

# MPRI, Fondations mathématiques de la théorie des automates

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**Avertissement :** On attachera une grande importance à la clarté, à la précision et à la concision de la rédaction.

## Partie 1. Monoïde syntaxique de $(abab)^*$ .

**Question 1.** Soit  $A = \{a, b\}$ . On considère le langage  $L = (abab)^*$ . Calculer son automate minimal.

**Question 2.** Calculer le monoïde syntaxique  $M$  de  $L$ . On donnera la liste des éléments et des relations permettant de définir  $M$  (vous devriez trouver 10 éléments, en comptant l'élément neutre).

**Question 3.** Donner la liste des idempotents de  $M$ .

**Question 4.** Déterminer la structure en  $\mathcal{D}$ -classes de  $M$  et dessiner les diagrammes boîtes à œufs.

**Question 5.** Le langage  $L$  est-il sans-étoile? Justifier votre réponse.

**Question 6.** Montrer que  $L$  vérifie l'équation profinie  $x^\omega y^\omega = y^\omega x^\omega$  pour tout  $x, y \in A^*$ .

## Partie 2. Ensembles préfixes maximaux finis.

On rappelle qu'un préfixe  $p$  d'un mot  $u$  est *propre* si  $p \neq u$ . Un ensemble  $X$  de mots non vides de  $A^+$  est dit *préfixe* si aucun mot de  $X$  n'est préfixe propre d'un autre mot de  $X$  (autrement dit, si  $u, uv \in X$  entraîne  $v = 1$ ). Un ensemble préfixe  $X$  est dit *maximal* si, pour tout mot  $u \in A^+$ , l'ensemble  $X \cup \{u\}$  n'est pas préfixe.

Le but de l'exercice est d'étudier certaines propriétés des ensembles préfixes maximaux finis.

**Question 7.** Soit  $X$  un ensemble préfixe et soit  $u \in X^*$ . Montrer que  $v \in X^*$  si et seulement si  $uv \in X^*$ .

**Question 8.** Montrer que l'ensemble

$$Z = \{aa, ab, baa, bab, bb\}$$

est un ensemble préfixe maximal.

Soit  $X$  un ensemble préfixe maximal fini et soit  $P$  l'ensemble des préfixes propres de ses mots. Par exemple, si  $X = Z$ , on a  $P = \{1, a, b, ba\}$ .

**Question 9.** En utilisant la maximalité de  $X$ , montrer que  $A^* = P \cup XA^*$  et que cette union est disjointe.

**Question 10.** Déduire de la question précédente que  $A^* = X^*P$  et que tout mot  $u$  de  $A^*$  admet une factorisation unique de la forme  $u = x_1 \cdots x_n p$  avec  $n \geq 0$ ,  $x_1, \dots, x_n \in X$  et  $p \in P$ .

**Question 11.** Soit  $\mathcal{A}$  l'automate minimal déterministe de  $X^*$ . Montrer que l'état initial  $q_0$  de  $\mathcal{A}$  est aussi le seul état final.

On suppose maintenant que  $X$  est un ensemble préfixe maximal fini qui est en plus *suffixe* (aucun mot de  $X$  n'est suffixe propre d'un autre mot de  $X$ ). Soit  $n$  la longueur maximale des mots de  $X$ .

**Question 12.** Soit  $q$  un état de  $\mathcal{A}$  et  $u$  un mot. Montrer que si  $q_0 \cdot u = q \cdot u$ , alors  $q = q_0$ .

**Question 13.** Soit  $u$  un mot de longueur  $\geq n$  et soient  $q_1, q_2$  deux états de  $\mathcal{A}$ . Montrer que s'il existe un mot  $v$  tel que  $q_1 \cdot uv = q_2 \cdot uv$ , alors  $q_1 \cdot u = q_2 \cdot u$ .

**Question 14.** Soit  $\eta : A^* \rightarrow M$  le monoïde syntactique de  $X^*$ . Montrer que si  $u$  est un mot de longueur  $\geq n$ , alors  $\eta(u)$  appartient à l'idéal minimal de  $M$ .

**Question 15.** Que peut-on en déduire sur la structure de  $M$ ?

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**Warning :** Clearness, accuracy and concision of the writing will be rewarded.

## Part 1. Syntactic monoid of $(abab)^*$ .

**Question 1.** Let  $A = \{a, b\}$  and let  $L = (abab)^*$ . Compute the minimal automaton of  $L$ .

**Question 2.** Compute the syntactic monoid  $M$  of  $L$ . Give the list of its elements (you should find 10 elements, including the identity) and of its defining relations.

**Question 3.** Give the set of idempotents of  $M$ .

**Question 4.** Compute the  $\mathcal{D}$ -class structure of  $M$  and draw the egg-box pictures.

**Question 5.** Is the language  $L$  star-free? Justify your answer.

**Question 6.** Show that  $L$  satisfies the profinite equations  $x^\omega y^\omega = y^\omega x^\omega$  for all  $x, y \in A^*$ .

## Part 2. Finite maximal prefix sets.

Recall that a prefix  $p$  of a word  $u$  is *proper* if  $p \neq u$ . A set  $X$  of nonempty words of  $A^+$  is *prefix* if no word of  $X$  is a proper prefix of another word of  $X$  (in other words, if  $u, uv \in X$  implies  $v = 1$ ). A prefix set  $X$  is *maximal* if, for all  $u \in A^+$ , the set  $X \cup \{u\}$  is not prefix.

The aim of this exercise is to study some properties of finite maximal prefix sets.

**Question 7.** Let  $X$  be a prefix set and let  $u \in X^*$ . Show that  $v \in X^*$  if and only if  $uv \in X^*$ .

**Question 8.** Show that the set

$$Z = \{aa, ab, baa, bab, bb\}$$

is a maximal prefix set.

Let  $X$  be a finite maximal prefix set and let  $P$  be the set of all proper prefixes of its words. For instance, if  $X = Z$ , then  $P = \{1, a, b, ba\}$ .

**Question 9.** Using the maximality of  $X$ , show that  $A^* = P \cup XA^*$  and that this union is a disjoint one.

**Question 10.** Deduce from the previous question that  $A^* = X^*P$  and that any word  $u \in A^*$  admits a unique factorization of the form  $u = x_1 \cdots x_n p$  with  $n \geq 0$ ,  $x_1, \dots, x_n \in X$  and  $p \in P$ .

**Question 11.** Let  $\mathcal{A}$  be the deterministic minimal automaton of  $X^*$ . Show that the initial state  $q_0$  of  $\mathcal{A}$  is also its unique final state.

Suppose now that  $X$  is a finite maximal prefix set which is also *suffix* (no word of  $X$  is a proper suffix of another word of  $X$ ). Let  $n$  be the maximal length of all words in  $X$ .

**Question 12.** Let  $q$  be a state of  $\mathcal{A}$  and let  $u$  be a word. Show that if  $q_0 \cdot u = q \cdot u$ , then  $q = q_0$ .

**Question 13.** Let  $u$  be a word of length  $\geq n$  and let  $q_1, q_2$  be two states of  $\mathcal{A}$ . Prove that if there exists a word  $v$  such that  $q_1 \cdot uv = q_2 \cdot uv$ , then  $q_1 \cdot u = q_2 \cdot u$ .

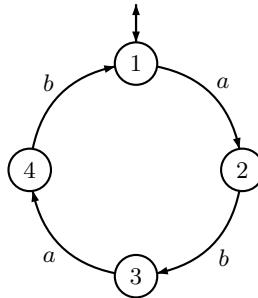
**Question 14.** Let  $\eta : A^* \rightarrow M$  be the syntactic monoid of  $X^*$ . Show that if  $u$  is a word of length  $\geq n$ , then  $\eta(u)$  belongs to the minimal ideal of  $M$ .

**Question 15.** What can be said about the structure of  $M$ ?

# Corrigé

## Part 1. Syntactic monoid of $(abab)^*$ .

**Question 1.** The minimal automaton of  $L$  is represented below



**Question 2.** The elements of  $M$  are given in the table below

	1	2	3	4
* 1	1	2	3	4
$a$	2	0	4	0
$b$	0	3	0	1
$* a^2$	0	0	0	0
$ab$	3	0	1	0
$ba$	0	4	0	2
$aba$	4	0	2	0
$bab$	0	1	0	3
$* abab$	1	0	3	0
$* baba$	0	2	0	4

This monoid has 10 elements. It has a zero since  $a^2 = 0$  and is defined by the relations  $a^2 = 0$ ,  $b^2 = 0$ ,  $ababa = a$  and  $babab = b$

**Question 3.** Idempotents:

$$E(S) = \{1, a^2, abab, baba\}$$

**Question 4.**  $\mathcal{D}$ -class structure:

$$\boxed{* 1}$$

$* baba$	$ba$	$bab$	$b$
$a$	$aba$	$* abab$	$ab$

$$\boxed{* a^2}$$

**Question 5.** The monoid  $M$  is not aperiodic, since the identity  $x^\omega = x^{\omega+1}$  is not satisfied for  $x = ab$ . Therefore  $L$  is not star-free.

**Question 6.** The idempotents of  $M$  commute and thus  $L$  satisfies the profinite identity  $x^\omega y^\omega = y^\omega x^\omega$ .

## Part 2. Finite maximal prefix sets.

**Question 7.** If  $v \in X^*$ , then clearly  $uv \in X^*$ . Since  $u \in X^*$ , one gets  $u = x_1 \cdots x_n$  for some  $x_1, \dots, x_n \in X$ . Suppose that  $uv \in X^*$ . Then  $uv = z_1 \cdots z_m$  for some  $z_1, \dots, z_m \in Z$ . Now, the words  $x_1$  and  $z_1$  are comparable in the prefix order. But they are both in  $X$ , which is a prefix set. Therefore  $x_1 = z_1$ . A similar argument would show that  $x_2 = z_2, \dots, x_n = z_n$ . Therefore  $v = z_{n+1} \cdots z_m \in X^*$ .

**Question 8.** Clearly, no word of  $Z$  is a proper prefix of another word of  $Z$ . Therefore  $Z$  is prefix. Suppose that  $Z \cup \{u\}$  is a prefix set. First suppose that  $|u| \geq 3$ . Since  $aa, ab, bb \in Z$ , the prefix of length 2 of  $u$  has to be  $ba$ . But then one of the words  $baa$  or  $bab$  is a prefix of  $u$ , a contradiction. Thus  $|u| \leq 2$ , but this is also impossible since all the words of length  $\leq 2$  are prefixes of some word of  $Z$ . Thus  $Z$  is maximal.

**Question 9.** Let  $u \in A^*$ . Since  $X \cup \{u\}$  is not a prefix set, either  $u$  is a prefix of some word of  $X$  or some word of  $X$  is a prefix of  $u$ . This gives the equality  $A^* = P \cup XA^*$ . The union is disjoint since if  $u \in P \cap XA^*$ , then  $u = xv$  for some  $x \in X$  and  $v \in A^*$ , which implies that  $x$  is a prefix of some word of  $X$ , a contradiction.

**Question 10.** Let us prove the result by induction on the length of  $u$ . The result is trivial if  $u \in P$  and in particular for  $u = 1$ . Suppose that  $u \notin P$ . Then since  $A^* = P \cup XA^*$ , one gets  $u \in XA^*$  and there exist  $x \in X$  and  $v \in A^*$  such that  $u = xv$ . Note that the pair  $(x, v)$  is unique since  $X$  is prefix. It suffices now to apply the induction hypothesis to  $v$  to conclude.

**Question 11.** Since  $1 \in X^*$ , the initial state is also a final state. Let  $q$  be a final state. Since  $q$  is accessible, there is a word  $u$  such that  $q_0 \cdot u = q$  and since  $q$  is final, one has  $u \in X^*$ . Now, by Question 7, the conditions  $v \in X^*$  and  $uv \in X^*$  are equivalent. It follows that the conditions  $q_0 \cdot v \in F$  and  $q \cdot v \in F$  are equivalent and thus the states  $q_0$  and  $q$  are equivalent. Since  $\mathcal{A}$  is minimal, they are equal. Therefore  $q_0$  is the unique final state.

**Question 12.** Since  $X$  is suffix, the dual version of Question 7 shows that the conditions  $u, vu \in X^*$  imply  $v \in X^*$ . Let  $q$  be a state. Since  $q$  is accessible and coaccessible, there are two words  $v$  and  $w$  such that  $q_0 \cdot v = q$  and  $q \cdot w = q_0$ . Now, if  $q_0 \cdot u = q \cdot u$ , then  $q_0 \cdot uw = q_0 \cdot vuw = q_0$ . It follows that  $uw, vuw \in X^*$  and thus  $v \in X^*$ . Therefore  $q_0 \cdot v = q_0$  and thus  $q = q_0$ .

**Question 13.** Suppose that  $q_1 \cdot uv = q_2 \cdot uv$ , where  $u$  is a word of length  $\geq n$ . Let  $w$  be a word such that  $q_0 \cdot w = q_1$ . Since  $A^* = X^*P$  by Question 10,  $wu = wrp$  for some words  $r$  and  $p$  such that  $rp = u$ ,  $wr \in X^*$  and  $p \in P$ . It follows that  $q_0 = q_0 \cdot wr = q_1 \cdot r$ . Since  $q_1 \cdot rpv = q_2 \cdot rpv$ , one gets  $q_0 \cdot pv = (q_2 \cdot r) \cdot pv$ . It follows by Question 12 that  $q_0 = q_2 \cdot r$ , whence  $q_1 \cdot r = q_2 \cdot r$  and finally  $q_1 \cdot u = q_2 \cdot u$ .

**Question 14.** It follows from Question 13 that all words of length  $\leq n$  have minimal rank in  $\mathcal{A}$ . Therefore their syntactic image belong to the minimal ideal of  $M$ .

**Question 15.** Consequently,  $M$  consists of the identity 1, some nonregular elements and a unique regular  $D$ -class, its minimal ideal.