

Learning Automata over Large Alphabets

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Outline

Learning Languages

The L* Algorithm

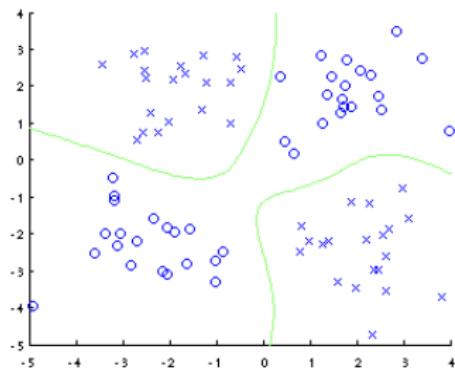
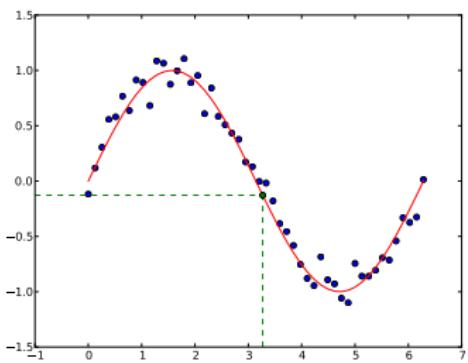
Learning over Large Alphabets

Learning with/without a Teacher

Conclusions

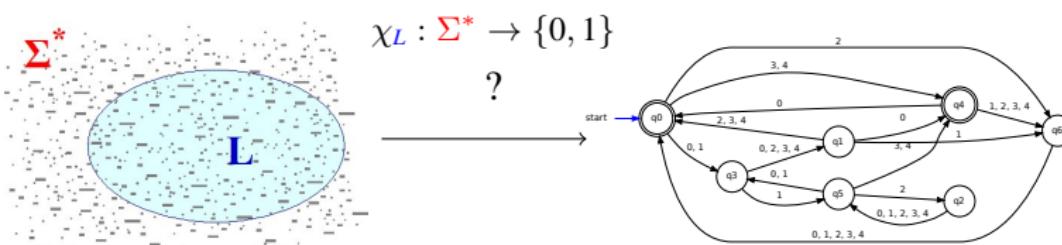
Machine Learning in General

- Given a sample $M = \{(x, y) \mid x \in X, y \in Y\}$
- Find $f : X \rightarrow Y$ such that $f(x) = y, \forall (x, y) \in M$
- Predict $f(x)$ for all $x \in X$



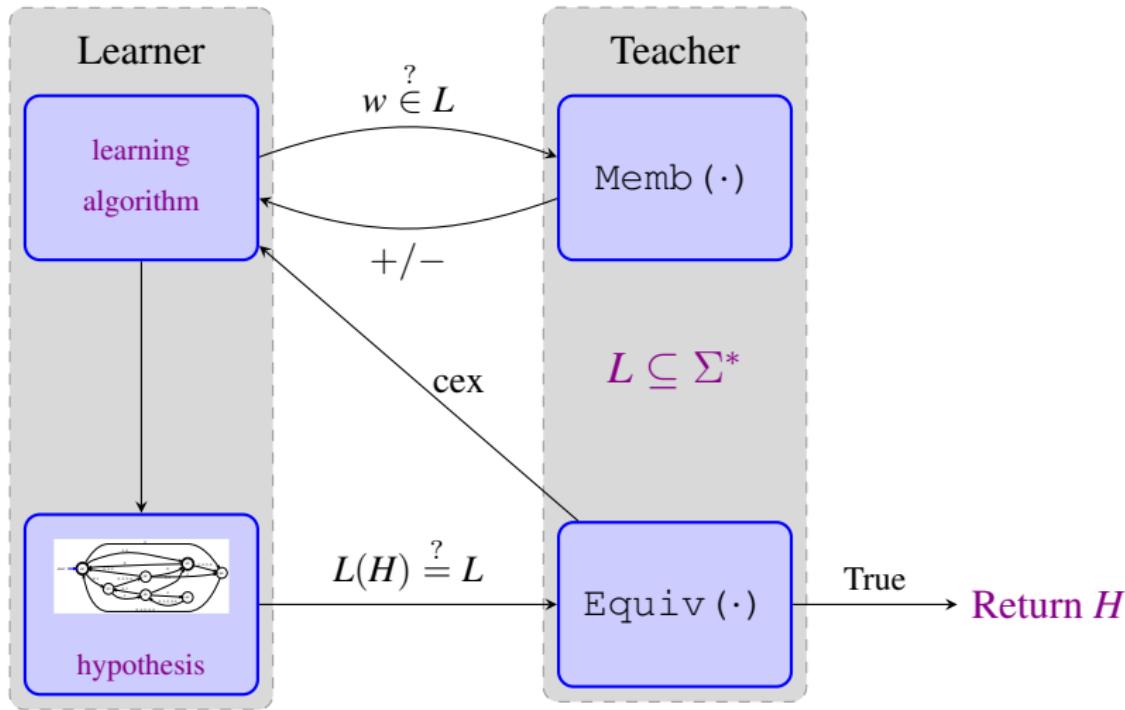
Learning Regular Languages

Let Σ be an alphabet and let $L \subseteq \Sigma^*$ be a regular language (*the target language*)



- Edward F Moore, *Gedanken-experiments on sequential machines*, 1956
- E. Mark Gold, *System Identification via State Characterization*, 1972
- Dana Angluin, *Learning regular sets from queries and counterexamples*, 1987

Active Learning



Regular Sets and their Syntactic Congruences

Equivalence relation \sim_L on Σ^* induced by $L \subseteq \Sigma^*$

$$u \sim_L v \text{ iff } \forall w \in \Sigma^* u \cdot w \in L \Leftrightarrow v \cdot w \in L$$

This relation is a right-congruence with respect to concatenation

$$u \sim v \text{ implies } u \cdot w \sim v \cdot w \text{ for all } u, v, w \in \Sigma^*$$

- $[u]$ is the equivalence class of u
- Σ^*/\sim is the set of all equivalence classes

Theorem (Myhill-Nerode)

The language L is regular iff \sim_L has finitely many congruence classes

Canonical Representation

The minimal automaton for L is $\mathcal{A}_L = (\Sigma, Q, q_0, \delta, F)$ where

- $Q = \Sigma^*/\sim$
- $q_0 = [\epsilon]$
- $\delta([u], a) = [u \cdot a]$
- $F = \{[u] : u \cdot \epsilon \in L\}$

\mathcal{A}_L is homomorphic to any other automaton accepting L

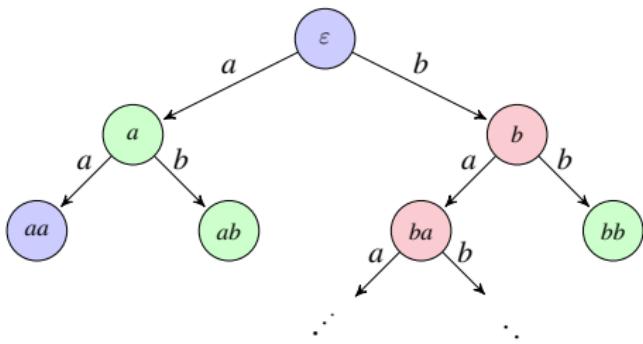
Observation Table - $T = (\Sigma, S, R, E, f)$

definition and properties

		<i>E</i>	
		ε	<i>a</i>
<i>S</i>	ε	-	+
	<i>a</i>	+	-
	<i>b</i>	-	-
	<i>ba</i>	-	-

		<i>E</i>	
		ε	<i>a</i>
<i>R</i>	aa	-	+
	ab	+	-
	bb	+	-
	baa	-	-
	bab	+	-

- *S* states of the canonical automaton
- The words/paths correspond to a spanning tree
- *R* cross- and back-edges/transitions



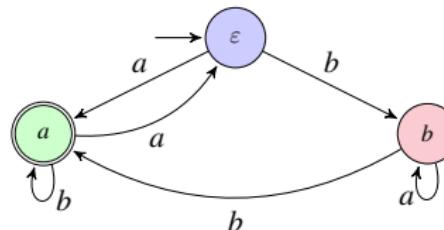
Observation Table - $T = (\Sigma, S, R, E, f)$

definition and properties

		E	
		ε	a
S	ε	-	+
	a	+	-
	b	-	-
	ba	-	-

		E	
		ε	a
R	aa	-	+
	ab	+	-
	bb	+	-
	baa	-	-
	bab	+	-

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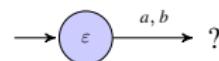
L^* Example ($\Sigma = \{a, b\}$)

L* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	—
a	
b	

hypothesis automaton

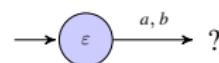


L* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	-
a	+
b	-

hypothesis automaton

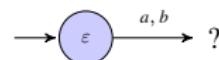


L* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	—
a	+
b	—

hypothesis automaton

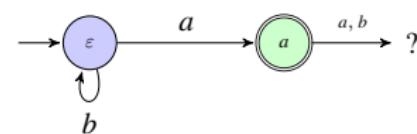


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	-
a	+
b	-

hypothesis automaton

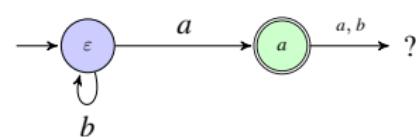


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	-
a	+
b	-
aa	
ab	

hypothesis automaton

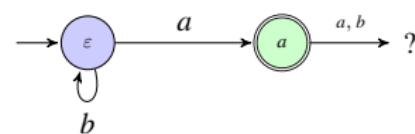


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	-
a	+
b	-
aa	-
ab	+

hypothesis automaton

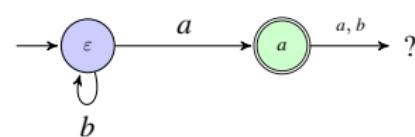


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	-
a	+
b	-
aa	-
ab	+

hypothesis automaton

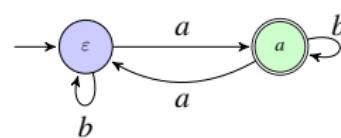


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	-
a	+
b	-
aa	-
ab	+

hypothesis automaton

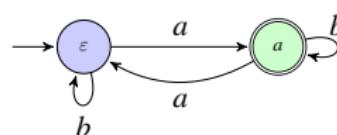


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε
ε	-
a	+
<hr/>	<hr/>
b	-
aa	-
ab	+

hypothesis automaton



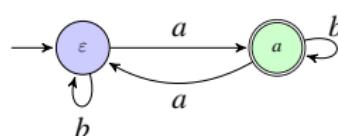
Ask Equivalence Query:
 counterexample: $\textcolor{red}{-ba}$
 $\textcolor{red}{a} \not\sim \textcolor{red}{ba} \longrightarrow \textcolor{red}{a}$ is a new distinguishing string

L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε	a
ε	—	
a	+	
<hr/>	<hr/>	<hr/>
b	—	
aa	—	
ab	+	

hypothesis automaton



counterexample: $-ba$

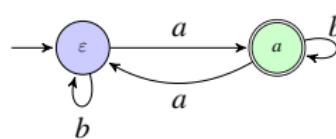
$a \not\sim ba \rightarrow a$ is a new distinguishing string

L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε	a
ε	-	+
a	+	-
b	-	-
aa	-	+
ab	+	-

hypothesis automaton



counterexample: $-ba$

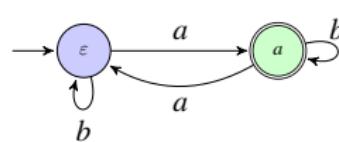
$a \not\sim ba \rightarrow a$ is a new distinguishing string

L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε	a
ε	-	+
a	+	-
b	-	-
aa	-	+
ab	+	-

hypothesis automaton



counterexample: $-ba$

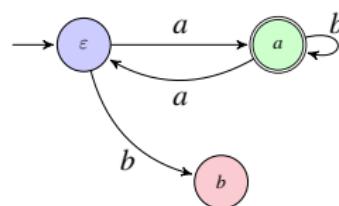
$a \not\sim ba \rightarrow a$ is a new distinguishing string

L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε	a
ε	-	+
a	+	-
b	-	-
aa	-	+
ab	+	-
ba		
bb		

hypothesis automaton

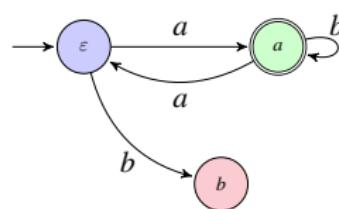


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε	a
ε	—	+
a	+	—
b	—	—
	aa	—
aa	—	+
ab	+	—
ba	—	—
bb	+	—

hypothesis automaton

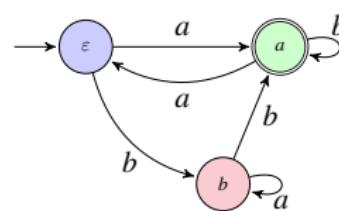


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε	a
ε	—	+
a	+	—
b	—	—
	aa	—
aa	—	+
ab	+	—
ba	—	—
bb	+	—

hypothesis automaton

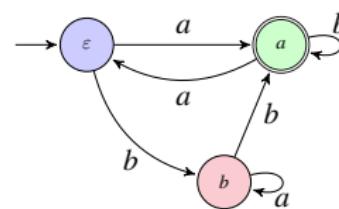


L^* Example ($\Sigma = \{a, b\}$)

observation table

	ε	a
ε	—	+
a	+	—
b	—	—
aa	—	+
ab	+	—
ba	—	—
bb	+	—

hypothesis automaton



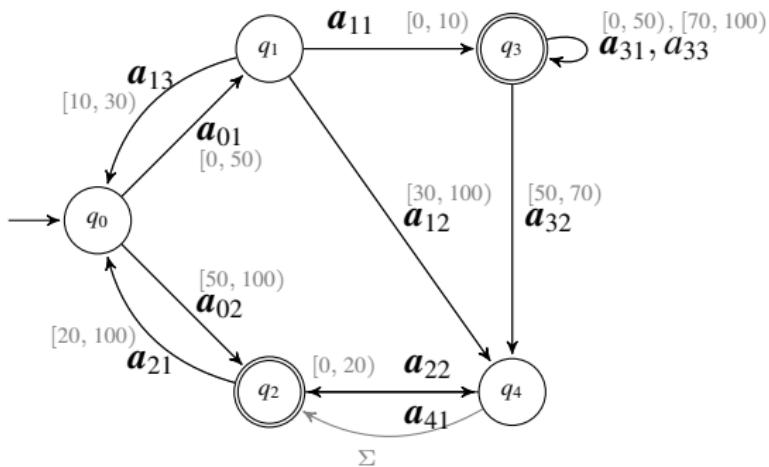
Ask Equivalence Query:

True

Languages over Large Alphabets

- Traditionally automata theory is flat, based on small alphabets, e.g. $\{a, b\}$
- In verification, for example, we have sequences over a huge state-space like \mathbb{B}^n for very large n
- Or we want to have languages over numbers or vectors
- We use symbolic automata with a modest number of states
- We do not want to enumerate all transitions but represent them symbolically using predicates on the alphabet
- We will use inequalities (intervals) for numbers or Boolean functions for Boolean vectors

Symbolic Automata



$$\llbracket a_{01} \rrbracket = [0, 50]$$

$$\Sigma = [0, 100] \subseteq \mathbb{R}$$

$$\llbracket a \rrbracket = \{a \in \Sigma \mid \psi(a) = a\}$$

$$w = 20 \cdot 40 \cdot 60 \quad +$$

$$\mathcal{A} = (\Sigma, \Sigma, \psi, Q, \delta, q_0, F)$$

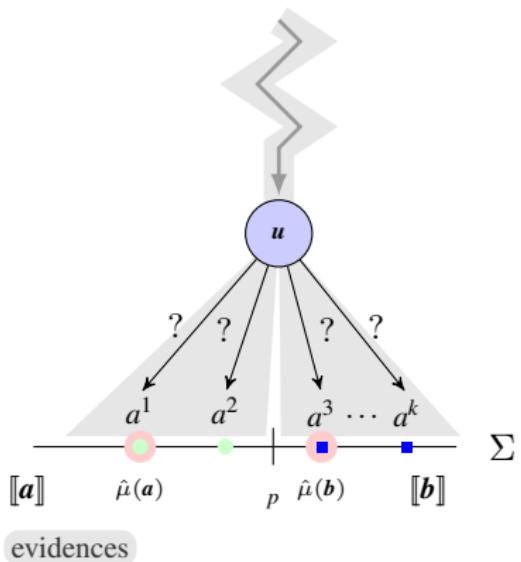
- Q finite set of states,
- q_0 initial state,
- F accepting states,
- Σ large concrete alphabet,
- $\delta \subseteq Q \times 2^\Sigma \times Q$
- Σ finite alphabet (symbols)
- $\psi_q : \Sigma \rightarrow \Sigma_q, q \in Q$

\mathcal{A} is **complete** and deterministic

if $\forall q \in Q \ \{\llbracket a \rrbracket \mid a \in \Sigma_q\}$

forms a *partition of Σ*

Learning using Evidences and Representatives



$$\mu(u \cdot a) = \{\hat{\mu}(u) \cdot a^i \mid a^i \in \llbracket a \rrbracket\}$$

representatives

$$\hat{\mu}(u \cdot a) = \hat{\mu}(u) \cdot \hat{\mu}(a)$$

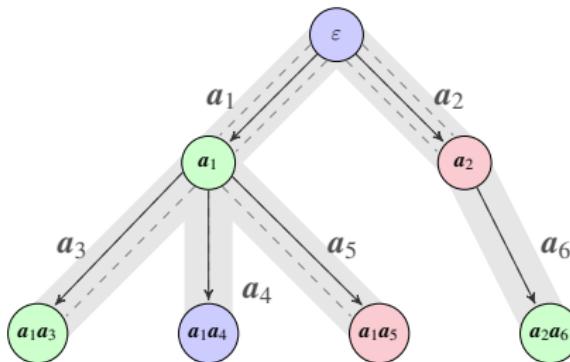
Let Σ be a subset of \mathbb{R}

- To characterize continuations of u , ask queries about $u \cdot a$ for a finite sample of Σ (evidence)
- Evidence can be a fixed set, random, or a result of binary search
- Form **evidence compatible** partitions
 - All evidences within a partition block behave the same
 - Estimate boundaries using *split*, *binary search*,...
- Associate a symbol to each partition block
- Choose one evidence as the **representative** for each new symbol

Symbolic Observation Table - $T = (\Sigma, \Sigma, S, R, \psi, E, f, \mu, \hat{\mu})$

		E		
		ε	a	
S		ε	-	+
R	a_1	+	-	
	a_2	-	-	
		a_1a_3	+	-
		a_1a_4	-	+
		a_1a_5	-	-
		a_2a_6	+	-

- Prefixes are symbolic words
 - Symbols represent sets of letters (“fat” edges)
- Suffixes are concrete words (distinguish states)
- Fill in the table according to the representatives



Symbolic Observation Table - $T = (\Sigma, \Sigma, S, R, \psi, E, f, \mu, \hat{\mu})$

		E	
		ε	a
S		ε	— +
a_1		+	—
a_2		—	—
R		a_1a_3	— +
a_1a_4		—	+
a_1a_5		—	—
a_2a_6		+	—

- $\psi = \{\psi_s\}_{s \in S}, \psi_s : \Sigma \rightarrow \Sigma_s$ semantics
 - $\llbracket a \rrbracket = \{a \in \Sigma \mid \psi(a) = a\}$
- $\mu : \Sigma \rightarrow 2^\Sigma$ evidences
 - $\mu(\varepsilon) = \{\varepsilon\}, \mu(s \cdot a) = \hat{\mu}(s) \cdot \mu(a)$
- $\hat{\mu} : \Sigma \rightarrow \Sigma$ representative
 - $\hat{\mu}(\varepsilon) = \varepsilon, \hat{\mu}(s \cdot a) = \hat{\mu}(s) \cdot \hat{\mu}(a)$
- $f : (S \cup R) \cdot E \rightarrow \{-, +\}$ classif. function
 - $f(s \cdot e) = f(\hat{\mu}(s) \cdot e), f_s(e) = f(s \cdot e)$

Counter-example Treatment (Symbolic Breakpoint)

Proposition

If w is a counter-example to \mathcal{A}_T then there exists an i -factorization of w such that either

$$f(\hat{\mu}(s_{i-1} \cdot a_i) \cdot v_i) \neq f(\hat{\mu}(s_i) \cdot v_i) \quad (1)$$

or

$$f(\hat{\mu}(s_{i-1}) \cdot a_i \cdot v_i) \neq f(\hat{\mu}(s_{i-1}) \cdot \hat{\mu}(a_i) \cdot v_i) \quad (2)$$

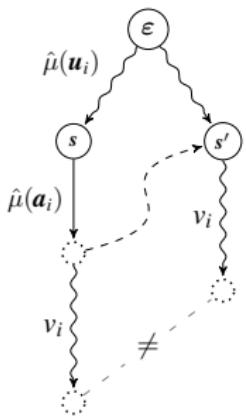
- If (1), then v_i is a new distinguishing word vertical expansion
 - Table not closed \rightarrow new state
- If (2), then a_i is a new evidence for a_i . horizontal expansion
 - Evidence incompatibility \rightarrow new transition / refinement

Counter-example Treatment (Symbolic Breakpoint)

Let $w = a_1 \cdots a_i \cdots a_{|w|} = u_i \cdot a_i \cdot v_i$ be a counter-example.

$$f(\hat{\mu}(\textcolor{violet}{s}_{i-1} \cdot \textcolor{violet}{a}_i) \cdot v_i) \neq f(\hat{\mu}(\textcolor{violet}{s}_i) \cdot v_i)$$

$$s_i = \delta(\varepsilon, u_i \cdot a_i)$$

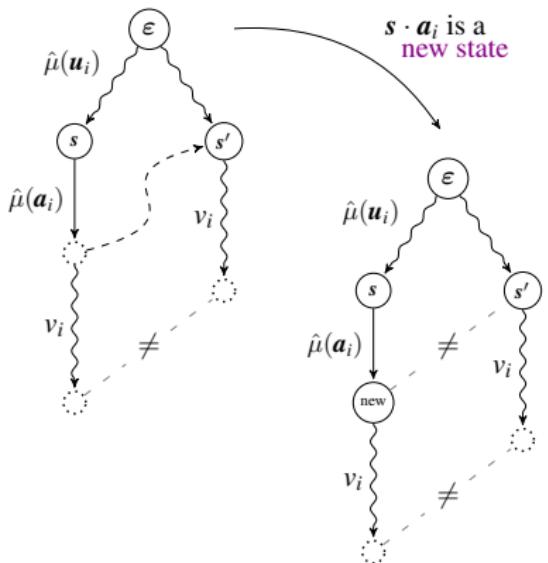


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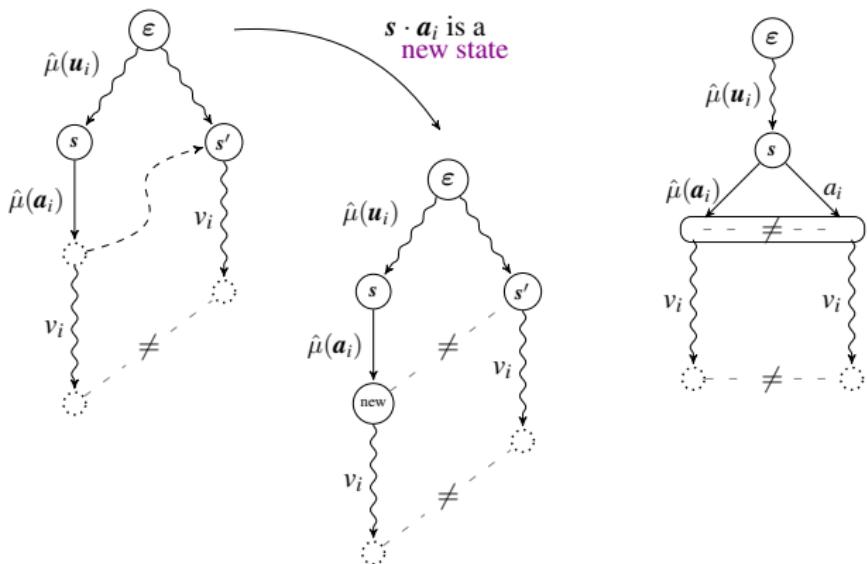


Counter-example Treatment (Symbolic Breakpoint)

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$$f(\hat{\mu}(s_{i-1} \cdot a_i) \cdot v_i) \neq f(\hat{\mu}(s_i) \cdot v_i) \quad f(\hat{\mu}(s_{i-1}) \cdot a_i \cdot v_i) \neq f(\hat{\mu}(s_{i-1}) \cdot \hat{\mu}(a_i) \cdot v_i)$$

$$s_i = \delta(\epsilon, u_i \cdot a_i)$$

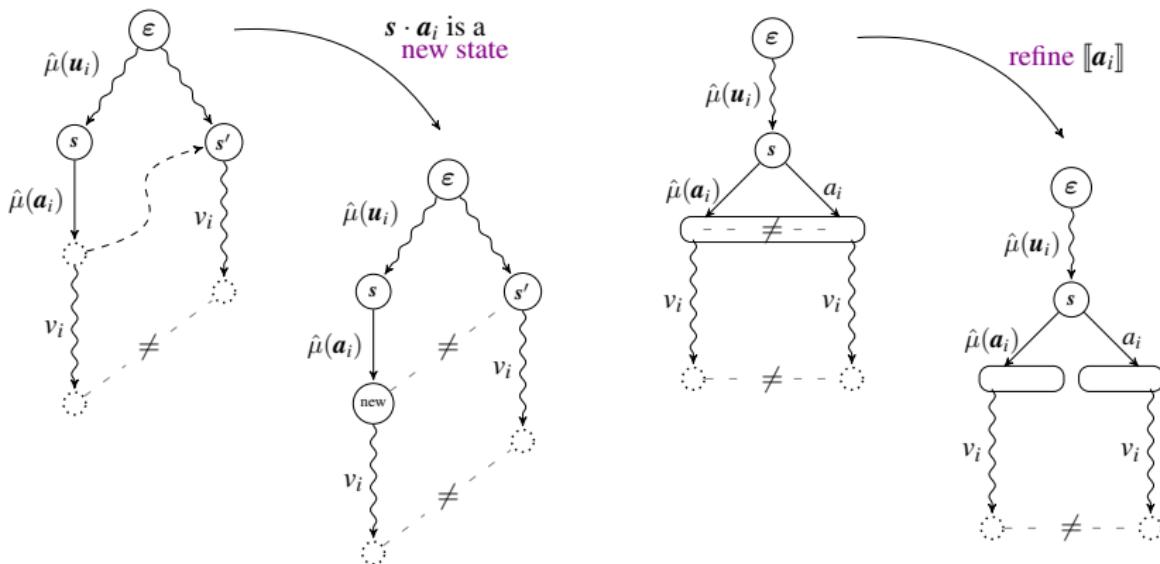


Counter-example Treatment (Symbolic Breakpoint)

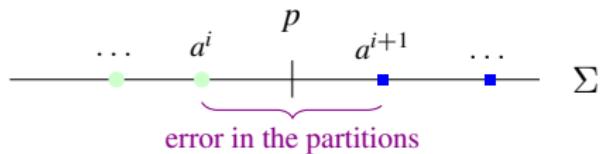
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$$f(\hat{\mu}(s_{i-1} \cdot a_i) \cdot v_i) \neq f(\hat{\mu}(s_i) \cdot v_i) \quad f(\hat{\mu}(s_{i-1}) \cdot a_i \cdot v_i) \neq f(\hat{\mu}(s_{i-1}) \cdot \hat{\mu}(a_i) \cdot v_i)$$

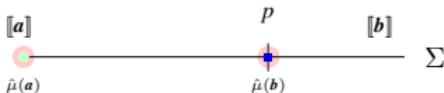
$$s_i = \delta(\epsilon, u_i \cdot a_i)$$



Learning with a Teacher



- Equivalence is checked by an oracle (teacher) returning a minimal counter-examples (in length and lexicographically)
- Choose as evidence the min element of the interval (Σ has min)
- The counter-example indicates the minimal element of a new transition (in horizontal expansion)
- The partition bounds are exact and **no error** is introduced



Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

observation table

semantics

hypothesis automaton

	ϵ
ϵ	

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

observation table

semantics

hypothesis automaton

	ε
ε	—



Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

observation table

semantics

hypothesis automaton

	ε
ε	—
a_0	a_0

ε
 $\llbracket a_0 \rrbracket = [1, 100)$

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

observation table

semantics

hypothesis automaton

	ε
ε	—
a_0	+

ε
 $\llbracket a_0 \rrbracket = [1, 100)$

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

observation table

semantics

hypothesis automaton

	ε
ε	—
a_0^1	+

ε
 $\llbracket a_0 \rrbracket = [1, 100)$

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

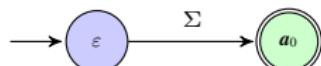
observation table

	ε
ε	-
a_0	+

semantics

ε
 $\llbracket a_0 \rrbracket = [1, 100)$
 a_0

hypothesis automaton



Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

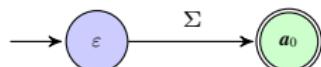
observation table

	ε
ε	—
a_0	+
$a_0 a_1$	$\begin{matrix} 1 & 1 \end{matrix}$

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 100)$
- a_0
- $\llbracket a_1 \rrbracket = [1, 100)$

hypothesis automaton



Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

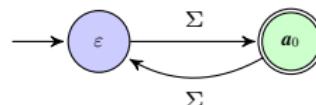
observation table

	ε
ε	—
a_0	+
a_0a_1	—

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 100)$
- a_0
- $\llbracket a_1 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

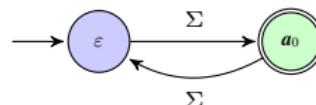
observation table

	ε
ε	—
a_0	+
a_0a_1	—

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 100)$
- a_0
- $\llbracket a_1 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $- 24$

$24 \in \llbracket a_0 \rrbracket$ but $1 \not\sim 24$

→ refine a_0

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

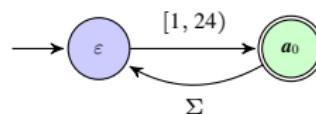
observation table

	ε
ε	—
a_0	+
a_0a_1	—
a_2	

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 24)$
- $\llbracket a_2 \rrbracket = [24, 100)$
- a_0
- $\llbracket a_1 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $- 24$

$24 \in \llbracket a_0 \rrbracket$ but $1 \not\sim 24$

→ refine a_0

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

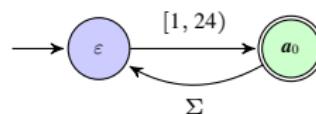
observation table

	ε
ε	—
a_0	+
a_0a_1	—
a_2	—

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 24)$
- $\llbracket a_2 \rrbracket = [24, 100)$
- a_0
- $\llbracket a_1 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $- 24$

$24 \in \llbracket a_0 \rrbracket$ but $1 \not\sim 24$

→ refine a_0

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

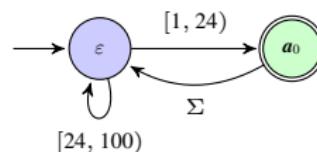
observation table

	ε
ε	—
a_0	+
$a_0 a_1$	—
a_2	—

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 24)$
- $\llbracket a_2 \rrbracket = [24, 100)$
- a_0

hypothesis automaton



Ask Equivalence Query:

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

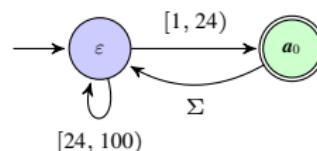
observation table

	ε
ε	—
a_0	+
a_0a_1	—
a_2	—

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 24)$
- $\llbracket a_2 \rrbracket = [24, 100)$
- a_0

hypothesis automaton



$$\llbracket a_1 \rrbracket = [1, 100)$$

Ask Equivalence Query:
counterexample $+ 1 \cdot 66$

$66 \in \llbracket a_1 \rrbracket$ but $1 \not\sim 66$

→ refine a_1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

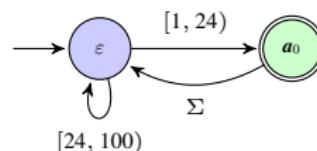
observation table

	ε
ε	—
a_0	+
a_0a_1	—
a_2	—

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 24)$
- $\llbracket a_2 \rrbracket = [24, 100)$
- a_0

hypothesis automaton



Ask Equivalence Query:
counterexample $+ 1 \cdot 66$

$66 \in \llbracket a_1 \rrbracket$ but $1 \not\sim 66$

→ refine a_1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

observation table

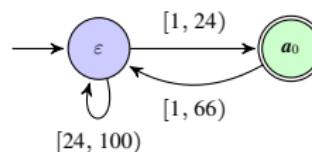
	ε
ε	—
a_0	+
a_0a_1	—
a_2	—
a_0a_3	—

semantics

ε
 $\llbracket a_0 \rrbracket = [1, 24)$
 $\llbracket a_2 \rrbracket = [24, 100)$
 a_0

$\llbracket a_1 \rrbracket = [1, 66)$
 $\llbracket a_3 \rrbracket = [66, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+1 \cdot 66$

$66 \in \llbracket a_1 \rrbracket$ but $1 \not\sim 66$

→ refine a_1

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

observation table

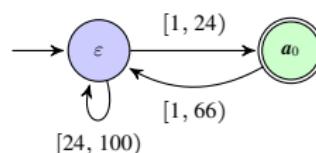
	ε
ε	—
a_0	+
a_0a_1	—
a_2	—
a_0a_3	+

semantics

ε
 $\llbracket a_0 \rrbracket = [1, 24)$
 $\llbracket a_2 \rrbracket = [24, 100)$
 a_0

$\llbracket a_1 \rrbracket = [1, 66)$
 $\llbracket a_3 \rrbracket = [66, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+1 \cdot 66$

$66 \in \llbracket a_1 \rrbracket$ but $1 \not\sim 66$

→ refine a_1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

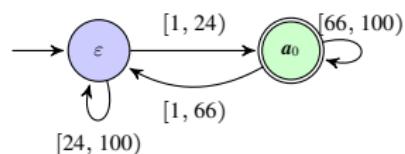
observation table

	ε
ε	—
a_0	+
a_0a_1	—
a_2	—
a_0a_3	+

semantics

$$\begin{aligned} \varepsilon &= [1, 24) \\ [a_0] &= [1, 24) \\ [a_2] &= [24, 100) \\ a_0 & \end{aligned}$$

hypothesis automaton



Ask Equivalence Query:

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

observation table

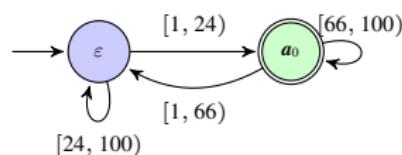
	ε
ε	—
a_0	+
a_0a_1	—
a_2	—
a_0a_3	+

semantics

$$\begin{aligned} \varepsilon &= [1, 24) \\ [a_0] &= [1, 24) \\ [a_2] &= [24, 100) \\ a_0 & \end{aligned}$$

$$\begin{aligned} [a_1] &= [1, 66) \\ [a_3] &= [66, 100) \end{aligned}$$

hypothesis automaton



Ask Equivalence Query:
counterexample $-24 \cdot 1$

$1 \not\sim 24 \cdot 1 \longrightarrow$ add

distinguishing suffix 1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

observation table

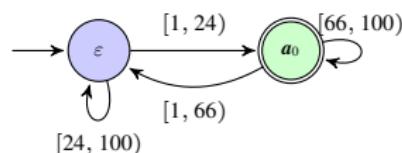
	ε	1
ε	—	
a_0	+	
$a_0 a_1$	—	
a_2	—	
$a_0 a_3$	+	

semantics

$$\begin{aligned} \varepsilon &= [1, 24) \\ [a_0] &= [1, 24) \\ [a_2] &= [24, 100) \\ a_0 & \end{aligned}$$

$$\begin{aligned} [a_1] &= [1, 66) \\ [a_3] &= [66, 100) \end{aligned}$$

hypothesis automaton



Ask Equivalence Query:
counterexample $-24 \cdot 1$

$1 \not\sim 24 \cdot 1 \longrightarrow$ add

distinguishing suffix 1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

observation table

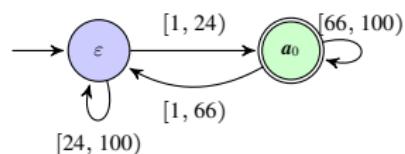
	ε	1
ε	—	+
a_0	+	—
a_0a_1	—	+
a_2	—	—
a_0a_3	+	—

semantics

$$\begin{aligned} \varepsilon &= [1, 24) \\ [a_0] &= [1, 24) \\ [a_2] &= [24, 100) \\ a_0 & \end{aligned}$$

$$\begin{aligned} [a_1] &= [1, 66) \\ [a_3] &= [66, 100) \end{aligned}$$

hypothesis automaton



Ask Equivalence Query:
counterexample $-24 \cdot 1$

$1 \not\sim 24 \cdot 1 \longrightarrow$ add

distinguishing suffix 1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

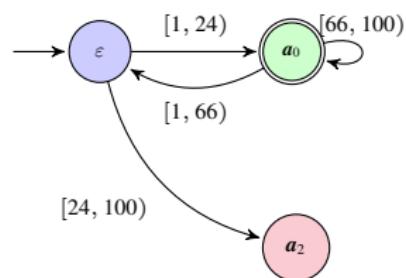
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—

semantics

- ε
- $\llbracket a_0 \rrbracket = [1, 24)$
- $\llbracket a_2 \rrbracket = [24, 100)$
- a_0
- $\llbracket a_1 \rrbracket = [1, 66)$
- $\llbracket a_3 \rrbracket = [66, 100)$
- a_2

hypothesis automaton



Ask Equivalence Query:
counterexample $-24 \cdot 1$

$1 \not\sim 24 \cdot 1 \longrightarrow$ add
distinguishing suffix 1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

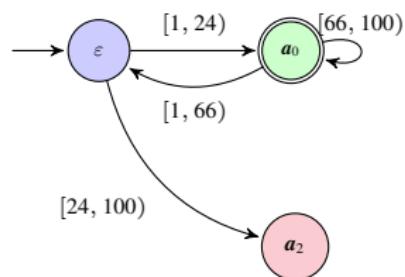
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
$a_0 a_1$	—	+
$a_0 a_3$	+	—
$a_2 a_4$	—	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $-24 \cdot 1$

$1 \not\sim 24 \cdot 1 \longrightarrow$ add

distinguishing suffix 1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

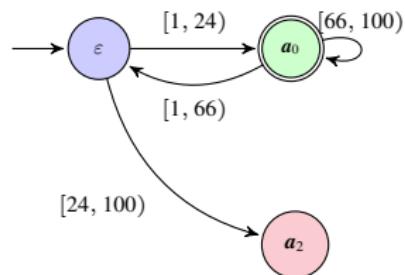
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
$a_0 a_1$	—	+
$a_0 a_3$	+	—
$a_2 a_4$	—	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $-24 \cdot 1$

$1 \not\sim 24 \cdot 1 \longrightarrow$ add

distinguishing suffix 1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

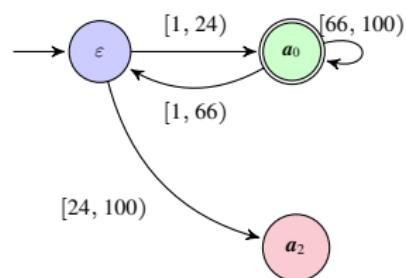
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $-24 \cdot 1$

$1 \not\sim 24 \cdot 1 \longrightarrow$ add

distinguishing suffix 1

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

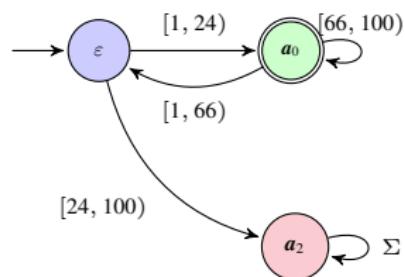
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

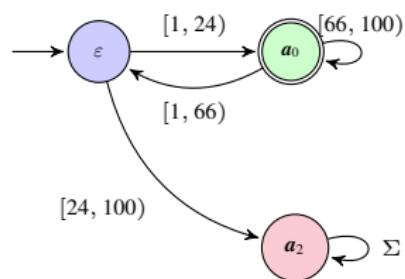
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
$a_0 a_1$	—	+
$a_0 a_3$	+	—
$a_2 a_4$	—	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+24 \cdot 51$

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

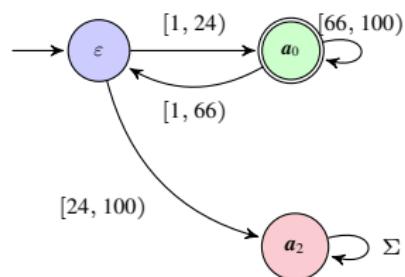
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+24 \cdot 51$

$51 \in \llbracket a_4 \rrbracket$ but $1 \not\sim 51$

→ refine a_4

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

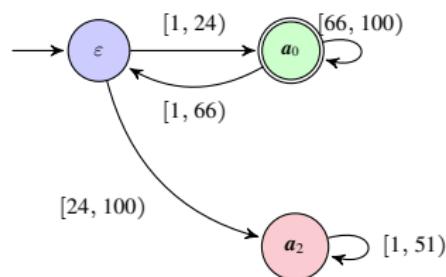
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—
a_2a_5	—	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 51)$, $\llbracket a_5 \rrbracket = [51, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+24 \cdot 51$

$51 \in \llbracket a_4 \rrbracket$ but $1 \not\sim 51$

→ refine a_4

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

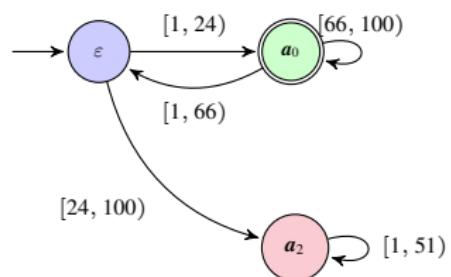
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—
a_2a_5	+	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 51)$, $\llbracket a_5 \rrbracket = [51, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+24 \cdot 51$

$51 \in \llbracket a_4 \rrbracket$ but $1 \not\sim 51$
 → refine a_4

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

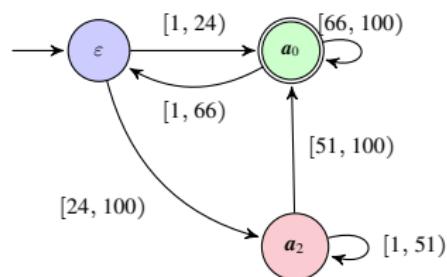
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
$a_0 a_1$	—	+
$a_0 a_3$	+	—
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 51)$, $\llbracket a_5 \rrbracket = [51, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+24 \cdot 51$

$51 \in \llbracket a_4 \rrbracket$ but $1 \not\sim 51$
 → refine a_4

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

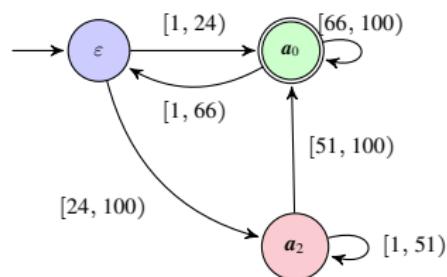
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—
a_2a_5	+	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 51)$, $\llbracket a_5 \rrbracket = [51, 100)$

hypothesis automaton



Ask Equivalence Query:
counterexample $+24 \cdot 51$

$51 \in \llbracket a_4 \rrbracket$ but $1 \not\sim 51$
 → refine a_4

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

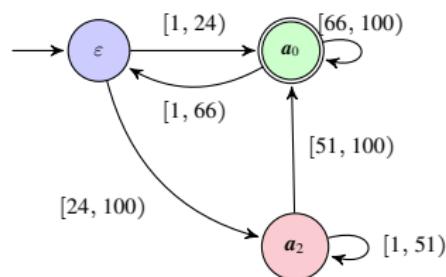
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—
a_2a_5	+	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 51)$, $\llbracket a_5 \rrbracket = [51, 100)$

hypothesis automaton



Ask Equivalence Query:
True

Example with Teacher ($\Sigma = [1, 100)$)

Teacher returns minimal counterexamples

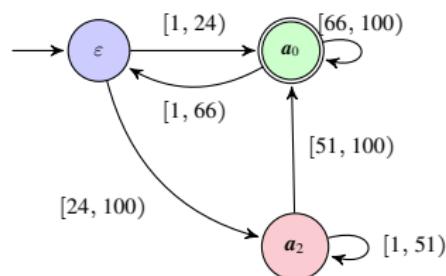
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—
a_2a_5	+	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 51)$, $\llbracket a_5 \rrbracket = [51, 100)$

hypothesis automaton



Ask Equivalence Query:
True

$$M = \{\varepsilon, 1, 24, 11, 166, 241, 2451, 111, 1661, 2411, 24511\}$$

$$|M| = 11, |MQ| = 7, |EQ| = 5, |S| = 3, |R| = 4$$

Example with Teacher ($\Sigma = [1, 100]$)

Teacher returns minimal counterexamples

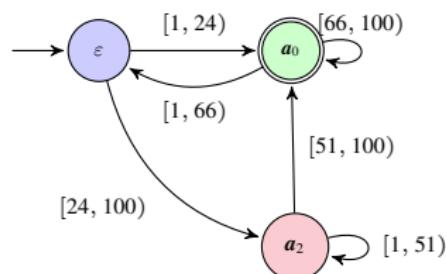
observation table

	ε	1
ε	—	+
a_0	+	—
a_2	—	—
a_0a_1	—	+
a_0a_3	+	—
a_2a_4	—	—
a_2a_5	+	—

semantics

- ε : $\llbracket a_0 \rrbracket = [1, 24)$, $\llbracket a_2 \rrbracket = [24, 100)$
- a_0 : $\llbracket a_1 \rrbracket = [1, 66)$, $\llbracket a_3 \rrbracket = [66, 100)$
- a_2 : $\llbracket a_4 \rrbracket = [1, 51)$, $\llbracket a_5 \rrbracket = [51, 100)$

hypothesis automaton



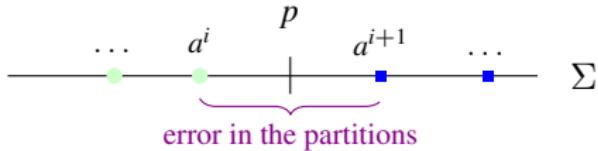
Ask Equivalence Query:
True

$$M = \{\varepsilon, 1, 24, 11, 166, 241, 2451, 111, 1661, 2411, 24511\}$$

L^* over $(\Sigma \cap \mathbb{N}) \rightarrow |M| = 790$, $|MQ| = 789$, $|EQ| = 2$, $|S| = 4$, $|R| = 396$

$$|M| = 11, |MQ| = 7, |EQ| = 5, |S| = 3, |R| = 4$$

Learning without a Teacher



- Equivalence is checked by testing random words selected using a probability distribution D
- Counter-examples are not minimal we may have errors in the boundaries
- Counter-examples may be missed terminate algorithm and return hypothesis after $r(\varepsilon, \delta, i)$ random words have been tested, none of which is a counter-example
- The final hypothesis \mathcal{A} is a good approximation of the target language L with high probability

$$P(d(L, L_{\mathcal{A}}) < \varepsilon) \geq 1 - \delta \quad (\text{PAC learnability})$$

Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

observation table

semantics

hypothesis automaton

Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

observation table

semantics

hypothesis automaton

	ε
ε	

Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

observation table

	ε
ε	—

semantics



hypothesis automaton

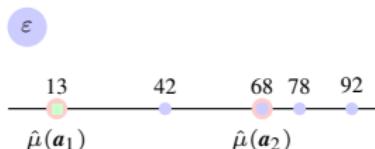
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

observation table

	ε
ε	—

semantics



hypothesis automaton

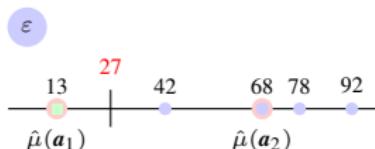
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

observation table

	ε
ε	—

semantics



hypothesis automaton

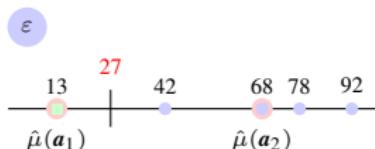
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

observation table

	ε
ε	—
a_1	+
a_2	—

semantics



hypothesis automaton

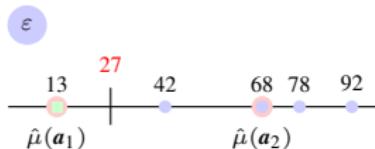
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Counterexamples are not minimal

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	ε
ε	—
a_1	+
a_2	—

semantics



hypothesis automaton

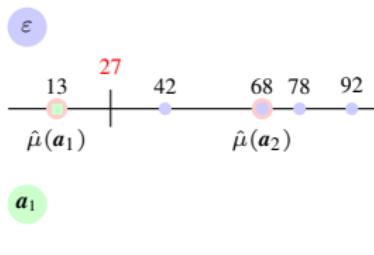
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Counterexamples are not minimal

observation table

	ε
ε	—
a_1	+
a_2	—

semantics



hypothesis automaton

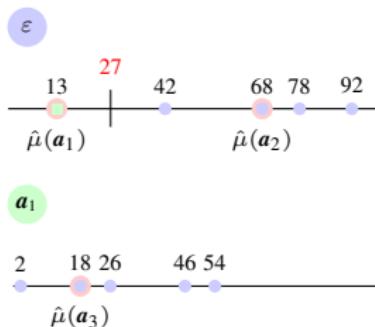
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

observation table

	ε
ε	—
a_1	+
a_2	—

semantics



hypothesis automaton

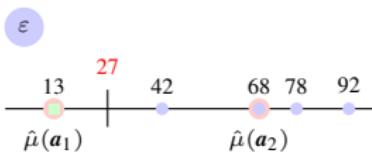
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

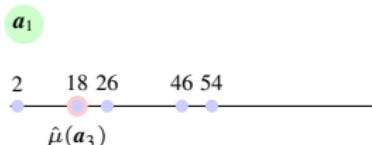
observation table

	ε
ε	—
a_1	+
a_2	—
$a_1 a_3$	—

semantics



hypothesis automaton



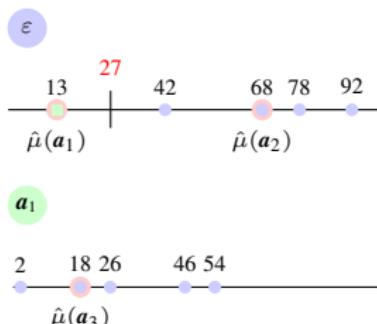
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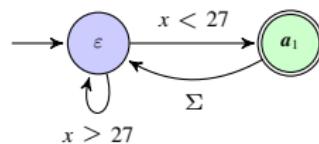
observation table

	ε
ε	—
a_1	+
a_2	—
$a_1 a_3$	—

semantics



hypothesis automaton



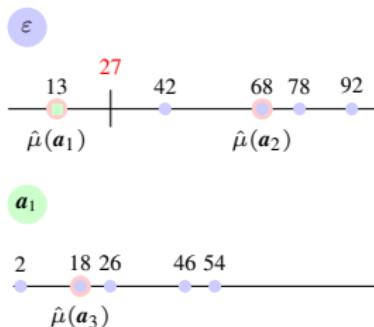
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Counterexamples are not minimal

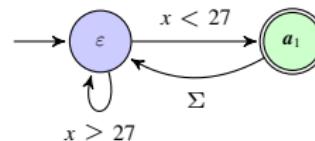
observation table

	ε
ε	—
a_1	+
a_2	—
$a_1 a_3$	—

semantics



hypothesis automaton



Ask Equivalence Query:
counterexample — 12 · 73 · 11

add distinguishing string 11

→ new state

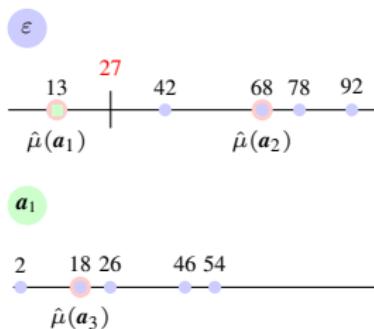
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Counterexamples are not minimal

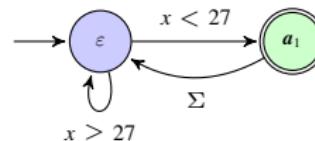
observation table

	ε	11
ε	—	
a_1	+	
a_2	—	
$a_1 a_3$	—	

semantics



hypothesis automaton



Ask Equivalence Query:
counterexample — 12 · 73 · 11

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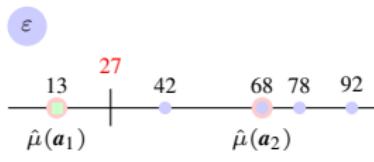
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Counterexamples are not minimal

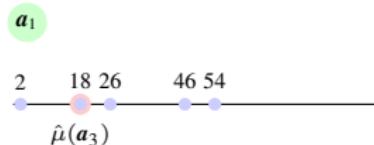
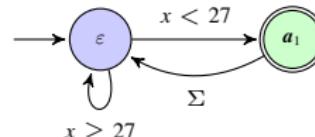
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+

semantics



hypothesis automaton



Ask Equivalence Query:
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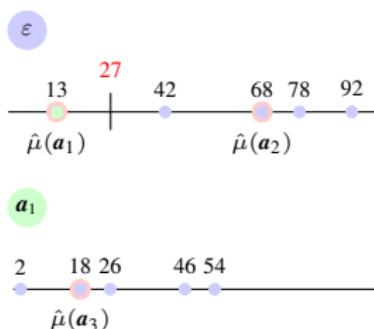
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

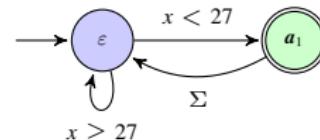
observation table

	ε	11
ε	—	+
13	+	—
a_1	—	—
68	—	—
a_2	—	—
13 18 $a_1 a_3$	—	+

semantics



hypothesis automaton



Ask Equivalence Query:
counterexample — 12 · 73 · 11

add distinguishing string 11

→ new state

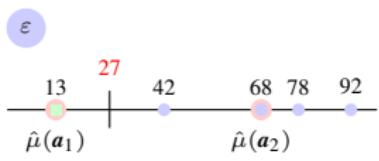
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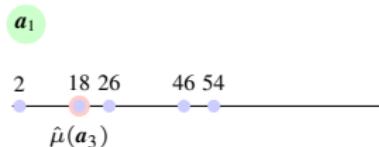
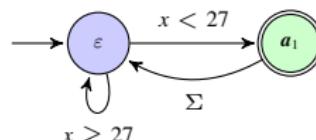
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+

semantics



hypothesis automaton



a_2



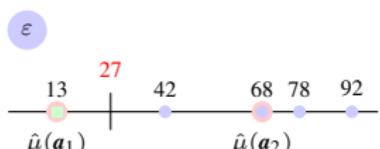
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

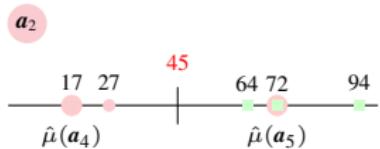
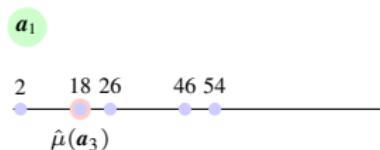
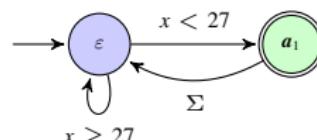
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+

semantics



hypothesis automaton



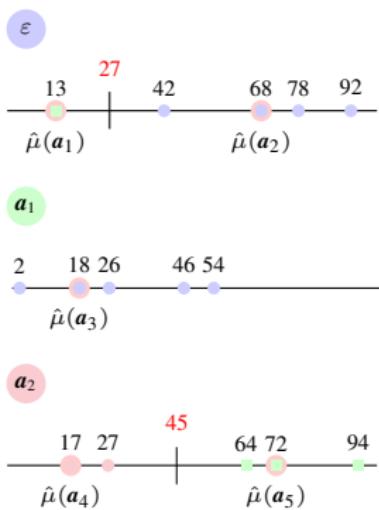
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

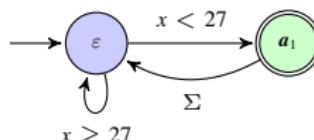
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_2 a_4$	$13\ 18$	$68\ 17$
$a_2 a_5$	$68\ 72$	

semantics



hypothesis automaton



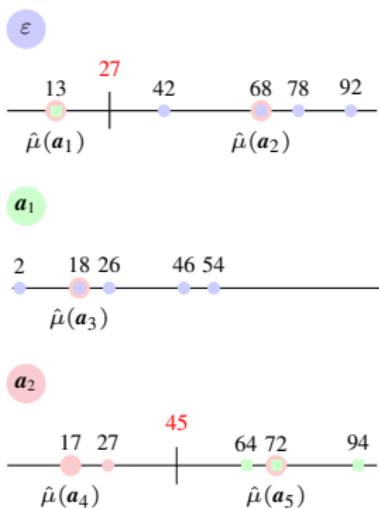
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

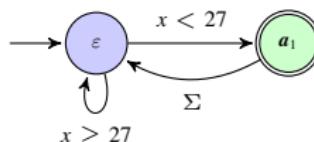
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



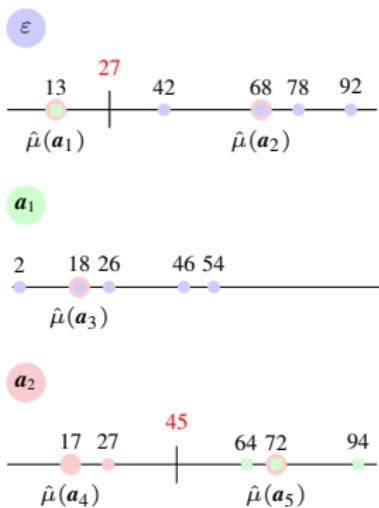
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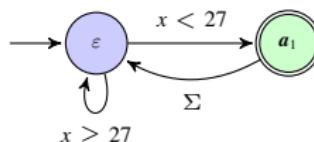
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



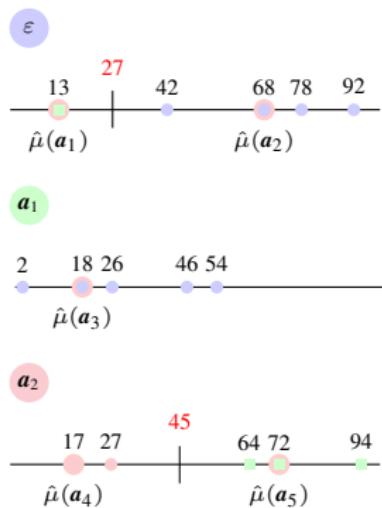
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

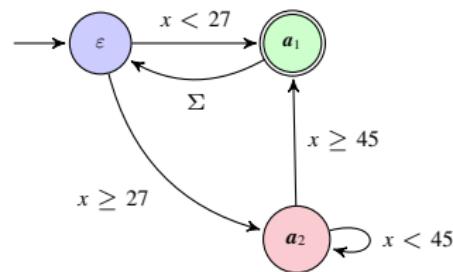
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



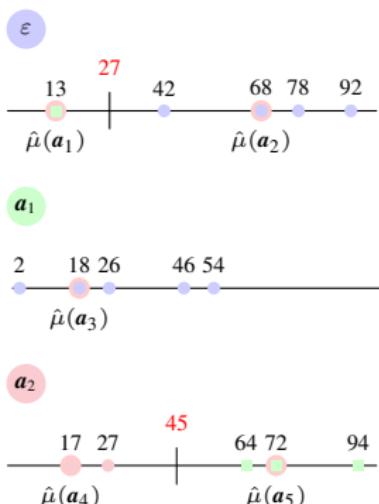
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Counterexamples are not minimal

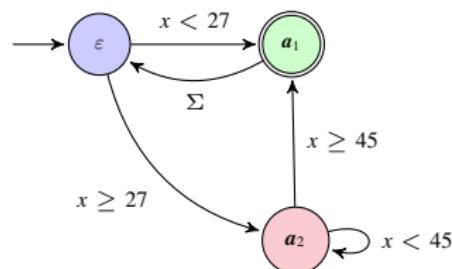
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



Ask Equivalence Query:
counterexample $-12 \cdot 73 \cdot 11$

add 73 as evidence of a_1

→ new transition

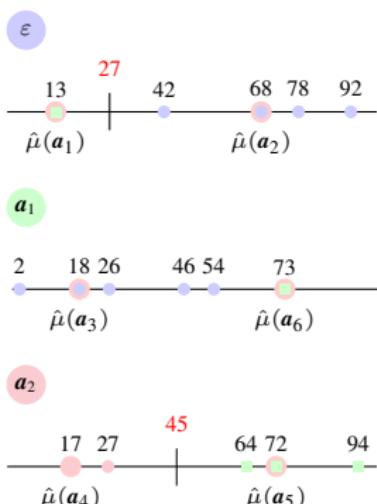
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

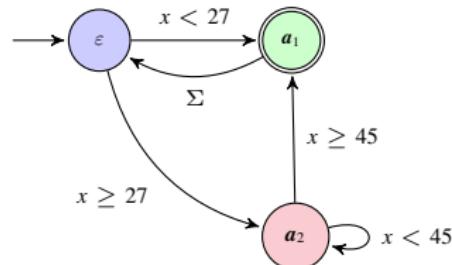
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



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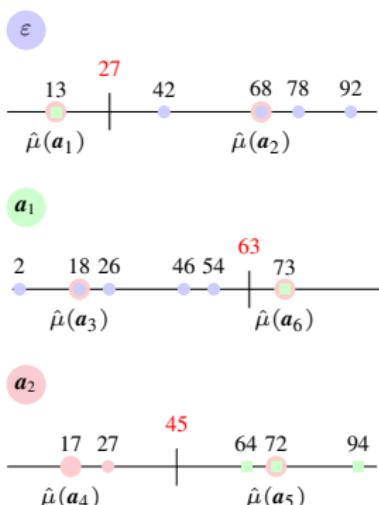
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Counterexamples are not minimal

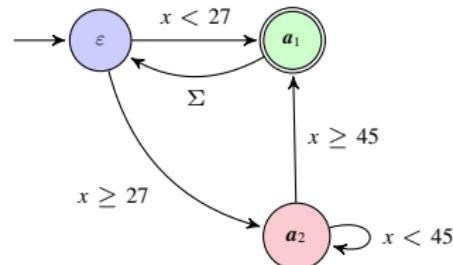
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



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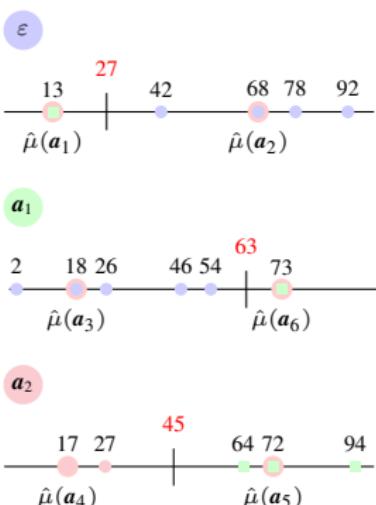
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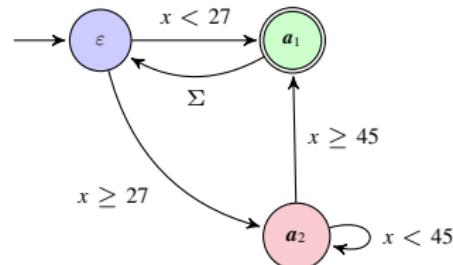
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_1 a_6$	+	—
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



Ask Equivalence Query:
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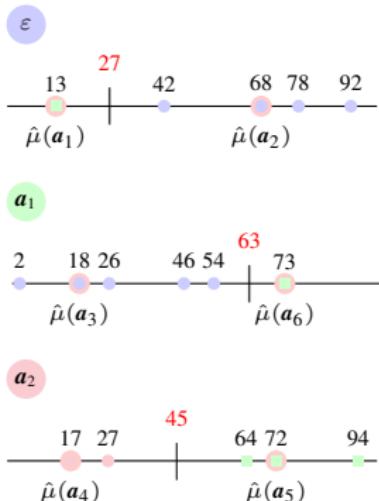
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Counterexamples are not minimal

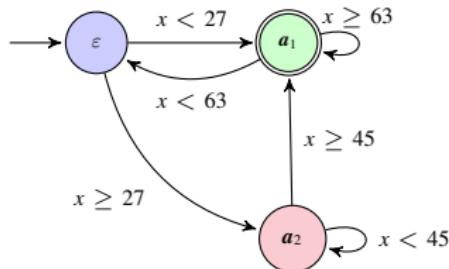
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_1 a_6$	+	—
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



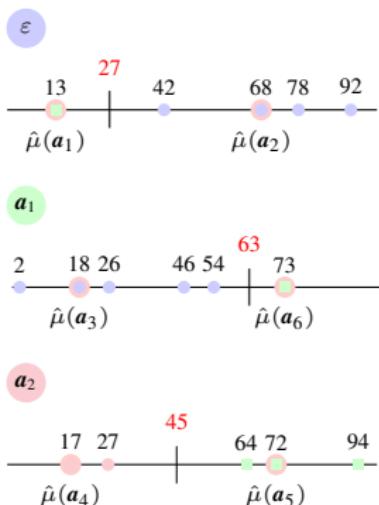
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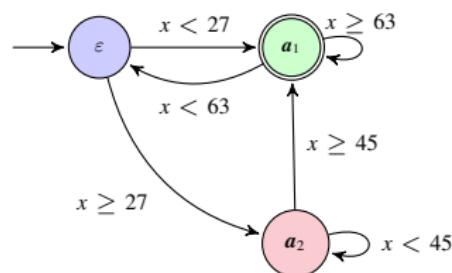
observation table

	ε	11
ε	—	+
13	+	—
a_1	+	—
68	—	—
a_2	—	—
13 18	—	+
$a_1 a_3$	—	+
13 73	+	—
$a_1 a_6$	+	—
68 17	—	—
$a_2 a_4$	—	—
68 72	+	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



Ask Equivalence Query:
counterexample $-52 \cdot 47$

add 47 as evidence of a_2

→ refine existing transition

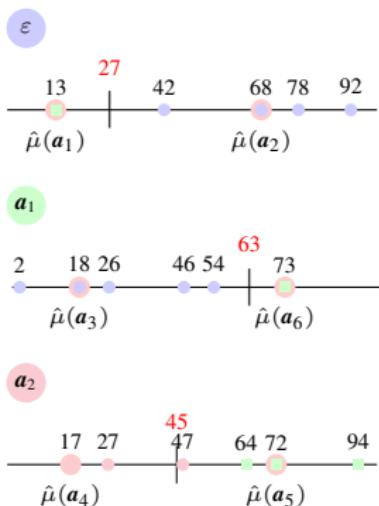
Example without Teacher ($\Sigma = [1, 100]$)

Counterexamples are not minimal

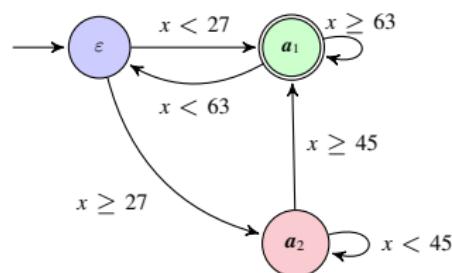
observation table

	ε	11
ε	—	+
13	+	—
a_1	+	—
68	—	—
a_2	—	—
13 18	—	+
$a_1 a_3$	—	+
13 73	+	—
$a_1 a_6$	+	—
68 17	—	—
$a_2 a_4$	—	—
68 72	+	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



Ask Equivalence Query:
counterexample — 52 · 47

add 47 as evidence of a_2

→ refine existing transition

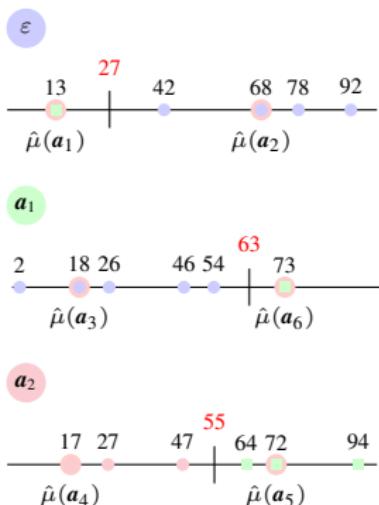
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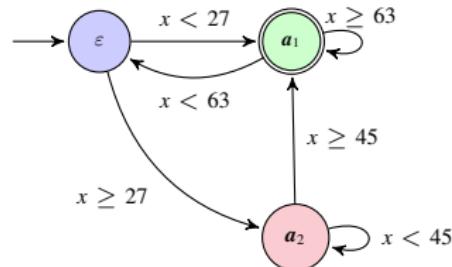
observation table

	ε	11
ε	—	+
13	+	—
a_1	+	—
68	—	—
a_2	—	—
13 18	—	+
$a_1 a_3$	—	+
13 73	+	—
$a_1 a_6$	+	—
68 17	—	—
$a_2 a_4$	—	—
68 72	+	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



Ask Equivalence Query:
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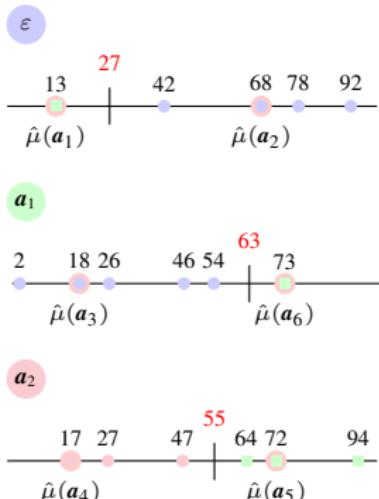
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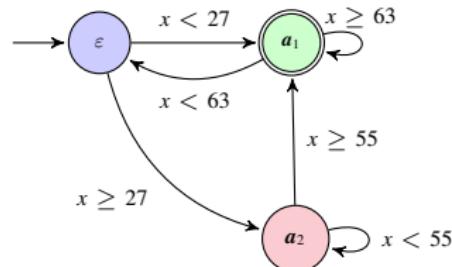
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_1 a_6$	+	—
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



hypothesis automaton



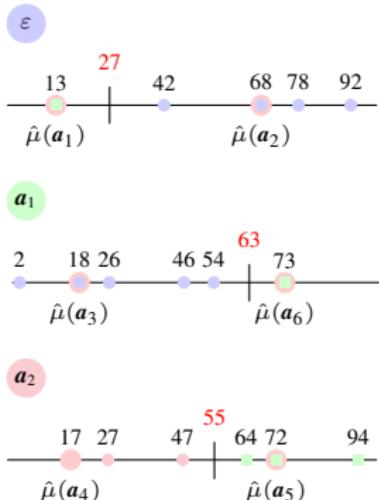
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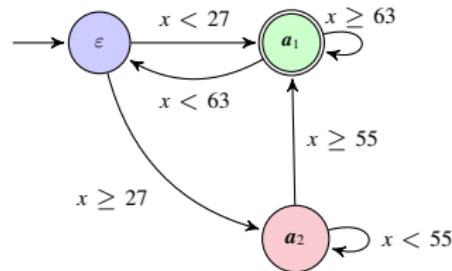
observation table

	ε	11
ε	—	+
a_1	+	—
a_2	—	—
$a_1 a_3$	—	+
$a_1 a_6$	+	—
$a_2 a_4$	—	—
$a_2 a_5$	+	—

semantics



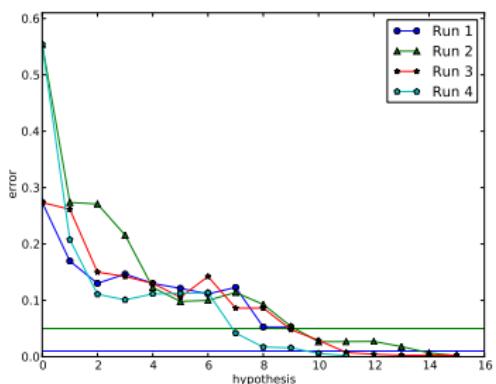
hypothesis automaton



Ask Equivalence Query:

...

Computing the Error



- The error of the hypotheses along several runs of the algorithm

The error is measured as volumes of the symmetric difference \mathcal{L} between the conjectured and the target language

$$\text{error} = D(\mathcal{L}) = \lim_{k \rightarrow \infty} D_k(\mathcal{L}),$$

where $D_k(\mathcal{L})$ is the k -volume of \mathcal{L} , i.e., $D_k(\mathcal{L}) = V(\mathcal{L}_k)/V(\Sigma_k)$

Current Status

- We implemented the algorithm for the case $\Sigma \subset \mathbb{R}$ and $\Sigma \subset \mathbb{N}$ with and without a teacher
- Experimental results on password rules over ASCII characters
- We developed a similar algorithm for $\Sigma = \mathbb{B}^n$ for large n
- We use bounded complexity alphabet partitions in the style of k -DNF or decision lists
- First results are encouraging, will be used to extend to $\Sigma \subset \mathbb{R}^n$

Discussion

- Symbolic automata: concrete values are used to select a transition but their exact value is not remembered by the automaton
- Is there a niche for it in the time series analysis science?
- Combination of temporal (automaton) and static concept learning
- Can be an alternative to (deep) recurrent neural networks
- Should relax full compatibility with the sample (noise) and be ready to drop negative examples

Discussion

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Thank you