) \cdot \lambda^{(n)}$.
We have $\tau^{(n)}=A_{n}(\pi, \lambda)^{t} \tau^{(0)}$.

## Improved lemma with induction

Replace $C_{0}(\pi, \lambda):=\mathbb{Z}_{\geq 0}^{d}$ by the subcone $C_{n}(\pi, \lambda):=A_{n}(\pi, \lambda) \mathbb{Z}^{d}$.

## Improved lemma with induction

Replace $C_{0}(\pi, \lambda):=\mathbb{Z}_{\geq 0}^{d}$ by the subcone $C_{n}(\pi, \lambda):=A_{n}(\pi, \lambda) \mathbb{Z}^{d}$.

## Lemma

Let $\pi$ be a permutation and $\lambda \in \mathbb{R}^{d}$. If the Rauzy induction is well defined up to step $n$ and $C_{n}(\pi, \lambda) \cap R(\tau)=\{0\}$ then the iet $T_{\pi, \lambda}$ has no periodic orbit.

## Improved lemma with induction

Replace $C_{0}(\pi, \lambda):=\mathbb{Z}_{\geq 0}^{d}$ by the subcone $C_{n}(\pi, \lambda):=A_{n}(\pi, \lambda) \mathbb{Z}^{d}$.

## Lemma

Let $\pi$ be a permutation and $\lambda \in \mathbb{R}^{d}$. If the Rauzy induction is well defined up to step $n$ and $C_{n}(\pi, \lambda) \cap \mathrm{R}(\tau)=\{0\}$ then the iet $T_{\pi, \lambda}$ has no periodic orbit.

Questions:
(1) under which condition Rauzy induction is well defined?
(2) what does look like $C_{n}(\pi, \lambda)$ ?

## Infinite orbits of the Rauzy induction

## Theorem (Rauzy, Veech)

Let $\pi$ be an irreducible permutation and $\lambda \in \mathbb{R}^{A}$. Then the following conditions are equivalent

- $(\pi, \lambda)$ satisfies the Keane condition,
- Rauzy induction is defined for all times and $\lambda^{(n)} \rightarrow 0$,
- for all $n \geq 0$, there exists $m$ such that $A_{m}\left(\pi^{(n)}, \lambda^{(n)}\right)>0$.

Short digression on Multidimensional continued fractions

## The shape of the cones $C_{n}(\pi, \lambda)$

## Theorem (Rauzy, Veech)

Let $\pi$ be an irreducible permutation and $\lambda \in \mathbb{R}^{A}$. Then

- if $(\pi, \lambda)$ does not satisfy Keane condition, then there exists $n$ such that the rightmost intervals of $\left(\pi^{(n)}, \lambda^{(n)}\right)$ have equal lengths ("trivial saddle connection"),
- if $(\pi, \lambda)$ satisfies the Keane condition, then $C_{n}(\pi, \lambda)$ converges to the cone of invariant measures of $T_{\pi, \lambda}$.


## The semi-algorithm

input: $(\pi, \lambda)$

- $\mathrm{n}=0$
- repeat
- if $C_{n}(\pi, \lambda) \cap R(\pi, \lambda)=\{0\}$ : return "no periodic orbit"
- if Keane condition is violated at the $n$-th step: return "found saddle connection"
- $\mathrm{n}=\mathrm{n}+1$


## bad news: $S A F=0$

Starting from 6 letters, there exists some $(\pi, \lambda)$ with

- Keane condition
- $C_{\infty}(\pi, \lambda) \cap \mathrm{R}(\tau) \neq\{0\}$.


## bad news: $S A F=0$

Starting from 6 letters, there exists some $(\pi, \lambda)$ with

- Keane condition
- $C_{\infty}(\pi, \lambda) \cap \mathrm{R}(\tau) \neq\{0\}$.

The most famous one is the so-called Arnoux-Rauzy example

$$
\pi=\left(\begin{array}{ccccc}
A_{1, \ell} & A_{1, r} & A_{2} & B_{1} & B_{2}
\end{array} C_{1} C_{2}\right) \quad \lambda=\left(\theta+1, \theta^{2}-\theta-1, \theta^{2}, \theta, \theta, 1,1\right)
$$

where $\theta^{3}-\theta^{2}-\theta-1=0$.

## Lemma

The semi-algorithm does not terminate on minimal $S A F=0$ instances.

## Hope (out of reach)

## Conjecture

Let $\pi$ be an irreducible permutation on $d$ letters and $\lambda \in\left(\mathbb{R}_{\geq 0} \cap \overline{\mathbb{Q}}\right)^{d}$. Then if $T_{\pi, \lambda}$ is minimal, it is uniquely ergodic.

If the conjecture was true, then the semi-algorithm would always terminate on algebraic $\lambda$ with $\operatorname{SAF}(\pi, \lambda) \neq 0$.

## It is an algorithm on quadratic fields

## Theorem (Boshernitzan 88)

For $\lambda$ in a quadratic number field $\mathbb{Q}[\sqrt{D}]^{d}$ the semi-algorithm always terminate

- $S A F=0$ implies completely periodic,
- minimal implies uniquely ergodic.

