

On the second Lyapunov exponent of some multidimensional continued fraction algorithms

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Multidimensional continued fraction algorithms

Write $\mathbf{x} \in \Delta \subseteq [0, 1]^d$ as $\mathbf{x} = \lim_{n \rightarrow \infty} \frac{\mathbf{p}^{(n)}}{q^{(n)}}$ with small $\|q^{(n)}\mathbf{x} - \mathbf{p}^{(n)}\|$.

We consider MCF algorithms

$$A : \Delta \rightarrow \mathrm{GL}(d+1, \mathbb{Z})$$

with associated transformations

$$T : \Delta \rightarrow \Delta, \quad \mathbf{x} \mapsto \pi(\iota(\mathbf{x}) A(\mathbf{x})^{-1})$$

$$\iota(x_1, \dots, x_d) = (1, x_1, \dots, x_d), \quad \pi(x_0, x_1, \dots, x_d) = \left(\frac{x_1}{x_0}, \dots, \frac{x_d}{x_0}\right)$$

Regular continued fractions, $d = 1$: $A(x) = \begin{pmatrix} \lfloor \frac{1}{x} \rfloor & 1 \\ 1 & 0 \end{pmatrix}$,

$$T(x) = \pi\left((1, x) \begin{pmatrix} 0 & 1 \\ 1 & -\lfloor \frac{1}{x} \rfloor \end{pmatrix}\right) = \pi\left(x, 1 - \lfloor \frac{1}{x} \rfloor x\right) = \frac{1}{x} - \left\lfloor \frac{1}{x} \right\rfloor$$

Jacobi–Perron algorithm (1868, 1907)

$$A_{\text{JP}}(x_1, \dots, x_d) = \begin{pmatrix} \lfloor \frac{1}{x_1} \rfloor & 1 & \lfloor \frac{x_2}{x_1} \rfloor & \dots & \lfloor \frac{x_d}{x_1} \rfloor \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \vdots & & \ddots & 1 \\ 1 & 0 & \dots & \dots & 0 \end{pmatrix}$$

$$A_{\text{JP}}(x_1, \dots, x_d)^{-1} = \begin{pmatrix} 0 & 0 & \dots & 0 & 1 \\ 1 & -\lfloor \frac{x_2}{x_1} \rfloor & \dots & -\lfloor \frac{x_d}{x_1} \rfloor & -\lfloor \frac{1}{x_1} \rfloor \\ 0 & 1 & 0 & \dots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & 1 & 0 \end{pmatrix}$$

$$T_{\text{JP}} : [0, 1]^d \rightarrow [0, 1]^d,$$

$$(x_1, \dots, x_d) \mapsto \left(\frac{x_2}{x_1} - \left\lfloor \frac{x_2}{x_1} \right\rfloor, \dots, \frac{x_d}{x_1} - \left\lfloor \frac{x_d}{x_1} \right\rfloor, \frac{1}{x_1} - \left\lfloor \frac{1}{x_1} \right\rfloor \right)$$

Convergents, linear cocycle

The convergents $\frac{\mathbf{p}_i^{(n)}}{q_i^{(n)}} = \frac{(p_{i,1}^{(n)}, \dots, p_{i,d}^{(n)})}{q_i^{(n)}}$ of $\mathbf{x}$ are given by the cocycle

$$A^{(n)}(\mathbf{x}) = A(T^{n-1}\mathbf{x}) \cdots A(T\mathbf{x}) A(\mathbf{x}) = \begin{pmatrix} q_0^{(n)} & p_{0,1}^{(n)} & \cdots & p_{0,d}^{(n)} \\ q_1^{(n)} & p_{1,1}^{(n)} & \cdots & p_{1,d}^{(n)} \\ \vdots & \vdots & \ddots & \vdots \\ q_d^{(n)} & p_{d,1}^{(n)} & \cdots & p_{d,d}^{(n)} \end{pmatrix}$$

Weak convergence: $\lim_{n \rightarrow \infty} \frac{\mathbf{p}_i^{(n)}}{q_i^{(n)}} = \mathbf{x}$ for all $0 \leq i \leq d$

Strong convergence: $\lim_{n \rightarrow \infty} \|\mathbf{p}_i^{(n)} - q_i^{(n)} \mathbf{x}\| = 0$ for all $0 \leq i \leq d$

We assume that (Δ, T, μ) is ergodic for an invariant probability measure μ , the cocycle A is log-integrable, with Lyapunov exponents

$$\lambda_1(A) > \lambda_2(A) \geq \lambda_3(A) \geq \cdots \geq \lambda_{d+1}(A)$$

$\lambda_2(A) < 0 \Rightarrow$ almost everywhere exponential convergence

Approximation cocycle

$$\begin{aligned} D^{(n)}(\mathbf{x}) &= \Pi A^{(n)}(\mathbf{x}) H(\mathbf{x}) \\ &= \begin{pmatrix} p_{1,1}^{(n)} - q_1^{(n)} x_1 & \cdots & p_{1,d}^{(n)} - q_1^{(n)} x_d \\ \vdots & \ddots & \vdots \\ p_{d,1}^{(n)} - q_d^{(n)} x_1 & \cdots & p_{d,d}^{(n)} - q_d^{(n)} x_d \end{pmatrix} \in \mathbb{R}^{d \times d}, \end{aligned}$$

$$\Pi = \begin{pmatrix} 0 & 1 & 0 & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 \end{pmatrix}, \quad H(x_1, \dots, x_d) = \begin{pmatrix} -x_1 & \cdots & -x_d \\ 1 & \cdots & 0 \\ 0 & \ddots & 0 \\ 0 & \cdots & 1 \end{pmatrix}.$$

Schweiger'00, Hardcastle–Khanin'02, ...: D is a cocycle of T

$$D^{(n)}(\mathbf{x}) = D(T^{n-1}\mathbf{x}) \cdots D(T\mathbf{x}) D(\mathbf{x}), \quad \text{with } D(\mathbf{x}) = \Pi A(\mathbf{x}) H(\mathbf{x})$$

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{1}{n} \log \|A^{(n)}(\mathbf{x}) \mathbf{y}\| &\leq \lambda_2(A) \text{ for all } \mathbf{y} \in H(\mathbf{x})\mathbb{R}^d = \iota(\mathbf{x})^\perp \text{ (a.e. } \mathbf{x} \in \Delta) \\ &\Rightarrow \lambda_2(A) = \lambda_1(D) \end{aligned}$$

Determining $\lambda_2(A) = \lambda_1(D)$

$$|\det A(\mathbf{x})| = 1 \Rightarrow \sum_{k=1}^{d+1} \lambda_k(A) = 0$$

$$d = 2: \lambda_2(A) = -\lambda_1(A) - \lambda_3(A) = \lambda_1({}^tA^{-1}) - \lambda_1(A)$$

can be used to show $\lambda_2(A) < 0$ for Jacobi–Perron (Paley–Ursell'30), ...

$d \geq 2$: ergodicity, subadditivity $\Rightarrow$

$$\lambda_1(D) = \inf_{n \in \mathbb{N}} \frac{1}{n} \int_{\Delta} \log \|D^{(n)}(\mathbf{x})\| \, d\mu(\mathbf{x})$$

gives upper bounds for $\lambda_2(A)$

Hardcastle'02: $\lambda_2(A) \leq -0.0053293$ for modified Jacobi–Perron
(d -dimensional Gauss), $d = 3 \Rightarrow \lambda_2(A) < 0$ for Brun, $d = 3$

Lower bounds?

Adapt method of Pollicott'10 (approximation by periodic orbits) to
matrices with negative entries?

Selmer algorithm (ordered)

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0, x_{d-1} + x_d \geq 1\}$$

$$T_S(x_1, \dots, x_d) = \pi(\text{ord}(1-x_d, x_1, \dots, x_d)) = \text{ord}\left(\frac{x_2}{x_1}, \dots, \frac{x_d}{x_1}, \frac{1-x_d}{x_1}\right)$$

$$A_S(\mathbf{x}) = \begin{cases} S_a & \text{if } x_d > 1/2, \\ S_b & \text{if } x_d < 1/2, \end{cases}$$

$$S_a = \begin{pmatrix} 0 & 1 & 0 & \cdots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 & 0 \\ 1 & 0 & \cdots & 0 & 1 \\ 1 & 0 & \cdots & 0 & 0 \end{pmatrix}, \quad S_b = \begin{pmatrix} 0 & 1 & 0 & \cdots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 & 0 \\ 1 & 0 & \cdots & 0 & 0 \\ 1 & 0 & \cdots & 0 & 1 \end{pmatrix}$$

invariant measure μ_S has density $\frac{c}{x_1 \cdots x_d}$

Selmer approximation cocycle

$$D_S(\mathbf{x}) = \begin{pmatrix} 0 & 1 & 0 & \cdots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 & 0 \\ -x_1 & -x_2 & \cdots & -x_{d-1} & 1 - x_d \\ -x_1 & -x_2 & \cdots & -x_{d-1} & -x_d \end{pmatrix} \quad \text{if } x_d > 1/2$$

$$D_S(\mathbf{x}) = \begin{pmatrix} 0 & 1 & 0 & \cdots & 0 \\ \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & 1 & 0 \\ -x_1 & -x_2 & \cdots & -x_{d-1} & -x_d \\ -x_1 & -x_2 & \cdots & -x_{d-1} & 1 - x_d \end{pmatrix} \quad \text{if } x_d < 1/2$$

Selmer algorithm, $d = 2$

$$D_S^{(2)}(\mathbf{x}) = \begin{cases} \begin{pmatrix} 1 - x_1 & -x_2 \\ 1 & 0 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_a^2 = \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \end{pmatrix} \\ \begin{pmatrix} 1 - x_1 & 1 - x_2 \\ 1 & 0 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_a S_b = \begin{pmatrix} 1 & 0 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix} \\ \begin{pmatrix} 1 & 0 \\ 1 - x_1 & -x_2 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_b S_a = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 1 & 0 \end{pmatrix} \\ \begin{pmatrix} 1 & 0 \\ 1 - x_1 & 1 - x_2 \end{pmatrix} & \text{if } A_S^{(2)}(\mathbf{x}) = S_b^2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 1 & 1 \end{pmatrix} \end{cases}$$

$$\|D_S^{(2)}(\mathbf{x})\|_\infty = 1 \text{ for all } \mathbf{x} \in \Delta \Rightarrow \lambda_2(A_S) \leq 0$$

$$\|D_S^{(4)}(\mathbf{x})\|_\infty < 1 \text{ for some } \mathbf{x} \in \Delta \Rightarrow \lambda_2(A_S) < 0$$

$$\lambda_2(A_S) \leq \frac{1}{42} \int_{\Delta} \log \|D_S^{(42)}(\mathbf{x})\| d\mu(\mathbf{x}) < -0.050393$$

Selmer matrices

Theorem (Berthé–St-Thuswaldner)

Let $d = 2$ and $M \in \{S_a, S_b\}^n$ for some $n \geq 1$. The following are equivalent.

- 1. M is a primitive matrix,*
- 2. M is a Pisot matrix,*
- 3. $M^2 \notin \{S_a S_b, S_b^2\}^n$.*

cf. results of Avila–Delecroix'15 (arXiv:1506.03692) for Brun matrices ($d = 2$) and Arnoux–Rauzy matrices ($d \geq 2$)

Selmer algorithm, computer simulations

Calculating $D_S^{(n)}(\mathbf{x})$ and $A_S^{(n)}(\mathbf{x})$ for $n = 2^{30}$ for randomly chosen points $\mathbf{x} \in \Delta$ gives the following estimates for $d \in \{2, 3, 4, 5\}$ (without guaranteed accuracy)

d	$\lambda_2(A_S)$	$1 - \frac{\lambda_2(A_S)}{\lambda_1(A_S)}$
2	-0.07072	1.3871
3	-0.02283	1.1444
4	+0.00176	0.9866
5	+0.01594	0.8577

Lagarias'93: uniform approximation exponent

$$\sup \left\{ \delta : \left\| \mathbf{x} - \frac{\mathbf{p}_i^{(n)}}{q_i^{(n)}} \right\| = O\left(\frac{1}{(q_i^{(n)})^\delta}\right), 0 \leq i \leq d \right\} = 1 - \frac{\lambda_2(A)}{\lambda_1(A)} \quad (\text{a.e. } \mathbf{x} \in \Delta)$$

Computer simulations

We calculate $D^{(n)}(\mathbf{x})$ for randomly chosen points $\mathbf{x} \in \Delta$ by a C program with double precision floating point arithmetic. We have to renormalize the matrices before they get too small or too large. For $k|n$, e.g., $k = 2^{10}$, $n = 2^{30}$, we have

$$\begin{aligned} & \frac{1}{D_{1,1}^{(n)}(\mathbf{x})} D^{(n)}(\mathbf{x}) \\ &= \frac{D_{1,1}^{(n-k)}(\mathbf{x})}{D_{1,1}^{(n)}(\mathbf{x})} D^{(k)}(T^{n-k}\mathbf{x}) \cdots \frac{D_{1,1}^{(k)}(\mathbf{x})}{D_{1,1}^{(2k)}(\mathbf{x})} D^{(k)}(T^k\mathbf{x}) \frac{1}{D_{1,1}^{(k)}(\mathbf{x})} D^{(k)}(\mathbf{x}) \\ \log |D_{1,1}^{(n)}(\mathbf{x})| &= \log \frac{|D_{1,1}^{(n)}(\mathbf{x})|}{|D_{1,1}^{(n-k)}(\mathbf{x})|} + \cdots + \log \frac{|D_{1,1}^{(2k)}(\mathbf{x})|}{|D_{1,1}^{(k)}(\mathbf{x})|} + \log |D_{1,1}^{(k)}(\mathbf{x})| \\ (D_{1,1}^{(n)}(\mathbf{x})) & \text{ is the top left coefficient of } D^{(n)}(\mathbf{x}) \end{aligned}$$

Jacobi–Perron algorithm, computer simulations

$$A_{\text{JP}}(x_1, \dots, x_d) = \begin{pmatrix} \lfloor \frac{1}{x_1} \rfloor & 1 & \lfloor \frac{x_2}{x_1} \rfloor & \dots & \lfloor \frac{x_d}{x_1} \rfloor \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \vdots & & \ddots & 1 \\ 1 & 0 & \dots & \dots & 0 \end{pmatrix}$$

$$D_{\text{JP}}(x_1, \dots, x_d) = \begin{pmatrix} 0 & 1 & \dots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & 1 \\ -x_1 & \dots & -x_{d-1} & -x_d \end{pmatrix}$$

d	$\lambda_2(A_{\text{JP}})$	$1 - \frac{\lambda_2(A_{\text{JP}})}{\lambda_1(A_{\text{JP}})}$	d	$\lambda_2(A_{\text{JP}})$	$1 - \frac{\lambda_2(A_{\text{JP}})}{\lambda_1(A_{\text{JP}})}$
2	-0.44841	1.3735	7	-0.02819	1.0243
3	-0.22788	1.1922	8	-0.01470	1.0127
4	-0.13062	1.1114	9	-0.00505	1.0044
5	-0.07880	1.0676	10	+0.00217	0.9981
6	-0.04798	1.0413	11	+0.00776	0.9933

$$D_{d,j}^{(n)}(\mathbf{x}) = - \sum_{i=1}^d x_{d-i+1}^{(n)} D_{d,j}^{(n-i)}(\mathbf{x}) \quad \text{with } (x_1^{(n)}, \dots, x_d^{(n)}) = T^n(\mathbf{x})$$

Results for random recurrences or products of random matrices?

Nearest integer Jacobi–Perron algorithm

$$\Delta = [-1/2, 1/2]^d$$

(coefficients can be negative)

$$A_{\text{NIJP}}(x_1, \dots, x_d) = \begin{pmatrix} \lfloor \frac{1}{x_1} + \frac{1}{2} \rfloor & 1 & \lfloor \frac{x_2}{x_1} + \frac{1}{2} \rfloor & \dots & \lfloor \frac{x_d}{x_1} + \frac{1}{2} \rfloor \\ 0 & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \vdots & & \ddots & 1 \\ 1 & 0 & \dots & \dots & 0 \end{pmatrix}$$

$$D_{\text{NIJP}}(x_1, \dots, x_d) = \begin{pmatrix} 0 & 1 & \dots & 0 \\ \vdots & \ddots & \ddots & \vdots \\ 0 & \dots & 0 & 1 \\ -x_1 & \dots & -x_{d-1} & -x_d \end{pmatrix}$$

d	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	d	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	d	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$	d	$1 - \frac{\lambda_2(A)}{\lambda_1(A)}$
2	1.40145	6	1.06898	10	1.01737	14	0.99924
3	1.22519	7	1.04944	11	1.01125	15	0.99657
4	1.14373	8	1.03551	12	1.00639		
5	1.09786	9	1.02521	13	1.00246		

Brun and modified Jacobi–Perron algorithms (ordered)

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0\}$$

$$T_B(x_1, \dots, x_d) = \pi(\text{ord}(1 - x_1, x_2, \dots, x_d))$$

$$A_B(\mathbf{x}) = B_k \quad \text{if } x_k > 1 - x_1 > x_{k+1} \quad (0 \leq k \leq d, x_0 = 1, x_{d+1} = 0)$$

$$B_0 = \begin{pmatrix} 1 & 0 & \cdots & \cdots & 0 \\ 1 & 1 & & & \vdots \\ 0 & 0 & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & 1 & 0 \\ 0 & 0 & \cdots & 0 & 1 \end{pmatrix}, \quad B_k = \begin{pmatrix} 1 & 1 & 0 & \cdots & \cdots & \cdots & \cdots & 0 \\ k-1 \left\{ \begin{array}{l} 0 & 0 & 1 & \ddots & & & & \vdots \\ \vdots & 0 & \ddots & \ddots & \ddots & & & \vdots \\ 0 & \vdots & \ddots & \ddots & 1 & \ddots & & \vdots \\ 1 & \vdots & & \ddots & 0 & 0 & \ddots & \vdots \end{array} \right. \\ d-k \left\{ \begin{array}{l} 0 & \vdots & & \ddots & 1 & \ddots & 0 \\ \vdots & \vdots & & & \ddots & \ddots & 0 \\ 0 & 0 & \cdots & \cdots & \cdots & \cdots & 0 & 1 \end{array} \right. \end{pmatrix}$$

$$A_{\text{MJP}}(\mathbf{x}) = A_B^{(n)}(\mathbf{x}) \quad \text{with } n = \min\{k \geq 1 : A_B(T^{k-1}\mathbf{x}) \neq B_0\}$$

d	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$	d	$\lambda_2(A_B)$	$1 - \frac{\lambda_2(A_B)}{\lambda_1(A_B)}$
2	-0.11216	1.3683	7	-0.01210	1.0493
3	-0.07189	1.2203	8	-0.00647	1.0283
4	-0.04651	1.1504	9	-0.00218	1.0102
5	-0.03051	1.1065	10	+0.00115	0.9943
6	-0.01974	1.0746	11	+0.00381	0.9799

Garrity's triangle algorithm (ordered)

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0\}$$

$$T_G(x_1, \dots, x_d) = \begin{cases} \pi(\text{ord}(1 - \sum_{j=1}^k x_j, x_1, \dots, x_d)) & \text{if } 0 < 1 - \sum_{j=1}^k x_j < x_{k+1}, 1 \leq k \leq d-2 \\ \frac{1}{x_1}(x_2, \dots, x_d, 1 - \sum_{j=1}^{d-1} x_j - \ell x_d) & \text{if } 0 < 1 - \sum_{j=1}^{d-1} x_j - \ell x_d < x_d, \ell \geq 0 \end{cases}$$

$d = 2 :$

$$A_G(\mathbf{x}) = \begin{pmatrix} 1 & 1 & 0 \\ \lfloor \frac{1-x_1}{x_2} \rfloor & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}, \quad D_G(\mathbf{x}) = \begin{pmatrix} -\lfloor \frac{1-x_1}{x_2} \rfloor x_1 & 1 - \lfloor \frac{1-x_1}{x_2} \rfloor x_2 \\ -x_1 & -x_2 \end{pmatrix}$$

Problem: $\lfloor \frac{1-x_1}{x_2} \rfloor x_1$ can be large

d	$\lambda_2(A_G)$	$1 - \frac{\lambda_2(A_G)}{\lambda_1(A_G)}$	d	$\lambda_2(A_G)$	$1 - \frac{\lambda_2(A_G)}{\lambda_1(A_G)}$
2	+0.34434	0.6859	7	-0.00644	1.0225
3	+0.37673	0.5798	8	-0.00768	1.0304
4	+0.25232	0.6286	9	-0.00435	1.0189
5	+0.10677	0.7778	10	-0.00074	1.0035
6	+0.01859	0.9468	11	+0.00237	0.9880

Intermediate algorithm between Arnoux–Rauzy and Brun

$$\Delta = \{(x_1, \dots, x_d) \in \mathbb{R}^d : 1 \geq x_1 \geq \dots \geq x_d \geq 0\}$$

$$T_{\text{BST}}(x_1, \dots, x_d) = \pi(\text{ord}(1 - \sum_{j=1}^k x_j, x_1, \dots, x_d))$$

if $0 < 1 - \sum_{j=1}^k x_j < x_{k+1}$ ($1 \leq k \leq d$, $x_{d+1} = 1$)

d	$\lambda_2(A_{\text{BST}})$	$1 - \frac{\lambda_2(A_{\text{BST}})}{\lambda_1(A_{\text{BST}})}$	d	$\lambda_2(A_{\text{BST}})$	$1 - \frac{\lambda_2(A_{\text{BST}})}{\lambda_1(A_{\text{BST}})}$
2	-0.13648	1.3606	7	-0.02033	1.0729
3	-0.10803	1.2430	8	-0.01175	1.0468
4	-0.07540	1.1817	9	-0.00563	1.0246
5	-0.05035	1.1388	10	-0.00114	1.0054
6	-0.03263	1.1034	11	+0.00224	0.9886

Comparison between (simulations of) algorithms

uniform approximation coefficients $1 - \frac{\lambda_2(A)}{\lambda_1(A)} \leq 1 + \frac{1}{d}$

d	Selmer	Brun	JP	BST	Garrity	NIJP
2	1.3871	1.3683	1.3735	1.3606	0.6859	1.40145
3	1.1444	1.2203	1.1922	1.2430	0.5798	1.22519
4	0.9866	1.1504	1.1114	1.1817	0.6286	1.14373
5	0.8577	1.1065	1.0676	1.1388	0.7778	1.09786
6	0.7442	1.0746	1.0413	1.1034	0.9468	1.06898
7	0.6437	1.0493	1.0243	1.0729	1.0225	1.04944
8	0.5561	1.0283	1.0127	1.0468	1.0304	1.03551
9	0.4810	1.0102	1.0044	1.0246	1.0189	1.02521
10	0.4173	0.9943	0.9981	1.0054	1.0035	1.01737
11	0.3636	0.9799	0.9933	0.9886	0.9880	1.01125
12						1.00639
13						1.00246
14						0.99924

bold face: best algorithm or best positive algorithm in dimension d
(among those that we considered)