

Derivatives and distances in probabilistic coherence spaces

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This talk

Try to use the regularity of morphisms in the cartesian closed category $\mathbf{Pcoh}_!$ (they are power series with coefficients ≥ 0 with > 0 radius of convergence).

In particular: these functions have differentials. What is their probabilistic operational meaning? Two (related) answers:

- Expected computation time
- Operational distance

in a probabilistic functional programming language: pPCF.

Syntax of probabilistic pPCF

Syntax:

$$\sigma, \tau, \dots := \iota \mid \sigma \Rightarrow \tau$$

$$M, N, P \dots := \underline{n} \mid \text{succ}(M) \mid \text{pred } M \mid x \mid \text{coin}(r) \mid \text{let}(x, M, N) \\ \mid \text{if}(M, N, P) \mid (M) N \mid \lambda x^\sigma M \mid \text{fix}(M)$$

with $r \in [0, 1]$.

Typing rules

$$\frac{}{\Gamma \vdash \underline{n} : \iota} \quad \frac{}{\Gamma, x : \sigma \vdash x : \sigma}$$

$$\frac{\Gamma \vdash M : \iota}{\Gamma \vdash \text{succ}(M) : \iota} \quad \frac{\Gamma \vdash M : \iota}{\Gamma \vdash \text{pred } M : \iota}$$

$$\frac{\Gamma \vdash M : \iota \quad \Gamma, x : \iota \vdash N : \sigma}{\Gamma \vdash \text{let}(x, M, N) : \sigma} \quad \text{Only for } M \text{ of type } \iota!$$

$$\frac{\Gamma, x : \sigma \vdash M : \tau}{\Gamma \vdash \lambda x^\sigma M : \sigma \Rightarrow \tau} \quad \frac{\Gamma \vdash M : \sigma \Rightarrow \tau \quad \Gamma \vdash N : \sigma}{\Gamma \vdash (M) N : \tau}$$

$$\frac{\Gamma \vdash M : \sigma \Rightarrow \sigma}{\Gamma \vdash \text{fix}(M) : \sigma} \quad \frac{r \in [0, 1]}{\Gamma \vdash \text{coin}(r) : \iota}$$

Intuitions on pPCF

$\vdash M : \iota$ means that M represents a subprobability distribution on the integers.

Example: $\text{coin}(1/3)$ weights $\underline{0}$ with probability $1/3$, $\underline{1}$ with probability $2/3$ and $\underline{n+2}$ with probability 0 .

$\text{let}(x, M, N)$: samples an integer according to M and feeds N through x with the obtained value (a \underline{n} for some $n \in \mathbb{N}$).

$\vdash M : \iota \Rightarrow \iota$ means that M is an (generally non linear) sub-probability distribution transformer.

Reduction rules

We define a weak head reduction strategy.

Deterministic reduction rules

$$\frac{}{(\lambda x^\sigma M) N \rightarrow_d M [N/x]}$$

$$\frac{}{\text{fix}(M) \rightarrow_d (M) \text{fix}(M)}$$

$$\frac{}{\text{succ}(\underline{n}) \rightarrow_d \underline{n+1}}$$

$$\frac{}{\text{if}(\underline{0}, M, N) \rightarrow_d M}$$

$$\frac{}{\text{if}(\underline{n+1}, M, N) \rightarrow_d N}$$

$$\frac{}{\text{let}(x, \underline{n}, N) \rightarrow_d N [\underline{n}/x]}$$

Probabilistic reductions

$$\frac{M \rightarrow_d M'}{M \xrightarrow{1} M'} \quad \frac{}{\text{coin}(p) \xrightarrow{p} \underline{0}} \quad \frac{}{\text{coin}(p) \xrightarrow{1-p} \underline{1}}$$

$$\frac{M \xrightarrow{p} M'}{(M) N \xrightarrow{p} (M') N} \quad \frac{M \xrightarrow{p} M'}{\text{succ}(M) \xrightarrow{p} \text{succ}(M')}$$

$$\frac{M \xrightarrow{p} M'}{\text{let}(x, M, N) \xrightarrow{p} \text{let}(x, M', N)}$$

$$\frac{M \xrightarrow{p} M'}{\text{if}(M, N, P) \xrightarrow{p} \text{if}(M', N, P)}$$

Probability of reduction

Given M such that $\vdash M : \iota$, we can consider all possible reductions from M to a given integer constant \underline{n} :

$$M = M_0 \xrightarrow{p_1} M_1 \xrightarrow{p_2} \dots \xrightarrow{p_k} M_k = \underline{n}$$

Summing up the probabilities $\prod_{i=1}^k p_i$ of all these paths we get the probability that M reduces to \underline{n} , denoted $\Pr(M \downarrow \underline{n})$.

Observational distance

Given M and N such that $\vdash M : \sigma$ and $\vdash N : \sigma$, one defines the *observational distance* $d_{\text{obs}}(M, N)$ between M and N as the sup of all

$$|\text{Pr}((C) M \downarrow \underline{0}) - \text{Pr}((C) N \downarrow \underline{0})|$$

for all possible “contexts” which are closed terms C such that $\vdash C : \sigma \Rightarrow \iota$.

M and N are *observationally equivalent* if this “distance” is 0.

$d_{\text{obs}}(_, _)$ is a distance on the observational classes of closed terms of type σ (for any type σ).

Part I

Reminder — Probabilistic coherence spaces: an “analytic” denotational model

If $u, u' \in (\mathbb{R}_{\geq 0})^I$ then $\langle u, u' \rangle = \sum_{i \in I} u_i u'_i \in \overline{\mathbb{R}_{\geq 0}}$.

If $P \subseteq \overline{\mathbb{R}_{\geq 0}}^I$ then

$$P^\perp = \{u' \in \overline{\mathbb{R}_{\geq 0}}^I \mid \forall u \in P \langle u, u' \rangle \leq 1\}.$$

A probabilistic coherence space (PCS) is a pair $X = (|X|, PX)$

where $PX \subseteq \overline{\mathbb{R}_{\geq 0}}^{|X|}$ such that

- $PX^{\perp\perp} = PX$
- $\forall a \in |X| \exists x \in PX \ x_a > 0$
- $\forall a \in |X| \exists m \in \mathbb{R}_{\geq 0} \forall x \in PX \ x_a \leq m$

So actually $PX \subseteq (\mathbb{R}_{\geq 0})^{|X|}$.

Dual of X : $X^\perp = (|X|, PX^\perp)$ so that $X^{\perp\perp} = X$.

Examples of PCS's.

- $1 = (\{*\}, [0, 1])$ with $1^\perp = 1$ for $0 < r < \infty$.
- $\mathbf{Bool} = 1 \oplus 1 = (\{\mathbf{t}, \mathbf{f}\}, \{(x_{\mathbf{t}}, x_{\mathbf{f}}) \in \mathbb{R}_{\geq 0}^2 \mid x_{\mathbf{t}} + x_{\mathbf{f}} \leq 1\})$
- $\mathbf{Bool}^\perp = 1 \& 1 = (\{\mathbf{t}, \mathbf{f}\}, \{(x_{\mathbf{t}}, x_{\mathbf{f}}) \in \mathbb{R}_{\geq 0}^2 \mid x_{\mathbf{t}}, x_{\mathbf{f}} \leq 1\})$
- $\mathbf{N} = (\mathbb{N}, \{x \in (\mathbb{R}_{\geq 0})^{\mathbb{N}} \mid \sum_{i=0}^{\infty} x_i \leq 1\})$

Linear morphisms in PCS's

Linear morphisms from X to Y : if $t \in (\mathbb{R}_{\geq 0})^{|X| \times |Y|}$ (a matrix) and $x \in PX$ (a vector) then we can apply the matrix to the vector:

$$tx \in (\mathbb{R}_{\geq 0})^{|Y|} \text{ with } (tx)_b = \sum_{a \in |X|} t_{a,b} x_a$$

Then t is a linear morphism from X to Y ($t \in \mathbf{Pcoh}(X, Y)$) if $\forall x \in PX \ tx \in PY$. This defines a category \mathbf{Pcoh} , a model of LL (with all fixpoints of types and term fixpoint operators at all types).

Fact

$\mathbf{Pcoh}(X, Y)$ is a PCS structure, that is $\mathbf{Pcoh}(X, Y) = P(X \multimap Y)$ for a PCS $X \multimap Y$ with $|X \multimap Y| = |X| \times |Y|$.

$P(\mathbb{N} \multimap \mathbb{N})$: sub-stochastic matrices on $\mathbb{N} \times \mathbb{N}$.

Non-linear morphisms in PCS's

The non-linear morphisms $X \rightarrow Y$ are the elements of $P(!X \multimap Y)$ where $!X$ is a PCS with $|!X| = \mathcal{M}_{\text{fin}}(|X|)$ (finite multisets).

If $x \in PX$ and $m \in |!X|$ define

$$x^m = \prod_{a \in |X|} x_a^{m(a)}.$$

Then $t \in P(!X \multimap Y) = \mathbf{Pcoh}_!(X, Y)$ is characterized by

$$\forall x \in PX \quad \widehat{t}(x) = \left(\sum_{m \in |!X|} t_{m,b} x^m \right)_{b \in |Y|} \in PY$$

$\mathbf{Pcoh}_!$ is a model of probabilistic pPCF. Morphisms are functions:

$$\widehat{t \circ s} = \widehat{t} \circ \widehat{s} \text{ and } (\forall x \in PX \widehat{s}(x) = \widehat{t}(x)) \Rightarrow s = t.$$

pPCF interpretation

Fact

$\mathbf{Pcoh}_!$ is a cartesian closed category with an object of integers \mathbb{N} and least fixpoints operators $(X \Rightarrow X) \rightarrow X$ for all X . An thus it is a model of pPCF.

$\llbracket \iota \rrbracket = \mathbb{N}$ and $\llbracket \sigma \Rightarrow \tau \rrbracket = !\llbracket \sigma \rrbracket \multimap \llbracket \tau \rrbracket$

If $\Gamma \vdash M : \sigma$ with $\Gamma = (x_1 : \sigma_1, \dots, x_k : \sigma_k)$ then

$$\llbracket M \rrbracket_{\Gamma} \in \mathbf{Pcoh}_!(\llbracket \sigma_1 \rrbracket \& \dots \& \llbracket \sigma_k \rrbracket, \llbracket \sigma \rrbracket)$$

so $\llbracket M \rrbracket_{\Gamma}$ can be seen as a function $\prod_{i=1}^k P[\llbracket \sigma_i \rrbracket] \rightarrow P[\llbracket \sigma \rrbracket]$.

Why no negative coefficients in power series?

Seems crucial for combining fixpoints and power series. Assume e.g. that we admit the “weak parallel or” function

$$\text{wpor} : [0, 1] \times [0, 1] \rightarrow [0, 1] \quad (u, v) \mapsto u + v - uv$$

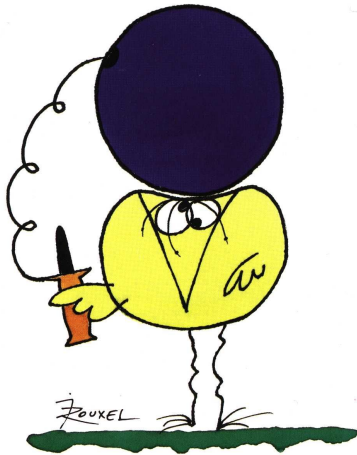
Spawns two threads, stops as soon as one of them stops.

Add it to pPCF. Then we can define $\vdash P : 1 \Rightarrow 1$ by

$$P = \text{fix}(\lambda f^{\iota \Rightarrow \iota} \lambda x^{\iota} \text{wpor}(x, (f) x)).$$

Spawns an unbounded number of copies of x , stops as soon as one of them does. Then $\llbracket P \rrbracket(0) = 0$ and $\llbracket P \rrbracket(u) = 1$ for $u > 0$.
Scott continuous, but far from being analytic!

Les devises Shadok



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Example of term interpretations

- $\llbracket x_i \rrbracket_{\Gamma}(\vec{u}) = u_i$
- $\llbracket \underline{n} \rrbracket_{\Gamma}(\vec{u}) = e_n$ where $(e_n)_j = \delta_{n,j} \in \text{PN}$.
- $\llbracket \text{if}(M, N, P) \rrbracket_{\Gamma} =$
 $\llbracket M \rrbracket_{\Gamma}(\vec{u})_0 \llbracket N \rrbracket_{\Gamma}(\vec{u}) + (\sum_{n=1}^{\infty} \llbracket M \rrbracket_{\Gamma}(\vec{u})_n) \llbracket P \rrbracket_{\Gamma}(\vec{u})$
- $\llbracket \text{coin}(p) \rrbracket_{\Gamma}(\vec{u}) = p e_0 + (1 - p) e_1$
- $\llbracket \text{let}(x, M, N) \rrbracket_{\Gamma}(\vec{u}) = \sum_{n=0}^{\infty} \llbracket M \rrbracket_{\Gamma}(\vec{u})_n \llbracket N \rrbracket_{\Gamma, x:\iota}(\vec{u}, e_n)$

Main properties of this interpretation

Fact

For all M with $\vdash M : \iota$ and $n \in \mathbb{N}$, we have $\Pr(M \downarrow \underline{n}) = \llbracket M \rrbracket_n$.

As a consequence

Fact (adequacy, Danos and E.)

For all M, N such that $\vdash M : \sigma$ and $\vdash N : \sigma$, we have $\llbracket M \rrbracket = \llbracket N \rrbracket \Rightarrow d_{\text{obs}}(M, N) = 0$

And also

Fact (full abstraction, Pagani, Tasson and E.)

The converse implication.

Part II

Derivatives and execution (on an example)

In a pPCF extension with unit type 1, for $r \in [0, 1]$

$$M_r = \text{fix}(\lambda f^{1 \Rightarrow 1} \lambda x^1 \text{if}(\text{coin}(r), (f) x; (f) x, x; x))$$

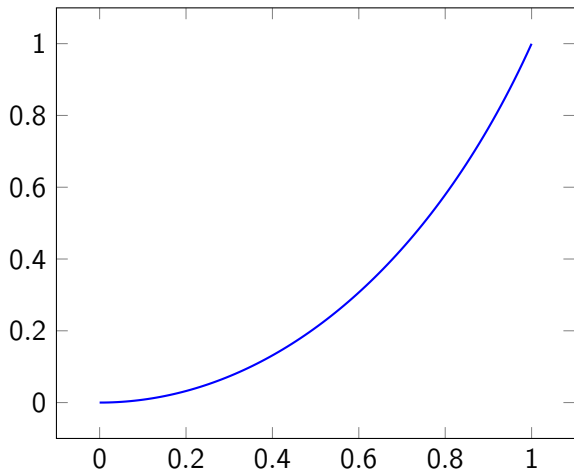
where $()$ is the unique value of type 1 and “;” is the “unit conditional” ($M; N$ reduces to $()$ if both M and N do).

Then $\llbracket M_r \rrbracket$ is a *monotonic* function $\varphi_r : [0, 1] \rightarrow [0, 1]$ minimal such that $\varphi_r(u) = (1 - r)u^2 + r\varphi_r(u)^2$. Hence

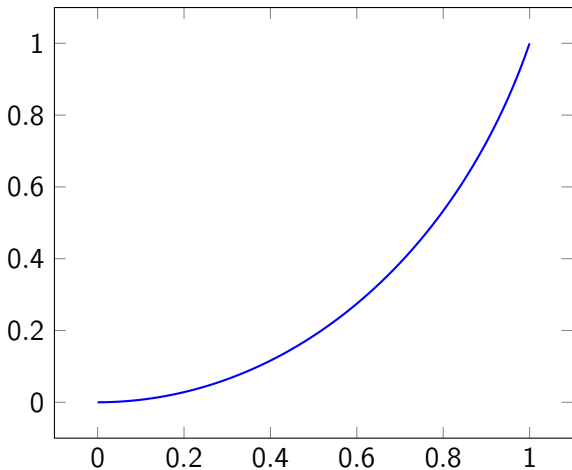
$$\varphi_r(u) = \begin{cases} \frac{1 - \sqrt{1 - 4r(1-r)u^2}}{2r} & \text{if } u > 0 \\ u^2 & \text{if } r = 0. \end{cases}$$

NB: by adequacy, $\varphi_r(1)$ is $\text{Pr}((M_r) () \downarrow ())$

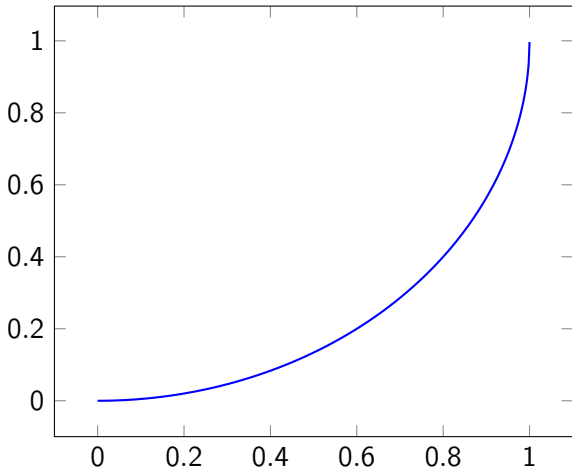
$r = 0.2$ — graph of $\varphi_{0.2}(u) = (1 - \sqrt{1 - 0.64u^2})/0.4$



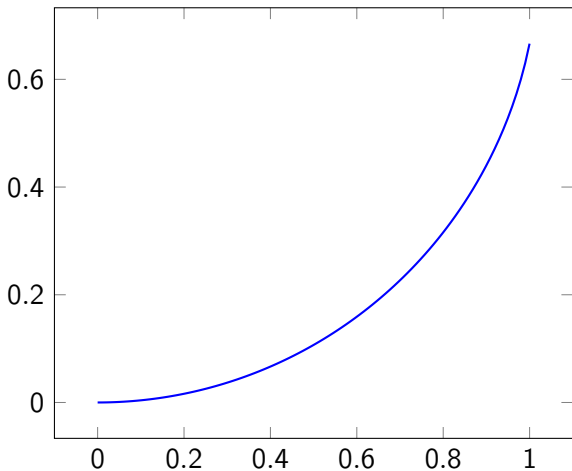
$r = 0.3$ — graph of $\varphi_{0.3}(u) = (1 - \sqrt{1 - 0.84u^2})/0.6$ (steeper slope at $u = 1$)



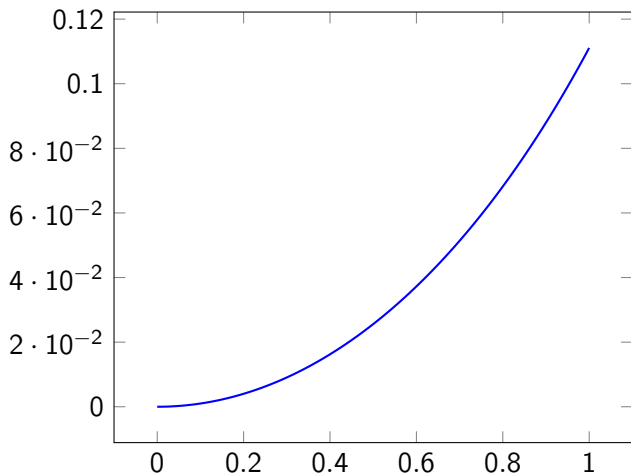
$r = 0.5$ — graph of $\varphi_{0.5}(u) = 1 - \sqrt{1 - u^2}$ (vertical slope at $u = 1$)



$r = 0.6$ — graph of $\varphi_{0.6}(u) = (1 - \sqrt{1 - 0.96u^2})/1.2$ ($\varphi_{0.6}(u) < 1$ but less steep slope at $u = 1$)



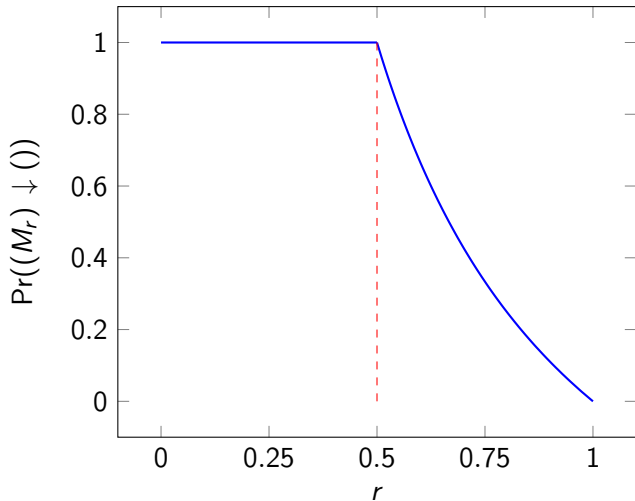
$r = 0.9$ — graph of $\varphi_{0.9}(u) = (1 - \sqrt{1 - 0.36u^2})/1.8$



Probability of termination of (M_r) for $r \in [0, 1]$:

$$\begin{aligned}\varphi_r(1) &= \frac{1 - \sqrt{1 - 4r(1 - r)}}{2r} \\ &= \frac{1 - |1 - 2r|}{2r} \\ &= \begin{cases} 1 & \text{if } r \leq 1/2 \\ \frac{1-r}{r} & \text{if } r > 1/2 \end{cases}\end{aligned}$$

Graph of $\varphi_r(1)$ for $0 \leq r \leq 1$: probability of termination of (M_r) (\cdot).



$$\varphi_r(u) = \sum_{n=0}^{\infty} a_n(r) u^n$$

Fact

$a_n(r) \in [0, 1]$ is the probability that the execution of (M_r) $()$ uses n times the argument $()$ in its reduction to $()$.

$$\varphi_r'(1) = \lim_{u \rightarrow 1^-} \frac{\varphi_r(1) - \varphi_r(u)}{1 - u} = \sum_{n=0}^{\infty} n a_n(r)$$

Fact

$\varphi_r'(1)/\varphi_r(1)$ is the conditional expectation of this execution time, under the condition that the computation terminates.

In the example we can compute this derivative. We have

$$\varphi_r(u) = (1 - r)u^2 + r\varphi_r(u)^2$$

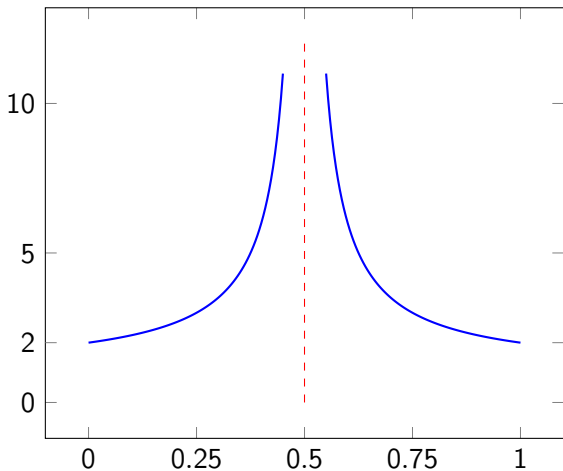
so

$$\varphi_r'(u) = 2(1 - r)u + 2r\varphi_r(u)\varphi_r'(u)$$

The conditional expectation of execution time is

$$\frac{\varphi_r'(1)}{\varphi_r(1)} = \frac{2(1 - r)}{(1 - 2r\varphi_r(1))\varphi_r(1)} = \begin{cases} \frac{2(1-r)}{1-2r} & \text{if } 0 \leq r < 1/2 \\ \frac{2r}{2r-1} & \text{if } 1/2 < r \leq 1 \end{cases}$$

Graph of $\varphi'_r(1)/\varphi_r(1)$: conditional expectation of the number of steps in the reduction of (M_r) for $0 \leq r \leq 1$.



Part III

Lipschitz property and distances

Amplification of probabilities

Fact

$$\forall \varepsilon \in [0, 1] \quad \varepsilon > 0 \Rightarrow d_{\text{obs}}(\text{coin}(0), \text{coin}(\varepsilon)) = 1$$

Take $C = \text{fix}(\lambda f^{\iota \Rightarrow \iota} \lambda x^{\iota} \text{if}(x, \underline{0}, (f) x))$ then

$$\Pr((C) \text{ coin}(0) \downarrow \underline{0}) = 0$$

$$\Pr((C) \text{ coin}(\varepsilon) \downarrow \underline{0}) = 1 \quad \forall \varepsilon > 0$$

The denotational distance

There is a “norm” on PX :

$$\|x\|_X = \sup\{\langle x, x' \rangle \mid x' \in PX^\perp\} \in [0, 1]$$

For instance $\|x\|_N = \sum_{n=0}^{\infty} x_n$.

Also PX is a lattice: $x \wedge y \in PX$ defined pointwise.

Then

$$d_X(x, y) = \|x - (x \wedge y)\|_X + \|y - (x \wedge y)\|_X.$$

defines a distance on PX .

Remember that $d_{\text{obs}}(\text{coin}(0), \text{coin}(\varepsilon)) = 1$ if $\varepsilon > 0$.

On the other hand $\llbracket \text{coin}(0) \rrbracket = e_1$ and $\llbracket \text{coin}(\varepsilon) \rrbracket = \varepsilon e_0 + (1 - \varepsilon)e_1$.
Also $e_1 \wedge (\varepsilon e_0 + (1 - \varepsilon)e_1) = (1 - \varepsilon)e_1$.

$$\begin{aligned}d_N(\llbracket \text{coin}(0) \rrbracket, \llbracket \text{coin}(\varepsilon) \rrbracket) &= \|e_1 - (1 - \varepsilon)e_1\|_N \\ &\quad + \|(1 - \varepsilon)e_1 + \varepsilon e_0 - (1 - \varepsilon)e_1\|_N \\ &= 2\varepsilon.\end{aligned}$$

Remember that by Full Abstraction, if $\vdash M : \sigma$ and $\vdash N : \sigma$,

$$d_{\llbracket \sigma \rrbracket}(\llbracket M \rrbracket, \llbracket N \rrbracket) = 0 \Leftrightarrow d_{\text{obs}}(M, N) = 0$$

we would like to say something not completely trivial in the case $d_{\text{obs}}(M, N) \neq 0$ by limiting the space of observation contexts (in the spirit of the work of Dal Lago and *al.* on probabilistic distances).

The local PCS

Given $x \in PX$, there is a PCS X_x such that

$$P(X_x) = \{u \in PX \mid x + u \in PX\}$$

We have $|X_x| = \{a \in |X| \mid \exists \varepsilon > 0 \ x + \varepsilon e_a \in PX\}$.

It is a PCS (not completely obvious).

This is the *local PCS* of X at x .

[$e_a \in (\mathbb{R}_{\geq 0})^{|X|}$ defined by $(e_a)_{a'} = \delta_{a,a'}$.]

Local derivatives

Let $t \in \mathbf{Pcoh}_!(X, Y)$ and $x \in PX$

Given $u \in P(X_x)$, we know that $x + u \in PX$ and hence we can compute $\widehat{t}(x + u) \in PY$:

$$\begin{aligned}\widehat{t}(x + u)_b &= \sum_{m \in |X|} t_{m,b}(x + u)^m \\ &= \sum_{m \in |X|} t_{m,b} \sum_{p \leq m} \binom{m}{p} x^{m-p} u^p.\end{aligned}$$

where

$$\binom{m}{p} = \prod_{a \in |X|} \binom{m(a)}{p(a)}$$

So keeping only the terms which are constant (p empty) and linear (p singleton) in u :

$$\hat{t}(x) + \sum_{a \in |X|} u_a \sum_{m \in |X|} (m(a) + 1) t_{m+[a], b} x^m \leq \hat{t}(x + u) \in PY$$

that is

$$\sum_{a \in |X|} \left(\sum_{m \in |X|} (m(a) + 1) t_{m+[a], b} x^m \right) u_a \in P(X_{\hat{t}(x)})$$

So we define the “Jacobian matrix” of t at x by

$$t'(x)_{a,b} = \sum_{m \in |X|} (m(a) + 1) t_{m+[a],b} x^m$$

and we have seen that

$$t'(x) \in P(X_x, Y_{\widehat{t}(x)}).$$

Fact (chain rule)

Let $s \in \mathbf{Pcoh}_!(X, Y)$ and $t \in \mathbf{Pcoh}_!(Y, Z)$. Let $x \in PX$ and $u \in PX_x$. Then we have $(t \circ s)'(x) u = t'(\widehat{s}(x)) s'(x) u$.

A Lipschitz property

We take some $p \in [0, 1)$. Let $t \in \mathbf{Pcoh}_!(X, 1)$.

Fact (basic observation)

If $x \in PX$ and $\|x\|_X \leq p$, then $\forall u \in PX$, one has

$$\|x + (1 - p)u\|_X \leq \|x\|_X + (1 - p)\|u\|_X \leq 1$$

so $(1 - p)u \in P(X_x)$.

So if $w \in P(X_x \multimap Y)$, we have $\forall u \in PX$ $\|w(1 - p)u\|_Y \leq 1$ and hence

$$(1 - p)w \in P(X \multimap Y).$$

In particular

$$(1 - p)t'(x) \in P(X \rightarrow 1).$$

Let $x \leq y \in PX$ such that $\|y\|_X \leq p$. Observe that $2 - p > 1$.

Then $x + (2 - p)(y - x) = y + (1 - p)(y - x) \in PX$ since $\|y - x\|_X \leq \|y\|_X \leq p \leq 1$.

Since $x + (2 - p)(y - x) \in PX$ we can define

$$h : [0, 2 - p] \rightarrow [0, 1] \\ \theta \mapsto \widehat{t}(x + \theta(y - x))$$

We have $h \in \mathbf{Pcoh}_!([0, 2 - p], [0, 1])$ by compositionality.

Remember that $2 - p > 1$ so $1 \in [0, 2 - p]$. By Chain Rule $\forall \theta \in [0, 1]$, $h'(\theta) = t'(x + \theta(y - x))(y - x)$.

Since $\|x + \theta(y - x)\|_X \leq \|y\|_X \leq p$ we have

$$\forall \theta \in [0, 1] \quad (1 - p)t'(x + \theta(y - x)) \in P(X \multimap 1).$$

Hence

$$\forall \theta \in [0, 1] \quad h'(\theta) = t'(x + \theta(y - x))(y - x) \leq \frac{\|y - x\|_X}{1 - p}.$$

We have

$$\begin{aligned} 0 \leq \widehat{t}(y) - \widehat{t}(x) &= h(1) - h(0) \\ &= \int_0^1 h'(\theta) d\theta \leq \frac{\|y - x\|_X}{1 - p}. \end{aligned}$$

From this we deduce easily

Fact (Lipschitz property of non-linear morphisms)

Let $p \in [0, 1)$ and $t \in \mathbf{Pcoh}_!(X, 1)$. Let $x, y \in PX$ with $\|x\|_X, \|y\|_X \leq p$. Then

$$|\widehat{t}(x) - \widehat{t}(y)| \leq \frac{d_X(x, y)}{1 - p}$$

A syntactic consequence

If $\Gamma \vdash C : \sigma \Rightarrow 1$ let $C^{(p)}$ with $\vdash C^{(p)} : \sigma \Rightarrow 1$ be

$$C^{(p)} = \lambda x^\sigma (C) \text{ if } (\text{coin}(p), M, \Omega^\sigma).$$

So that $\llbracket (C^{(p)}) M \rrbracket = \widehat{\llbracket C \rrbracket}(p \llbracket M \rrbracket)$ when $\vdash M : \sigma$. If $\vdash M : \sigma$ and $\vdash N : \sigma$.

Definition (p -tamed observational distance)

$$d_{\text{obs}}^{(p)}(M, N) = \sup \left\{ \left| \Pr \left((C^{(p)}) M \downarrow \underline{0} \right) - \Pr \left((C^{(p)}) N \downarrow \underline{0} \right) \right| \mid \vdash C : \sigma \Rightarrow \iota \right\}$$

$$\begin{aligned}
& d_{\text{obs}}^{\langle p \rangle}(M, N) \\
&= \sup\left\{ \left| \widehat{[[C]]}(p[[M]])_0 - \widehat{[[C]]}(p[[N]])_0 \right| \mid \vdash C : \sigma \Rightarrow \iota \right\} \\
&\leq \sup\left\{ \left| \widehat{t}(p[[M]]) - \widehat{t}(p[[N]]) \right| \mid t \in \mathcal{P}(![[\sigma]] \multimap 1) \right\} \\
&\leq \frac{d_{[[\sigma]]}(p[[M]], p[[N]])}{1 - p} = \frac{p}{1 - p} d_{[[\sigma]]}([[M]], [[N]]).
\end{aligned}$$

We have proven:

Fact (metric adequacy of PCS's)

If $\vdash M : \sigma$, $\vdash N : \sigma$ and $0 \leq p < 1$, then

$$d_{\text{obs}}^{\langle p \rangle}(M, N) \leq \frac{p}{1-p} d_{\llbracket \sigma \rrbracket}(\llbracket M \rrbracket, \llbracket N \rrbracket)$$

So for instance

$$d_{\text{obs}}^{\langle p \rangle}(\text{coin}(0), \text{coin}(\varepsilon)) \leq \frac{2\varepsilon p}{1-p}$$

Perspectives

- Understand better what happens on the border of PCS's (elements such that $\|x\| = 1$)
- Extend to “continuous types” (using positive cones and Crubillé's Theorem on stable functions on positive cones)
- A differential pPCF and probabilistic LL? What is the proof-theoretic status of these local derivatives?