Glueability of resource proof-structures: inverting the Taylor expansion

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Abstract

A Multiplicative-Exponential Linear Logic (MELL) proof-structure can be expanded into a set of resource proof-structures: its Taylor expansion. We introduce a new criterion characterizing those sets of resource proof-structures that are part of the Taylor expansion of some MELL proof-structure, through a rewriting system acting both on resource and MELL proof-structures.

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1 Introduction

Resource λ-calculus and the Taylor expansion Girard's linear logic (LL, [15]) is a refinement of intuitionistic and classical logic that isolates the infinitary parts of reasoning in two (dual) modalities: the *exponentials* ! and ?. They give a logical status to the operations of memory management such as *copying* and *erasing*: a linear proof corresponds—via Curry–Howard isomorphism—to a program that uses its argument *linearly*, *i.e.* exactly once, while an exponential proof corresponds to a program that can use its argument at will.

The intuition that linear programs are analogous to linear functions (as studied in linear algebra) while exponential programs mirror a more general class of analytic functions got a technical incarnation in Ehrhard's work [9, 10] on LL-based denotational semantics for the λ-calculus. This investigation has been then internalized in the syntax, yielding the *resource λ-calculus* [5, 11, 14]: there, copying and erasing are forbidden and replaced by the possibility to apply a function to a *bag* of resource λ-terms which specifies how many times an argument can be linearly passed to the function, so as to represent only bounded computations.

The *Taylor expansion* associates with an ordinary λ-term a (generally infinite) set of resource λ-terms, recursively approximating the usual application: the Taylor expansion of the λ-term MN is made of resource λ-terms of the form $t[u_1, \dots, u_n]$, where t is a resource λ-term in the Taylor expansion of M, and $[u_1, \dots, u_n]$ is a bag of arbitrarily finitely many (possibly 0) resource λ-terms in the Taylor expansion of N. Roughly, the idea is to decompose a program into a set of purely "resource-sensitive programs", all of them containing only bounded (although possibly non-linear) calls to inputs. The notion of Taylor expansion has many applications in the theory of the λ-calculus, *e.g.* in the study of linear head reduction [12], normalization [24, 27], Böhm trees [4, 19], λ-theories [20], intersection types [22]. More

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generally, understanding the relation between a program and its Taylor expansion renews the logical approach to the quantitative analysis of computation started with the inception of LL.

A natural question is the *inverse Taylor expansion problem*: how to characterize which sets of resource λ-terms are contained in the Taylor expansion of a same λ-term? Ehrhard and Regnier [14] defined a simple *coherence* relation such that a finite set of resource λ-terms is included in the Taylor expansion of a λ-term if and only if the elements of this set are pairwise coherent. Coherence is crucial in many structural properties of the resource λ-calculus, such as in the proof that in the λ-calculus normalization and Taylor expansion commute [12, 14].

We aim to solve the inverse Taylor expansion problem in the more general context of LL, more precisely in the *multiplicative-exponential fragment* MELL of LL, being aware that for MELL no coherence relation can solve the problem (see below).

Proof-nets, proof-structures and their Taylor expansion: seeing trees behind graphs In MELL, linearity and the sharp analysis of computations naturally lead to represent proofs in a more general *graph*-like syntax instead of a term-like or tree-like one.[1] Indeed, linear negation is involutive and classical duality can be interpreted as the possibility of juggling between different conclusions, without a distinguished output. Graphs representing proofs in MELL are called *proof-nets*: their syntax is richer and more expressive than the λ-calculus. Contrary to λ-terms, proof-nets are special inhabitants of the wider land of *proof-structures*: they can be characterized, among proof-structures, by abstract (geometric) conditions called correctness criteria [15]. The procedure of cut-elimination can be applied to proof-structures, and proof-nets can also be seen as the proof-structures with a good behavior with respect to cut-elimination [1]. Proof-structures can be interpreted in denotational models and proof-nets can be characterized among them by semantic means [25]. It is then natural to attack problems in the general framework of proof-structures. In this work, correctness plays no role at all, hence we will consider proof-structures and not only proof-nets. MELL proof-structures are a particular kind of graphs, whose edges are labeled by MELL formulæ and vertices by MELL connectives, and for which special subgraphs are highlighted, the *boxes*, representing the parts of the proof-structure that can be copied and discarded (*i.e.* called an unbounded number of times). A box is delimited from the rest of a proof-structure by exponential modalities: its border is made of one !-cell, its principal door, and arbitrarily many ?-cells, its auxiliary doors. Boxes are nested or disjoint (they cannot partially overlap), so as to add a tree-like structure to proof-structures *aside* from their graph-like nature.

As in λ-calculus, one can define [13] box-free *resource proof-structures*[2], where !-cells make resources available boundedly, and the *Taylor expansion* of MELL proof-structures into these resource proof-structures, that recursively copies the content of the boxes an arbitrary number of times. In fact, as somehow anticipated by Boudes [3], such a Taylor expansion operation can be carried on any tree-like structure. This primitive, abstract, notion of Taylor expansion can then be pulled back to the structure of interest, as shown in [18] and put forth again here.

The question of coherence for proof-structures The inverse Taylor expansion problem has a natural counterpart in the world of MELL proof-structures: given a set of resource proof-structures, is there a MELL proof-structure the expansion of which contains the set? Pagani and Tasson [23] give the following answer: it is possible to decide whether a finite set of resource proof-structures is a subset of the Taylor expansion of a same MELL proof-structure

[1] A term-like object is essentially a tree, with one output (its root) and many inputs (its other leaves).

[2] Also known as differential proof-structures [6] or differential nets [13, 21, 7] or simple nets [23].

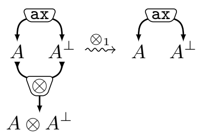
(and even possible to do it in non-deterministic polynomial time); but unlike the λ-calculus, the structure of the relation "being part of the Taylor expansion of a same proof-structure" is *much more* complicated than a binary (or even n-ary) coherence. Indeed, for any $n > 1$, it is possible to find $n + 1$ resource proof-structures such that any n of them are in the Taylor expansion of some MELL proof-structure, but there is no MELL proof-structure whose Taylor expansion has all the $n+1$ as elements (see our Example 21 and [26, pp. 244-246]).

In this work, we introduce a new combinatorial criterion, *glueability*, for deciding whether a set of resource proof-structures is a subset of the Taylor expansion of some MELL proof structure, based on a rewriting system on sequences of MELL formulæ. Our criterion is more general (and, we believe, simpler) than the one of [23], which is limited to the *cut-free* case with *atomic axioms* and characterizes only *finite* sets: we do not have these limitations. We believe that our criterion is a useful tool for studying proof-structures. We conjecture that it can be used to show that, for a suitable geometric restriction, a binary coherence relation does exist for resource proof-structures. It might also shed light on correctness and sequentialization.

As the proof-structures we consider are typed, an unrelated difficulty arises: a resource proof-structure might not be in the Taylor expansion of any MELL proof-structure, not because it does not respect the structure imposed by the Taylor expansion, but because its type is impossible.[3] To solve this issue we enrich the MELL proof-structure syntax with a "universal" proof-structure: a special $\maltese$-cell (*daimon*) that can have any number of outputs of any types, and we allow it to appear inside a box, representing information plainly missing (see Section 8 for more details and the way this matter is handled by Pagani and Tasson [23]).

2 Outline and technical issues

The rewritings The essence of our rewriting system is not located on proof-structures but on lists of MELL formulæ (Definition 9). In a very down-to-earth way, this rewriting system is generated by elementary steps akin to rules of sequent calculus read from the *bottom up*: they act on a list of conclusions, analogous to a monolaterous right-handed sequent. These steps are actually more sequentialized than sequent calculus rules, as they do not allow for commutation. For instance, the rule corresponding to the introduction of a $\otimes$ on the i-th formula, is defined as $\otimes_i : (\gamma_1, \dots, \gamma_{i-1}, A \otimes B, \gamma_{i+1}, \dots, \gamma_n) \to (\gamma_1, \dots, \gamma_{i-1}, A, B, \gamma_{i+1}, \dots, \gamma_n)$.



These rewrite steps then act on MELL proof-structures, coherently with their type, by modifying (most of the times, erasing) the cells directly connected to the conclusion of the proof-structure. Formally, this means that there is a functor $\mathbf{qMELL}^{\maltese}$ from the rewrite steps into the category **Rel** of sets and relations, associating with a list of formulæ the set of MELL proof-structures with these conclusions, and with a rewrite step a relation implementing it (Definition 12). The rules *deconstruct* the proof-structure, starting from its conclusions. The rule $\otimes_1$ acts by removing a $\otimes$-cell on the first conclusion, replacing it by two conclusions.

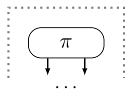
These rules can only act on specific proof-structures, and indeed, capture a lot of their structure: $\otimes_i$ can be applied to a MELL proof-structure R if and only if R has a $\otimes$-cell in the conclusion i (as opposed to, say, an axiom). So, in particular, every proof-structure is completely characterized by any sequence rewriting it to the empty proof-structure.

[3] Similarly, in the λ-calculus, there is no closed λ-term of type $X \to Y$ with $X \neq Y$ atomic, but the resource λ-term $(\lambda f.f)[\,]$ can be given that type: the empty bag $[\,]$ kills any information on the argument.

Naturality The same rules act also on sets of resource proof-structures, defining the functor $\mathfrak{P}\mathbf{qDiLL}_{\mathbf{0}}^{\maltese}$ from the rewrite steps into the category **Rel** (Definition 17). When carefully defined, the Taylor expansion induces a *natural transformation* from $\mathfrak{P}\mathbf{qDiLL}_{\mathbf{0}}^{\maltese}$ to $\mathbf{qMELL}^{\maltese}$ (Theorem 18). By applying this naturality repeatedly, we get our characterization (Theorem 20): a set of resource proof-structures Π is a subset of the Taylor expansion of a MELL proof-structure iff there is a sequence rewriting Π to the singleton of the *empty* proof-structure.

The naturality property is not only a mean to get our characterization, but also an interesting result in itself: natural transformations can often be used to express fundamental properties in a mathematical context. In this case, the *Taylor expansion is natural* with respect to the possibility to build a (MELL or resource) proof-structure by adding a cell to its conclusions or boxing it. Said differently, naturality of the Taylor expansion roughly means that the rewrite rules that deconstruct a MELL proof-structure R and a set of resource proof-structures in the Taylor expansion of R mimic each other.

Quasi-proof-structures and mix Our rewrite rules consume proof-structures from their conclusions. The rule corresponding to boxes in MELL opens a box by deleting its principal door (a !-cell) and its border, while for a resource proof-structure it deletes a !-cell and separates the different copies of the content of the box (possibly) represented by such a !-cell. This operation is problematic in a twofold way. In a resource proof-structure, where the border of boxes is not marked, it is not clear how to identify such copies. On the other side, in a MELL proof-structure the content of a box is not to be treated as if it were at the same level as what is outside of the box: it can be copied many times or erased, while what is outside boxes cannot, and treating the content in the same way as the outside suppresses this distinction, which is crucial in LL. So, we need to remember that the content of a box, even if it is at depth 0 (*i.e.* not contained in any other box) after erasing the box wrapping it by means of our rewrite rules, is not to be mixed with the rest of the structure at depth 0.



In order for our proof-structures to provide this information, we need to generalize them and consider that a proof-structure can have not just a tree of boxes, but a *forest*: this yields the notion of *quasi-proof-structure* (Definition 1). In this way, according to our rewrite rules, opening a box by deleting its principal door amounts to taking a box in the tree and disconnecting it from its root, creating a new tree. We draw this in a quasi-proof-structure by surrounding elements having the same root with a dashed line, open from the bottom, remembering the phantom presence of the border of the box, even if it was erased. This allows one to open the box only when it is “alone”, surrounded by a dashed line (see Definition 11).

This is not merely a technical remark, as this generalization gives a status to the `mix` rule of LL: indeed, mixing two proofs amounts to taking two proofs and considering them as one, without any other modifications. Here, it amounts to taking two proofs, each with its box-tree, and considering them as one by merging the roots of their trees (see the mix step in Definition 11). We embed this design decision up to the level of formulæ, which are segregated in different zones that have to be mixed before interacting (see the notion of partition of a finite sequence of formulæ in Section 3).

Geometric invariance and emptiness: the filled Taylor expansion The use of forests instead of trees for the nesting structure of boxes, where the different roots are thought of as the contents of long-gone boxes, has an interesting consequence in the Taylor expansion: indeed, an element of the Taylor expansion of a proof-structure contains an arbitrary number of copies of the contents of the boxes, in particular *zero*. If we think of the part at depth

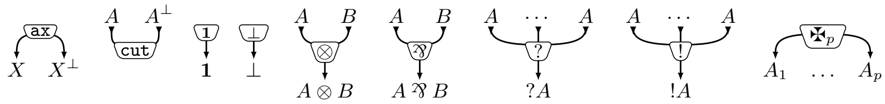


Figure 1 Cells, with their labels and their typed inputs and outputs (ordered from left to right).

0 of a MELL proof-structure as inside an invisible box, its content can be deleted in some elements of the Taylor expansion just as any other box.[4] As erasing completely conclusions would cause the Taylor expansion not preserve the conclusions (which would lead to technical complications), we introduce the *filled Taylor expansion* (Definition 8), which contains not only the elements of the usual Taylor expansion, but also elements of the Taylor expansion where one component has been erased and replaced by a $\maltese$-cell (*daimon*), representing a lack of information, apart from the number and types of the conclusions.

Atomic axioms Our paper first focuses on the case where proof-structures are restricted to *atomic axioms.* In Section 7 we sketch how to adapt our method to the non-atomic case.

3 Proof-structures and the Taylor expansion

MELL **formulæ and (quasi-)proof-structures** Given a countably infinite set of propositional variables $X, Y, Z, \dots$, MELL *formulæ* are defined by the following inductive grammar:

$$A, B ::= X \mid X^\perp \mid \mathbf{1} \mid \perp \mid A \otimes B \mid A ⅋ B \mid\, !A \mid ?A$$

Linear negation is defined via De Morgan laws $\mathbf{1}^\perp = \perp$, $(A \otimes B)^\perp = A^\perp ⅋ B^\perp$ and $(!A)^\perp = ?A$, so as to be involutive, *i.e.* $A^{\perp\perp} = A$. Given a list $\Gamma = (A_1, \dots, A_m)$ of MELL formulæ, a *partition* of Γ is a list $(\Gamma_1, \dots, \Gamma_n)$ of lists of MELL formulæ such that there are $0 = i_0 < \cdots < i_n = m$ with $\Gamma_j = (A_{i_{j-1}+1}, \dots, A_{i_j})$ for all $1 \leqslant j \leqslant n$; such a partition of Γ is also denoted by $(A_1, \dots, A_{i_1}; \cdots ; A_{i_{n-1}+1}, \dots, A_m)$, with lists separated by semi-colons.

We reuse the syntax of proof-structures given in [18] and sketch here its main features. We suppose known definitions of (directed) graph, rooted tree, and morphism of these structures. In what follows we will speak of *tails* in a graph: "hanging" edges with only one vertex. This can be implemented either by adding special vertices or using [2]'s graphs.

If an edge e is incoming in (resp. outgoing from) a vertex v, we say that e is a *input* (resp. *output*) of v. The reflexive-transitive closure of a tree τ is denoted by $\tau^{\circlearrowleft}$: the operator $(\cdot)^{\circlearrowleft}$ lifts to a functor from the category of trees to the category of directed graphs.

▶ **Definition 1.** *A* module M *is a (finite) directed graph with:*
- *vertices v labeled by $\ell(v) \in \{\mathtt{ax}, \mathtt{cut}, \mathbf{1}, \perp, \otimes, ⅋, ?, !\} \cup \{\maltese_p \mid p \in \mathbb{N}\}$, the* type *of v;*
- *edges e labeled by a* MELL *formula $\mathsf{c}(e)$, the* type *of e;*
- *an order $<_M$ that is total on the tails of $|M|$ and on the inputs of each vertex of type ⅋, ⊗.*

Moreover, all the vertices verify the conditions of Figure 1.[5]

A quasi-proof-structure *is a triple $R = (|R|, \mathcal{F}, \mathsf{box})$ where:*
- *$|R|$ is a module with no input tails, called the* module *of R;*
- *$\mathcal{F}$ is a forest of rooted trees with no input tails, called the* box-forest *of R;*

[4] The dual case, of copying the contents of a box, poses no problem in our approach.
[5] Note that there are no conditions on the types of the outputs of vertices of type $\maltese$ (*i.e.* of type $\maltese_p$ for some $p \in \mathbb{N}$); and the outputs of vertices of type $\mathtt{ax}$ must have *atomic* types.

- $\mathsf{box}\colon |R| \to \mathcal{F}^{\circlearrowleft}$ *is a morphism of directed graphs, the* box-function *of R, which induces a partial bijection from the inputs of the vertices of type* ! *and the edges in* $\mathcal{F}$*, and such that:*
  - *for any vertices* v, v' *with an edge from* v' *to* v*, if* $\mathsf{box}(v) \neq \mathsf{box}(v')$ *then* $\ell(v) \in \{!, ?\}$.[6]

  Moreover, for any output tails e_1, e_2, e_3 *in* $|R|$ *which are outputs of the vertices* v_1, v_2, v_3*, respectively, if* $e_1 <_{|R|} e_2 <_{|R|} e_3$ *then it is impossible that* $\mathsf{box}(v_1) = \mathsf{box}(v_3) \neq \mathsf{box}(v_2)$.[7]

A quasi-proof-structure $R = (|R|, \mathcal{F}, \mathsf{box})$ *is:*

1. $\mathsf{MELL}^{\maltese}$ *if all vertices in* $|R|$ *of type* ! *have exactly one input, and the partial bijection induced by* box *from the inputs of the vertices of type* ! *in* $|R|$ *and the edges in* $\mathcal{F}$ *is total.*
2. MELL *if it is* $\mathsf{MELL}^{\maltese}$ *and, for every vertex* v *in* $|R|$ *of type* $\maltese$*, one has* $\mathsf{box}^{-1}(\mathsf{box}(v)) = \{v\}$ *and* $\mathsf{box}(v)$ *is not a root of the box-forest* $\mathcal{F}$ *of* R.
3. $\mathsf{DiLL}_0^{\maltese}$ *if the box-forest* $\mathcal{F}$ *of* R *is just a juxtaposition of roots.*
4. DiLL_0 *(or* resource*) if it is* $\mathsf{DiLL}_0^{\maltese}$ *and there is no vertex in* $|R|$ *of type* $\maltese$.

For the previous systems, a proof-structure *is a quasi-proof-structure whose box-forest is a tree.*

Our MELL proof-structure (*i.e.* a MELL quasi-proof-structure that is also a proof-structure) corresponds to the usual notion of MELL proof-structure (as in [8]) except that we also allow the presence of a box filled only by a *daimon* (*i.e.* a vertex of type $\maltese$). The *empty* (DiLL_0 and MELL) proof-structure—whose module and box-forest are empty graphs—is denoted by ε.

Given a quasi-proof-structure $R = (|R|, \mathcal{F}, \mathsf{box})$, the output tails of $|R|$ are the *conclusions* of R. So, the pre-images of the roots of $\mathcal{F}$ via box partition the conclusions of R in a list of lists of such conclusions. The *type* of R is the list of lists of the types of these conclusions. We often identify the conclusions of R with a finite initial segment of $\mathbb{N}$.

By definition of graph morphism, two conclusions in two distinct lists in the type of a quasi-proof-structure R are in two distinct connected components of $|R|$; so, if R is not a proof-structure then $|R|$ contains several connected components. Thus, R can be seen as a list of proof-structures, its *components*, one for each root in its box-forest.

A non-root vertex v in the box-forest $\mathcal{F}$ induces a subgraph of $\mathcal{F}^{\circlearrowleft}$ of all vertices above it and edges connecting them. The pre-image of this subgraph through box is the *box* of v and the conditions on box in Definition 1 translate the usual nesting condition for LL boxes.

In quasi-proof-structures, we speak of *cells* instead of vertices, and, for a cell of type ℓ, of a *ℓ-cell.* A $\maltese$-cell is a $\maltese_p$-cell for some $p \in \mathbb{N}$. An *hypothesis cell* is a cell without inputs.

▶ **Example 2.** The graph in Figure 2 is a MELL quasi-proof-structure. The colored areas represent the pre-images of boxes, and the dashed boxes represent the pre-images of roots.

The Taylor expansion Proof-structures have a tree structure made explicit by their box-function. Following [18], the definition of the Taylor expansion uses this tree structure: first, we define how to "*expand*" a tree—and more generally a forest—via a generalization of the notion of thick subtree [3] (Definition 3; roughly, a thick subforest of a box-forest says the number of copies of each box to be taken, iteratively), we then take all the expansions of the tree structure of a proof-structure and we *pull* the approximations *back* to the underlying graphs (Definition 5), finally we *forget* the tree structures associated with them (Definition 6).

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[6] Roughly, it says that the border of a box is made of (inputs of) vertices of type ! or ?.

[7] This is a technical condition that simplifies the definition of the rewrite rules in Section 4. Note that $\mathsf{box}(v_1), \mathsf{box}(v_2), \mathsf{box}(v_3)$ are necessarily roots in $\mathcal{F}$, since box is a morphism of directed graphs.

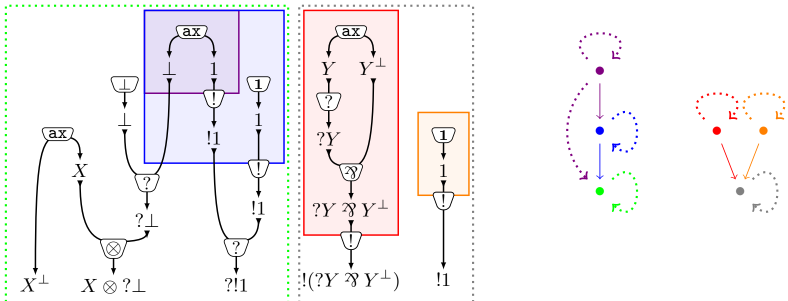
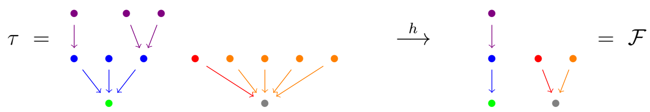


Figure 2 A MELL quasi-proof-structure R, its box-forest $\mathcal{F}_R$ (without dotted lines) and the reflexive-transitive closure $\mathcal{F}_R^{\circlearrowleft}$ of $\mathcal{F}_R$ (with also dotted lines).

▶ **Definition 3** (thick subforest)**.** *Let τ be a forest of rooted trees. A* thick subforest *of τ is a pair (σ, h) of a forest σ of rooted trees and a graph morphism $h\colon \sigma \to \tau$ whose restriction to the roots of σ is bijective.*

▶ **Example 4.** The following is a graphical presentation of a thick subforest (τ, h) of the box-forest $\mathcal{F}$ of the quasi-proof-structure in Figure 2, where the graph morphism $h\colon \tau \to \mathcal{F}$ is depicted chromatically (same color means same image via h).



Intuitively, it means that τ is obtained from $\mathcal{F}$ by taking 3 copies of the blue box, 1 copy of the red box and 4 copies of the orange box; in the first (resp. second; third) copy of the blue box, 1 copy (resp. 0 copies; 2 copies) of the purple box has been taken.

▶ **Definition 5** (proto-Taylor expansion)**.** *Let $R = (|R|, \mathcal{F}_R, \mathsf{box}_R)$ be a quasi-proof-structure. The* proto-Taylor expansion *of R is the set $\mathcal{T}^{\mathsf{proto}}(R)$ of thick subforests of $\mathcal{F}_R$.*

Let $t = (\tau_t, h_t) \in \mathcal{T}^{\mathsf{proto}}(R)$. The t-expansion *of R is the pullback (R_t, p_t, p_R) below, computed in the category of directed graphs and graph morphisms.*

$$\begin{array}{ccc} R_t & \xrightarrow{\;p_t\;} & \tau_t^{\circlearrowleft} \\ \Big\downarrow{\scriptstyle p_R} & \lrcorner & \Big\downarrow{\scriptstyle h_t^{\circlearrowleft}} \\ |R| & \xrightarrow{\;\mathsf{box}_R\;} & \mathcal{F}_R^{\circlearrowleft} \end{array}$$

Given a quasi-proof-structure R and $t = (\tau_t, h_t) \in \mathcal{T}^{\mathsf{proto}}(R)$, the directed graph R_t inherits labels on vertices and edges by composition with the graph morphism $p_R\colon R_t \to |R|$.

Let $[\tau_t]$ be the forest made up of the roots of τ_t and $\iota\colon \tau_t \to [\tau_t]$ be the graph morphism sending each vertex of τ_t to the root below it; $\iota^{\circlearrowleft}$ induces by post-composition a morphism $\overline{h_t} = \iota^{\circlearrowleft} \circ p_t\colon R_t \to [\tau_t]^{\circlearrowleft}$. The triple $(R_t, [\tau_t], \overline{h_t})$ is a DiLL_0 quasi-proof-structure, and it is a DiLL_0 proof-structure if R is a proof-structure. We can then define the *Taylor expansion* $\mathcal{T}(R)$ of a quasi-proof-structure R (an example of an element of a Taylor expansion is in Figure 3).

▶ **Definition 6** (Taylor expansion)**.** *Let R be a quasi-proof-structure. The* Taylor expansion *of R is the set of DiLL_0 quasi-proof-structures $\mathcal{T}(R) = \{(R_t, [\tau_t], \overline{h_t}) \mid t = (\tau_t, h_t) \in \mathcal{T}^{\mathsf{proto}}(R)\}$.*

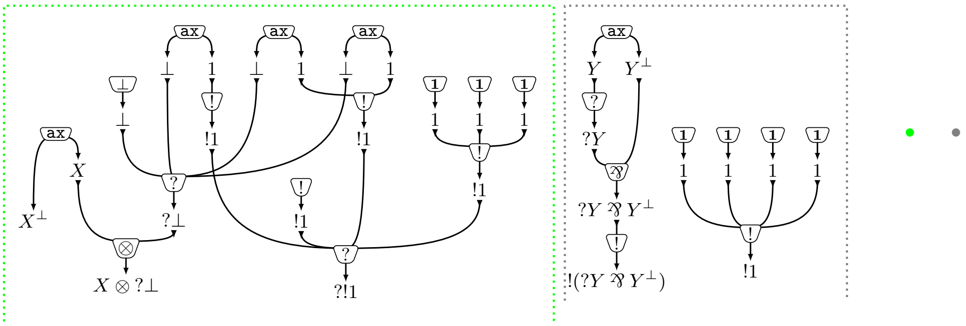


Figure 3 The element of the Taylor expansion of the MELL quasi-proof-structure R in Figure 2, obtained from the element of $\mathcal{T}^{\mathsf{proto}}(R)$ depicted in Example 4.

An element $(R_t, [\tau_t], \overline{h_t})$ of the Taylor expansion of a quasi-proof-structure R has much less structure than the pullback (R_t, p_t, p_R): the latter indeed is a DiLL_0 quasi-proof-structure R_t coming with its projections $|R| \xleftarrow{p_R} R_t \xrightarrow{p_t} \tau_t^{\circlearrowleft}$, which establish a precise correspondence between cells and edges of R_t and cells and edges of R: a cell in R_t is labeled (via the projections) by both the cell of $|R|$ and the branch of the box-forest of R it arose from. But $(R_t, [\tau_t], \overline{h_t})$ where R_t is without its projections p_t and p_R loses the correspondence with R.

▶ Remark 7. By definition, the Taylor expansion preserves conclusions: there is a bijection φ from the conclusions of a quasi-proof-structure R to the ones in each element ρ of $\mathcal{T}(R)$ such that i and $\varphi(i)$ have the same type and the same root (*i.e.* $\mathsf{box}_R(i) = \mathsf{box}_\rho(\varphi(i))$ up to isomorphism). Therefore, the types of R and ρ are the same (as a list of lists).

The filled Taylor expansion As discussed in Section 2 (p. 4), our method needs to "represent" the emptiness introduced by the Taylor expansion (taking 0 copies of a box) so as to preserve the conclusions. So, an element of the *filled Taylor expansion* $\mathcal{T}^{\maltese}(R)$ of a quasi-proof-structure R (an example is in Figure 4) is obtained from an element of $\mathcal{T}(R)$ where a whole component can be erased and replaced by a $\maltese$-cell with the same conclusions (hence $\mathcal{T}(R) \subseteq \mathcal{T}^{\maltese}(R)$).

▶ **Definition 8** (filled Taylor expansion). *An* emptying *of a* DiLL_0 *quasi-proof-structure* $\rho = (|\rho|, \mathcal{F}, \mathsf{box})$ *is the* DiLL_0 *quasi-proof-structure with the same conclusions as* ρ*, obtained from* ρ *by replacing each of the components of some roots of* $\mathcal{F}$ *with a* $\maltese$*-cell whose outputs are tails.*

The filled Taylor expansion $\mathcal{T}^{\maltese}(R)$ *of a quasi-proof-structure* R *is the set of all the emptyings of every element of its Taylor expansion* $\mathcal{T}(R)$*.*

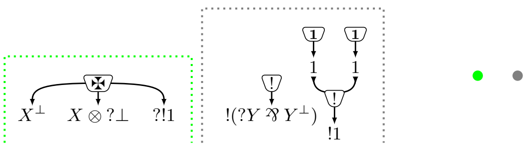


Figure 4 An element of the filled Taylor expansion of the MELL quasi-proof-structure in Figure 2.

4 Means of destruction: unwinding MELL quasi-proof-structures

Our aim is to deconstruct proof-structures (be they $\mathsf{MELL}^{\maltese}$ or DiLL_0) from their conclusions. To do that, we introduce a category of rules of deconstruction. The morphisms of this category are sequences of deconstructing rules, acting on lists of lists of formulæ. These morphisms act through functors on quasi-proof-structures, exhibiting their sequential structure.

$$\begin{array}{rcll}
(\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i),\mathsf{c}(i{+}1),\Gamma'_k;\cdots;\Gamma_n) & \xrightarrow{\mathtt{exc}_i} & (\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i{+}1),\mathsf{c}(i),\Gamma'_k;\cdots;\Gamma_n) & \\
(\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i),\mathsf{c}(i{+}1),\Gamma'_k;\cdots;\Gamma_n) & \xrightarrow{\mathtt{mix}_i} & (\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i);\mathsf{c}(i{+}1),\Gamma'_k;\cdots;\Gamma_n) & \\
(\Gamma_1;\cdots;\Gamma_k;\mathsf{c}(i),\mathsf{c}(i{+}1);\Gamma_{k+2};\cdots;\Gamma_n) & \xrightarrow{\mathtt{ax}_i} & (\Gamma_1;\cdots;\Gamma_k;\Gamma_{k+2};\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = A = \mathsf{c}(i{+}1)^\perp \\
(\Gamma_1;\cdots;\Gamma_k;\cdots;\Gamma_n) & \xrightarrow{\mathtt{cut}^i} & (\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i),\mathsf{c}(i{+}1);\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = A = \mathsf{c}(i{+}1)^\perp \\
(\Gamma_1;\cdots;\Gamma_k;\Gamma_{k+1},\mathsf{c}(i);\Gamma_{k+2};\cdots;\Gamma_n) & \xrightarrow{\maltese_i} & (\Gamma_1;\cdots;\Gamma_k;\Gamma_{k+2};\cdots;\Gamma_n) & \\
(\Gamma_1;\cdots;\Gamma_k;\mathsf{c}(i);\Gamma_{k+2};\cdots;\Gamma_n) & \xrightarrow{\mathbf{1}_i} & (\Gamma_1;\cdots;\Gamma_k;\Gamma_{k+2};\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = \mathbf{1} \\
(\Gamma_1;\cdots;\Gamma_k;\mathsf{c}(i);\Gamma_{k+2};\cdots;\Gamma_n) & \xrightarrow{\perp_i} & (\Gamma_1;\cdots;\Gamma_k;\Gamma_{k+2};\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = \perp \\
(\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i);\cdots;\Gamma_n) & \xrightarrow{\otimes_i} & (\Gamma_1;\cdots;\Gamma_k,A,B;\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = A\otimes B \\
(\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i);\cdots;\Gamma_n) & \xrightarrow{\parr_i} & (\Gamma_1;\cdots;\Gamma_k,A,B;\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = A\parr B \\
(\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i);\cdots;\Gamma_n) & \xrightarrow{?_{c\,i}} & (\Gamma_1;\cdots;\Gamma_k,?A,?A;\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = ?A \\
(\Gamma_1;\cdots;\Gamma_k,\mathsf{c}(i);\cdots;\Gamma_n) & \xrightarrow{?_{d\,i}} & (\Gamma_1;\cdots;\Gamma_k,A;\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = ?A \\
(\Gamma_1;\cdots;\Gamma_k;\mathsf{c}(i);\Gamma_{k+2};\cdots;\Gamma_n) & \xrightarrow{?_{w\,i}} & (\Gamma_1;\cdots;\Gamma_k;\Gamma_{k+2};\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = ?A \\
(\Gamma_1;\cdots;?\Gamma_k,\mathsf{c}(i);\cdots;\Gamma_n) & \xrightarrow{\mathsf{Box}_i} & (\Gamma_1;\cdots;?\Gamma_k,A;\cdots;\Gamma_n) & \text{with } \mathsf{c}(i) = !A
\end{array}$$

Figure 5 The generators of **Path**. In the source $\Gamma = (A_1,\ldots,A_{i_1};\cdots;A_{i_{m-1}+1},\ldots,A_{i_n})$ of each arrow, $\mathsf{c}(i)$ denotes the i^{th} formula in the flattening $(A_1,\ldots,A_{i_1},\ldots,A_{i_{m-1}+1},\ldots,A_{i_n})$ of Γ.

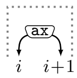
▶ **Definition 9** (the category **Path**). *Let* **Path** *be the category whose*

- *objects are lists* $\Gamma = (\Gamma_1;\ldots;\Gamma_n)$ *of lists of* MELL *formulæ;*
- *arrows are freely generated by the* elementary paths *in Figure 5.*

We call a path *any arrow* $\xi\colon \Gamma \to \Gamma'$. *We write the composition of paths without symbols and in the diagrammatic order, so, if* $\xi\colon \Gamma \to \Gamma'$ *and* $\xi'\colon \Gamma' \to \Gamma''$, $\xi\xi'\colon \Gamma \to \Gamma''$.

▶ **Example 10.** $\parr_1\ \parr_2\ \parr_3\ \otimes_1\ \otimes_3\ \mathtt{exc}_1\ \mathtt{exc}_2\ \mathtt{mix}_2\ \mathtt{ax}_1\ \mathtt{exc}_2\ \mathtt{mix}_2\ \mathtt{ax}_1\ \mathtt{ax}_1$ is a path of type $\big((X\otimes Y^\perp)\parr((Y\otimes Z^\perp)\parr(X^\perp\parr Z))\big) \longrightarrow \varepsilon$, where ε is the *empty list* of lists of formulæ.

We will tend to forget about exchanges and perform them silently (as it is customary, for instance, in most presentations of sequent calculi).

The category **Path** acts on $\mathsf{MELL}^{\maltese}$ quasi-proof-structures, exhibiting a sequential structure in their construction. For Γ a list of lists of MELL formulæ, $\mathbf{qMELL}^{\maltese}(\Gamma)$ is the set of $\mathsf{MELL}^{\maltese}$ quasi-proof-structures of type Γ. To ease the reading of the rewrite rules acting on a $\mathsf{MELL}^{\maltese}$ quasi-proof-structures R, we will only draw the parts of R belonging to the relevant component; *e.g.*, if we are interested in an $\mathtt{ax}$-cell whose outputs are the conclusions i and $i{+}1$, and it is the only cell in a component, we will write  ignoring the rest.

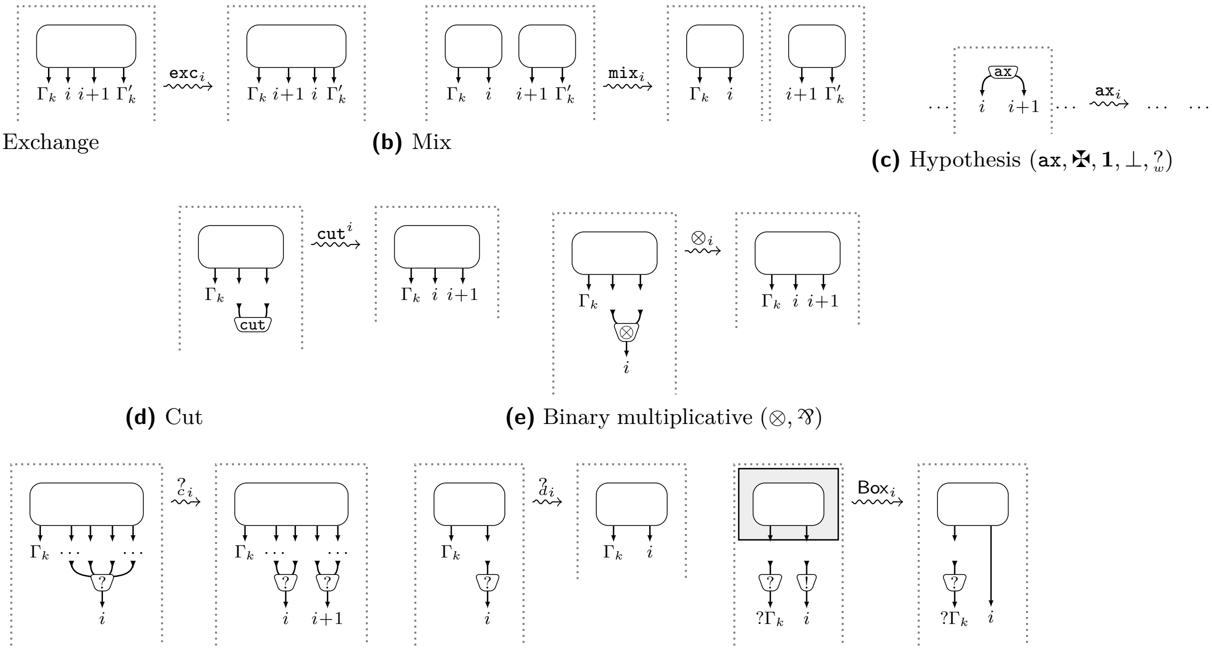
▶ **Definition 11** (action of paths on MELL quasi-proof-structures). *An elementary path* $a\colon \Gamma \to \Gamma'$ *defines a relation* $\overset{a}{\rightsquigarrow} \subseteq \mathbf{qMELL}^{\maltese}(\Gamma) \times \mathbf{qMELL}^{\maltese}(\Gamma')$ *(the* action *of* a*) as the smallest relation containing all the cases in Figure 6, with the following remarks:*

mix *read in reverse, a quasi-proof-structure with two components is in relation with a quasi-proof-structure with the same module but the two roots of such components merged.*

hypothesis *if* $a \in \{\mathtt{ax}_i, \maltese_i, \mathbf{1}_i, \perp_i, ?_{w\,i}\}$, *the rules have all in common to act by deleting a cell without inputs that is the only cell in its component. We have drawn the axiom case in Figure 6c, the others vary only by their number of conclusions.*

cut *read in reverse, a quasi-proof-structure with two conclusions* i *and* $i+1$ *is in relation with the quasi-proof-structure where these two conclusions are cut. This rule, from left to right, is non-deterministic (as there are many possible cuts).*

binary multiplicatives *these rules delete a binary connective. We have only drawn the* $\otimes$ *case in Figure 6e, the* $\parr$ *case is similar.*



(a) Exchange **(b)** Mix **(c)** Hypothesis ($\mathtt{ax}, \maltese, \mathbf{1}, \bot, ?_w$)

(d) Cut **(e)** Binary multiplicative ($\otimes, ⅋$)

(f) Contraction **(g)** Dereliction **(h)** Box

Figure 6 Actions of elementary paths on MELL$^{\maltese}$ quasi-proof-structures.

contraction *splits a ?-cell with $h{+}k{+}2$ inputs into two ?-cells with $h{+}1$ and $k{+}1$ inputs, respectively.*

dereliction *only applies if the ?-cell (with 1 input) does not shift a level in the box-forest.*

box *only applies if a box (and its border) is alone in its component.*

This definition of the rewrite system is extended to define a relation $\overset{\xi}{\rightsquigarrow} \subseteq \mathbf{qMELL}^{\maltese}(\Gamma) \times \mathbf{qMELL}^{\maltese}(\Gamma')$ (the action *of any path $\xi\colon \Gamma \to \Gamma'$) by composition of relations.*

Given two MELL$^{\maltese}$ quasi-proof-structures R and R', we say that a rule a *applies* to R if there is a finite sequence of exchanges $\mathtt{exc}_{i_1} \cdots \mathtt{exc}_{i_n}$ such that $R \overset{\mathtt{exc}_{i_1} \cdots \mathtt{exc}_{i_n} a}{\rightsquigarrow} R'$.

▶ **Definition 12** (the functor $\mathbf{qMELL}^{\maltese}$). *We define a functor $\mathbf{qMELL}^{\maltese}\colon \mathbf{Path} \to \mathbf{Rel}$ by:*

- on objects: $\mathbf{qMELL}^{\maltese}(\Gamma)$ *is the set of* MELL$^{\maltese}$ *quasi-proof-structures of type Γ;*
- on morphisms: *for $\xi\colon \Gamma \to \Gamma'$, $\mathbf{qMELL}^{\maltese}(\xi) = \overset{\xi}{\rightsquigarrow}$ (see Definition 11).*

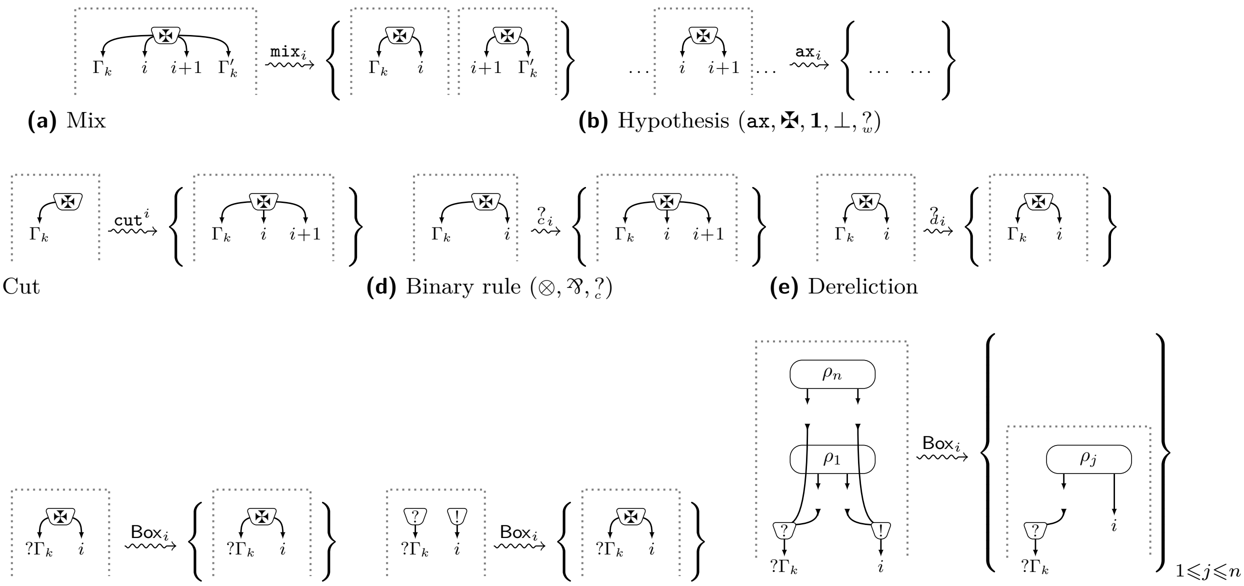
Our rewrite rules enjoy two useful properties, expressed by Propositions 13 and 15.

▶ **Proposition 13** (co-functionality). *Let $\xi\colon \Gamma \to \Gamma'$ be a path. The relation $\overset{\xi}{\rightsquigarrow}$ is a co-function on the sets of underlying graphs, that is, a function $\overset{\xi}{\rightsquigarrow}^{\mathrm{op}}\colon \mathbf{qMELL}^{\maltese}(\Gamma') \to \mathbf{qMELL}^{\maltese}(\Gamma)$.*

▶ **Lemma 14** (applicability of rules). *Let R be a non-empty* MELL$^{\maltese}$ *quasi-proof-structure. There exists a conclusion i such that:*

- *either a rule in $\{\mathtt{ax}_i, \mathbf{1}_i, \bot_i, \otimes_i, ⅋_i, ?_{c\,i}, ?_{d\,i}, ?_{w\,i}, \mathtt{cut}^i, \maltese_i, \mathsf{Box}_i\}$ applies to R;*
- *or $R \overset{\mathtt{mix}_i}{\rightsquigarrow} R'$ (where the conclusions affected by $\mathtt{mix}_i$ are $i{-}k, \dots, i, i{+}1, \dots, i{+}\ell$) and $i{-}k, \dots, i$ are all the conclusions of either a box or an hypothesis cell, and one of the components of R' coincides with this cell or box (and its border).*

Proposition 13 and Lemma 14 are proven by simple inspection of the rewrite rules of Figure 6.



(a) Mix (b) Hypothesis ($\mathtt{ax}, \maltese, \mathbf{1}, \bot, ?_w$) (c) Cut (d) Binary rule ($\otimes, \parr, ?_c$) (e) Dereliction (f) Daimoned box (g) Empty box (h) Non-empty box ($n > 0$)

Figure 7 Actions of elementary paths on $\maltese$-cells and on a box in $\mathbf{qDiLL}_0^{\maltese}$.

▶ **Proposition 15** (termination). *Let R be a* MELL$^{\maltese}$ *quasi-proof-structure of type Γ. There exists a path $\xi\colon \Gamma \to \varepsilon$ such that $R \overset{\xi}{\rightsquigarrow} \varepsilon$.*

To prove Proposition 15, it is enough to apply Lemma 14 and show that the size of MELL$^{\maltese}$ quasi-proof-structures decreases for each application of the rules in Figure 6, according to the following definition of size. The *size* of a proof-structure R is the couple (p, q) where

- p is the (finite) multiset of the number of inputs of each ?-cell in R;
- q is the number of cells not labeled by $\maltese$ in R.

The *size* of a quasi-proof-structure R is the (finite) multiset of the sizes of its components. Multisets are ordered as usual, couples are ordered lexicographically.

5 Naturality of unwinding DiLL$_0^{\maltese}$ quasi-proof-structures

For Γ a list of lists of MELL formulæ, $\mathbf{qDiLL}_0^{\maltese}(\Gamma)$ is the set of DiLL$_0^{\maltese}$ quasi-proof-structures of type Γ. For any set X, its powerset is denoted by $\mathfrak{P}(X)$.

▶ **Definition 16** (action of paths on DiLL$_0^{\maltese}$ quasi-proof-structures). *An elementary path $a\colon \Gamma \to \Gamma'$ defines a relation $\overset{a}{\rightsquigarrow} \subseteq \mathbf{qDiLL}_0^{\maltese}(\Gamma) \times \mathfrak{P}(\mathbf{qDiLL}_0^{\maltese}(\Gamma'))$ (the* action *of a) by the rules in Figure 6 (except Figure 6h, and with all the already remarked notes) and in Figure 7.*

We extend this relation on $\mathfrak{P}(\mathbf{qDiLL}_0^{\maltese}(\Gamma)) \times \mathfrak{P}(\mathbf{qDiLL}_0^{\maltese}(\Gamma'))$ by the monad multiplication of $X \mapsto \mathfrak{P}(X)$ and define $\overset{\xi}{\rightsquigarrow}$ (the action *of any path $\xi\colon \Gamma \to \Gamma'$) by composition of relations.*

Roughly, all the rewrite rules in Figure 7—except Figure 7h—mimic the behavior of the corresponding rule in Figure 6 using a $\maltese$-cell. Note that in Figure 7g a $\maltese$-cell is created.

The non-empty box rule in Figure 7h requires that, on the left of $\overset{\mathsf{Box}_i}{\rightsquigarrow}$, ρ_j is not connected to $\rho_{j'}$ for $j \neq j'$, except for the !-cell and the ?-cells in the conclusions. Read in reverse, the rule associates with a non-empty finite set of DiLL$_0$ quasi-proof-structures $\{\rho_1, \dots, \rho_n\}$ the merging of $\rho_1, \dots, \rho_n$, that is the DiLL$_0$ quasi-proof-structure depicted on the left of $\overset{\mathsf{Box}_i}{\rightsquigarrow}$.

▶ **Definition 17** (the functor $\mathfrak{P}\mathbf{qDiLL}_0^{\maltese}$). *We define a functor* $\mathfrak{P}\mathbf{qDiLL}_0^{\maltese} : \mathbf{Path} \to \mathbf{Rel}$ *by:*

- on objects: *for* Γ *a list of lists of* MELL *formulæ,* $\mathfrak{P}\mathbf{qDiLL}_0^{\maltese}(\Gamma) = \mathfrak{P}(\mathbf{qDiLL}_0^{\maltese}(\Gamma))$*, the set of sets of* $\mathsf{DiLL}_0^{\maltese}$ *quasi-proof-structures of type* Γ*;*
- on morphisms: *for* $\xi : \Gamma \to \Gamma'$*,* $\mathfrak{P}\mathbf{qDiLL}_0^{\maltese}(\xi) = \overset{\xi}{\rightsquigarrow}$ *(see Definition 16).*

▶ **Theorem 18** (naturality). *The filled Taylor expansion defines a natural transformation* $\mathfrak{T}^{\maltese} : \mathfrak{P}\mathbf{qDiLL}_0^{\maltese} \Rightarrow \mathbf{qMELL}^{\maltese} : \mathbf{Path} \to \mathbf{Rel}$ *by:* $(\Pi, R) \in \mathfrak{T}_\Gamma^{\maltese}$ *iff* $\Pi \subseteq \mathcal{T}^{\maltese}(R)$ *and the type of* R *is* Γ*. Moreover, if* Π *is a set of* DiLL_0 *proof-structures with* $\Pi \overset{\xi}{\rightsquigarrow} \Pi'$ *and* $\Pi' \subseteq \mathcal{T}(R')$*, then* R *is a* MELL *proof-structure and* $\Pi \subseteq \mathcal{T}(R)$*, where* R *is such that* $R \overset{\xi}{\rightsquigarrow} R'$*.*[8]

In other words, the following diagram commutes for every path $\xi \colon \Gamma \to \Gamma'$.

$$\begin{array}{ccc} \mathfrak{P}\mathbf{qDiLL}_0^{\maltese}(\Gamma) & \xrightarrow{\mathfrak{P}\mathbf{qDiLL}_0^{\maltese}(\xi)} & \mathfrak{P}\mathbf{qDiLL}_0^{\maltese}(\Gamma') \\ \downarrow \mathfrak{T}_\Gamma^{\maltese} & & \downarrow \mathfrak{T}_{\Gamma'}^{\maltese} \\ \mathbf{qMELL}^{\maltese}(\Gamma) & \xrightarrow{\mathbf{qMELL}^{\maltese}(\xi)} & \mathbf{qMELL}^{\maltese}(\Gamma') \end{array}$$

It means that given $\Pi \overset{\xi}{\rightsquigarrow} \Pi'$, where $\Pi' \subseteq \mathcal{T}^{\maltese}(R')$, we can simulate backwards the rewriting to R (this is where the co-functionality of the rewriting steps expressed by Proposition 13 comes handy) so that $R \overset{\xi}{\rightsquigarrow} R'$ and $\Pi \subseteq \mathcal{T}^{\maltese}(R)$; and conversely, given $R \overset{\xi}{\rightsquigarrow} R'$, we can simulate the rewriting for any $\Pi \subseteq \mathcal{T}^{\maltese}(R)$, so that $\Pi \overset{\xi}{\rightsquigarrow} \Pi'$ for some $\Pi' \subseteq \mathcal{T}^{\maltese}(R')$.

6 Glueability of DiLL_0 quasi-proof-structures

Naturality (Theorem 18) allows us to characterize the sets of DiLL_0 proof-structures that are in the Taylor expansion of some MELL proof-structure (Theorem 20 below).

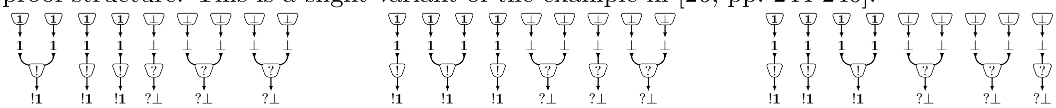
▶ **Definition 19** (glueability). *We say that a set* Π *of* $\mathsf{DiLL}_0^{\maltese}$ *quasi-proof-structures is* glueable*, if there exists a path* ξ *such that* $\Pi \overset{\xi}{\rightsquigarrow} \{\varepsilon\}$*.*

▶ **Theorem 20** (glueability criterion). *Let* Π *be a set of* DiLL_0 *proof-structures:* Π *is glueable if and only if* $\Pi \subseteq \mathcal{T}(R)$ *for some* MELL *proof-structure* R*.*

Proof. If $\Pi \subseteq \mathcal{T}(R)$ for some MELL proof-structure R, then by termination (Proposition 15) $R \overset{\xi}{\rightsquigarrow} \varepsilon$ for some path ξ, and so $\Pi \overset{\xi}{\rightsquigarrow} \{\varepsilon\}$ by naturality (Theorem 18, as $\mathcal{T}^{\maltese}(\varepsilon) = \{\varepsilon\}$).

Conversely, if $\Pi \overset{\xi}{\rightsquigarrow} \{\varepsilon\}$ for some path ξ, then by naturality (Theorem 18, as $\mathcal{T}(\varepsilon) = \{\varepsilon\}$ and Π is a set of DiLL_0 proof-structures) $\Pi \subseteq \mathcal{T}(R)$ for some MELL proof-structure R. ◀

▶ **Example 21.** The three DiLL_0 proof-structures ρ_1, ρ_2, ρ_3 below are not glueable as a whole, but are glueable two by two. In fact, there is no MELL proof-structure whose Taylor expansion contains ρ_1, ρ_2, ρ_3, but any pair of them is in the Taylor expansion of some MELL proof-structure. This is a slight variant of the example in [26, pp. 244-246].



An example of the action of a path starting from a DiLL_0 proof-structure ρ and ending in $\{\varepsilon\}$ can be found in Figures 8 and 9. Note that it is by no means the shortest possible path. When replayed backwards, it induces a MELL proof-structure R such that $\rho \in \mathcal{T}(R)$.

[8] The part of the statement after "moreover" is our way to control the presence of ✠-cells.

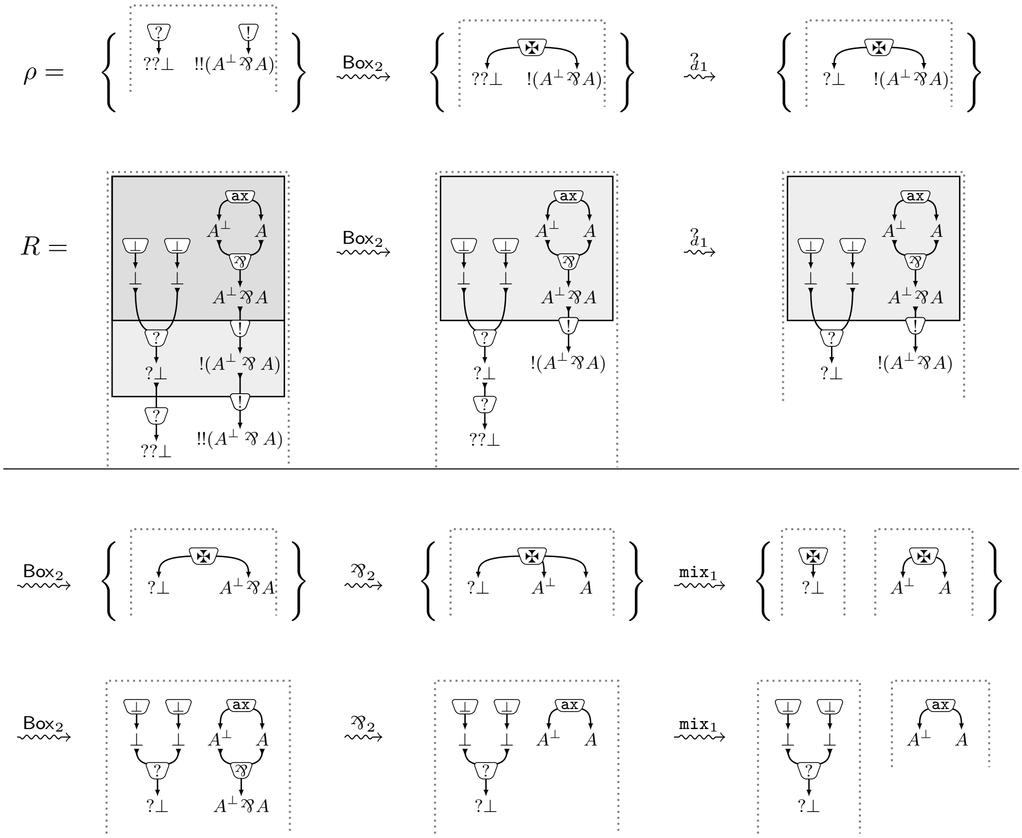
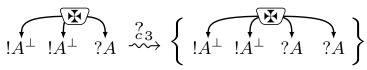
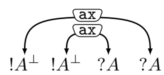


Figure 8 The path $\mathsf{Box}_2\, ?_{d1}\, \mathsf{Box}_2\, \parr_2\, \mathtt{mix}_1\, \mathtt{ax}_2\, ?_{c1}\, ?_{d2}\, \mathtt{mix}_1\, \bot_2\, ?_{d1}\, \bot_1$ witnessing that $\rho \in \mathcal{T}(R)$ (to be continued on Figure 9).

7 Non-atomic axioms

From now on, we relax the definition of quasi-proof-structure (Definition 1 and Figure 1) so that the outputs of any $\mathtt{ax}$-cell are labeled by dual MELL formulæ, not necessarily atomic. We can extend our results to this more general setting, with some technical complications. Indeed, the rewrite rule for contraction has to be modified. Consider a set of DiLL_0 proof-structures consisting of just a singleton which is a ✠-cell. The contraction rule rewrites it as:

 which is then in the Taylor expansion of 

on which no contraction rewrite rule $?_c$ can be applied backwards, breaking the naturality. The failure of the naturality is actually due to the failure of Proposition 13 in the case of the rewrite rule $?_c$: $\overset{?_c}{\rightsquigarrow}^{\mathrm{op}}$ (*i.e.* $\overset{?_c}{\rightsquigarrow}$ read from the right to the left) is functional but not total.

The solution to this conundrum lies in changing the contraction rule for $\mathsf{DiLL}_0^{✠}$ quasi-proof-structures, by explicitly adding ?-cells. Hence, the application of a contraction step $?_c$ in the $\mathsf{DiLL}_0^{✠}$ quasi-proof-structures precludes the possibility of anything else but a ?-cell on the $\mathsf{MELL}^{✠}$ side, which allows the contraction step $?_c$ to be applied backwards.

In turn, this forces us to change the definition of the filled Taylor expansion into a *η-filled Taylor expansion*, which has to include elements where a ✠-cell (representing an empty component) has some of its outputs connected to ?-cells.

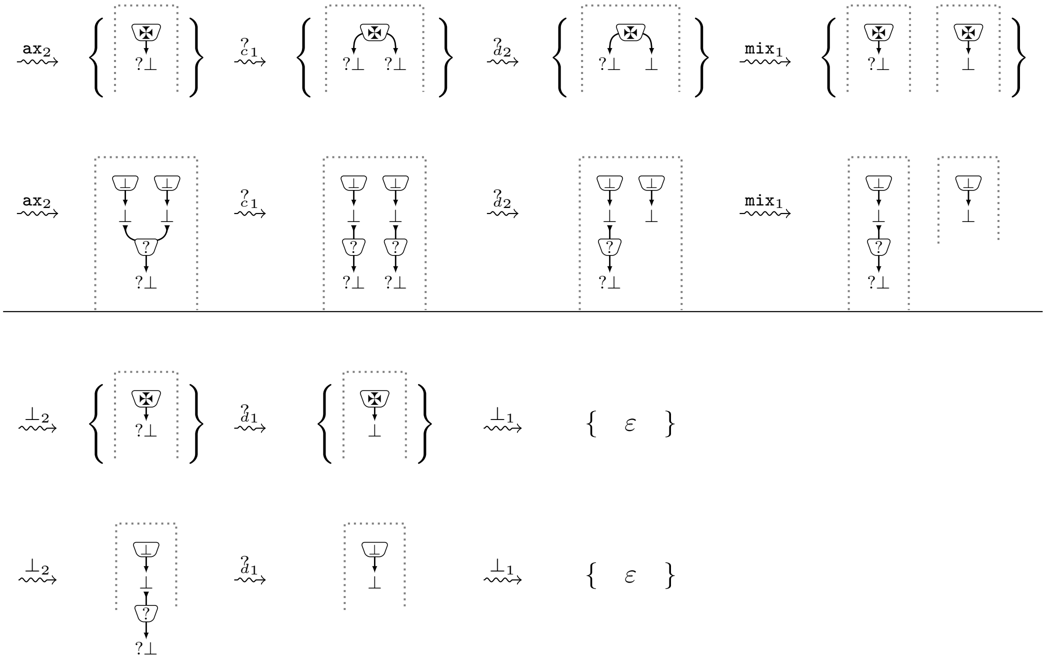


Figure 9 The path $\mathsf{Box}_2 \, ?_{d1} \, \mathsf{Box}_2 \, ⅋_2 \, \mathtt{mix}_1 \, \mathtt{ax}_2 \, ?_{c1} \, ?_{d2} \, \mathtt{mix}_1 \, \bot_2 \, ?_{d1} \, \bot_1$ witnessing that $\rho \in \mathcal{T}(R)$ (continued from Figure 8).

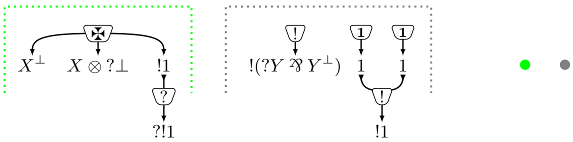


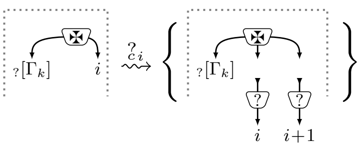
Figure 10 An element of the η-filled Taylor expansion of the MELL quasi-proof-structure in Fig. 2.

▶ **Definition 22** (η-filled Taylor expansion). *An* η-emptying *of a* DiLL_0 *quasi-proof-structure* $\rho = (|\rho|, \mathcal{F}, \mathsf{box})$ *is a* DiLL_0 *quasi-proof-structure with the same conclusions as* ρ*, obtained from* ρ *by replacing each of the components of some roots of* $\mathcal{F}$ *with a* ✠*-cell whose outputs are either tails or inputs of a* $?$*-cell whose output* i *is a tail, provided that* i *is the output tail of a* $?$*-cell in* ρ*.*

The η-filled Taylor expansion $\mathcal{T}_\eta^{✠}(R)$ *of a quasi-proof-structure* R *is the set of all the* η*-emptyings of every element of its Taylor expansion* $\mathcal{T}(R)$*.*

Note that the η-filled Taylor expansion contains all the elements of the filled Taylor expansion and some more, such as the one in Figure 10.

Functors $\mathbf{qMELL}^{✠}$ and $\mathfrak{P}\mathbf{qDiLL}_0^{✠}$ are defined as before (Def. 12 and 17, respectively),[9] except that the image of $\mathfrak{P}\mathbf{qDiLL}_0^{✠}$ on the generator $?_{ci}$ (Figure 7d) is changed to



[9] Remember that now, for Γ a list of lists of MELL formulæ, $\mathbf{qMELL}^{✠}(\Gamma)$ (resp. $\mathbf{qDiLL}_0^{✠}(\Gamma)$) is the set of $\mathsf{MELL}^{✠}$ (resp. $\mathsf{DiLL}_0^{✠}$) quasi-proof-structures of type Γ, possibly with non-atomic axioms.

where $_?[\Gamma_k]$ signifies that some of the conclusions of Γ_k might be connected to the ✠-cell through a ?-cell. We can prove similarly our main results.

▶ **Theorem 23** (naturality with η)**.** *The η-filled Taylor expansion defines a natural transformation $\mathfrak{T}_\eta^{\maltese} : \mathfrak{P}\mathbf{qDiLL}_0^{\maltese} \Rightarrow \mathbf{qMELL}^{\maltese} : \mathbf{Path} \to \mathbf{Rel}$ by: $(\Pi, R) \in \mathfrak{T}^{\maltese}_{\eta\,\Gamma}$ iff $\Pi \subseteq \mathcal{T}_\eta^{\maltese}(R)$ and the type of R is Γ. Moreover, if Π is a set of* DiLL$_0$ *proof-structures with $\Pi \overset{\xi}{\rightsquigarrow} \Pi'$ and $\Pi' \subseteq \mathcal{T}(R')$, then R is a* MELL *proof-structure and $\Pi \subseteq \mathcal{T}(R)$, where R is such that $R \overset{\xi}{\rightsquigarrow} R'$.*

▶ **Theorem 24** (glueability criterion with η)**.** *Let Π be a set of* DiLL$_0$ *proof-structures, not necessarily with atomic axioms: Π is glueable iff $\Pi \subseteq \mathcal{T}(R)$ for some* MELL *proof-structure R.*

8 Conclusions and perspectives

✠-cells inside boxes Our glueability criterion (Theorem 20) solves the inverse Taylor expansion problem in a “asymmetric” way: we characterize the sets of DiLL$_0$ proof-structures that are included in the Taylor expansion of some MELL proof-structure, but DiLL$_0$ proof-structures have no occurrences of ✠-cells, while a MELL proof-structure possibly contains ✠-cells inside boxes (see Definition 1). Not only this asymmetry is technically inevitable, but it reflects on the fact that some glueable set of DiLL$_0$ proof-structures might not contain any information on the content of some box (which is reified in MELL by a ✠-cell), or worse that, given the types, no content can fill that box. Think of the DiLL$_0$ proof-structure ρ made only of a !-cell with no inputs and one output of type $!X$, where X is atomic: $\{\rho\}$ is glueable but the only MELL proof-structure R such that $\{\rho\} \subseteq \mathcal{T}(R)$ is made of a box containing a ✠-cell.

This asymmetry is also present in Pagani and Tasson’s characterization [23], even if not particularly emphasized: their Theorem 2 (analogous to the left-to-right part of our Theorem 20) assumes not only that the rewriting starting from a finite set of DiLL$_0$ proof-structures terminates but also that it ends on a MELL proof-structure (without ✠-cells, which ensures that there exists a MELL proof-structure without ✠-cells filling all the empty boxes).

The λ-calculus, connectedness and coherence Our rewriting system and glueability criterion should help to prove the existence of a binary coherence for elements of the Taylor expansion of a fragment of MELL proof-structures (despite the impossibility for full MELL proved in [26]), extending the one that exists for resource λ-terms. We can remark that our glueability criterion is actually an extension of the criterion for resource λ-terms. Indeed, in the case of the λ-calculus, there are three rewrite steps, corresponding to abstraction, application and variable (which can be encoded in our rewrite steps), and coherence is defined inductively: if a set of resource λ-terms is coherent, then any set of resource λ-term that rewrites to it is also coherent.

Presented in this way, the main difference between the λ-calculus and MELL (concerning the inverse Taylor expansion problem) would not be because of the rewriting system but because the structure of any resource λ-term univocally determines the rewriting path, while, for DiLL$_0$ proof-structures, we have to quantify existentially over all possible paths. This is an unavoidable consequence of the fact that proof-structures do not have a tree-structure, contrary to λ-terms and resource λ-terms.

Moreover, it is possible to match and mix different sequences of rewriting. Indeed, consider three DiLL$_0$ proof-structures pairwise glueable. Proving that they are glueable as a whole amounts to computing a rewriting path from the rewriting paths witnessing the three glueabilities. Our paths were designed with that mixing-and-matching operation in mind, in the particular case where the boxes are connected. This is reminiscent of [16], where we also showed that a certain property enjoyed by the λ-calculus can be extended to proof-structures, provided they are connected inside boxes. We leave that work to a subsequent paper.

Functoriality and naturality Our functorial point of view on proof-structures might unify many results. Let us cite two of them:

- a sequent calculus proof of $\vdash \Gamma$ can be translated into a path from the empty sequence into Γ. This could be the starting point for the formulation of a new correctness criterion;
- the category **Path** can be extended with higher structure, allowing to represent cut-elimination. The functors $\mathbf{qMELL}^{\maltese}$ and $\mathfrak{P}\mathbf{qDiLL}_0^{\maltese}$ can also be extended to such higher functors, proving via naturality that cut-elimination and the Taylor expansion commute.

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