

An introduction to Yoneda structures

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Catégories supérieures, polygraphes et homotopie

Paris 21 May 2010

Bibliography

Ross Street and Bob Walters
Yoneda structures on 2-categories
Journal of Algebra
50:350-379, 1978

Mark Weber
Yoneda structures from 2-toposes
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15:259-323, 2007

Covariant and contravariant presheaves

A few opening words on Isbell conjugacy

Ideal completion

Every partial order A generates a free complete \vee -lattice $\mathcal{P}A$

$$A \longrightarrow \mathcal{P}A$$

whose elements are the downward closed subsets of A , with

$$\varphi \leq_{\mathcal{P}A} \psi \iff \varphi \subseteq \psi.$$

$$\mathcal{P}A = A^{op} \Rightarrow \{0, 1\}$$

Free colimit completions of categories

Every small category \mathcal{A} generates a free cocomplete category $\mathcal{P}\mathcal{A}$

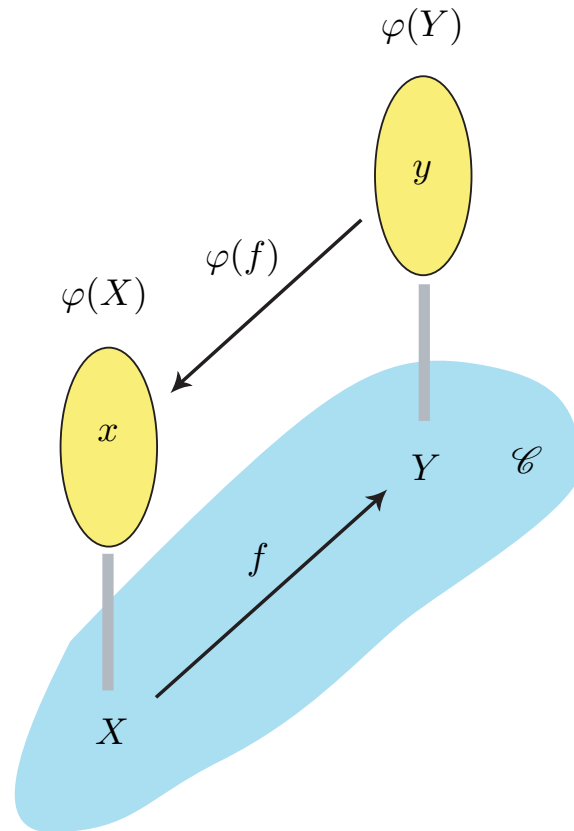
$$\mathcal{A} \longrightarrow \mathcal{P}\mathcal{A}$$

whose elements are the presheafs over \mathcal{A} , with

$$\varphi \longrightarrow \mathcal{P}\mathcal{A} \quad \psi \quad \Longleftrightarrow \quad \varphi \xrightarrow{\text{natural}} \psi.$$

$$\mathcal{P}\mathcal{A} = \mathcal{A}^{\text{op}} \Rightarrow \text{Set}$$

Contravariant presheaves



Replaces downward closed sets

Filter completion

Every partial order A generates a free complete \wedge -lattice $\mathcal{Q}A$

$$A \longrightarrow \mathcal{Q}A$$

whose elements are the upward closed subsets of A , with

$$\varphi \leq_{\mathcal{Q}A} \psi \iff \varphi \supseteq \psi.$$

$$\mathcal{Q}A = (A \Rightarrow \{0, 1\})^{op}$$

Free limit completions of categories

Every small category \mathcal{A} generates a free complete category $\mathcal{Q}\mathcal{A}$

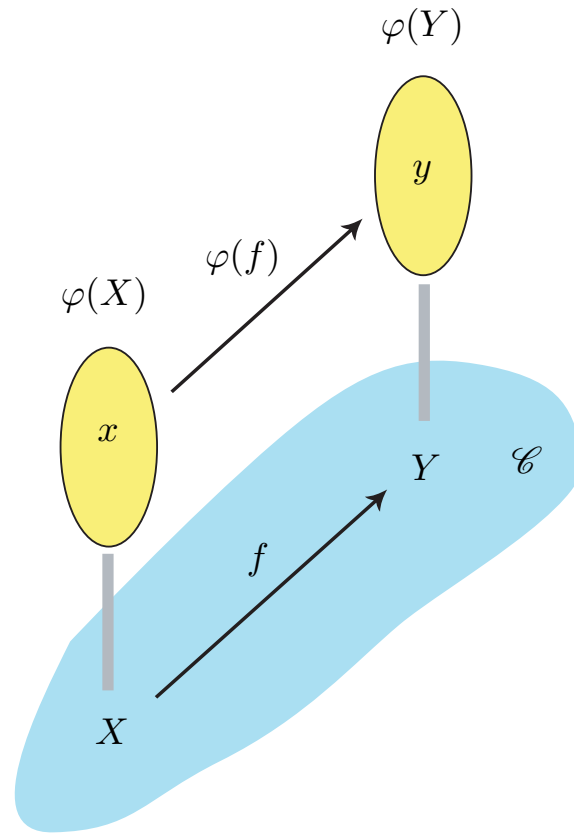
$$\mathcal{A} \longrightarrow \mathcal{Q}\mathcal{A}$$

whose elements are the covariant presheafs over \mathcal{A} , with

$$\varphi \longrightarrow_{\mathcal{Q}\mathcal{A}} \psi \iff \varphi \xleftarrow{\text{natural}} \psi.$$

$$\mathcal{Q}\mathcal{A} = (\mathcal{P}\mathcal{A}^{op})^{op} = (\mathcal{A} \Rightarrow \text{Set})^{op}$$

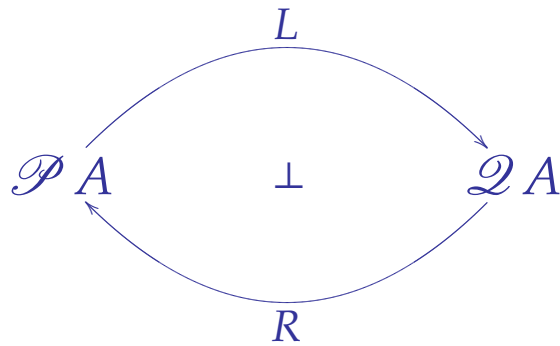
Covariant presheaves



Replaces upward closed sets

Related to the Dedekind-MacNeille completion

A Galois connection



$$L(\varphi) = \{ y \mid \forall x \in \varphi, x \leq_A y \}$$

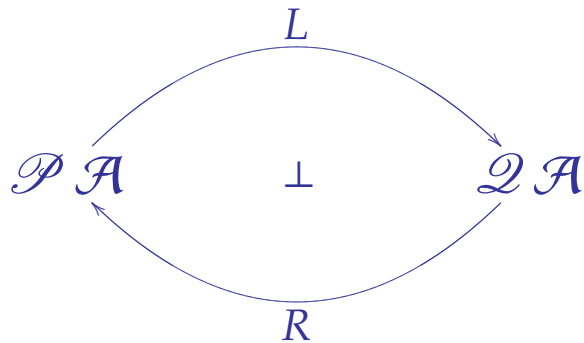
$$R(\psi) = \{ x \mid \forall y \in \psi, x \leq_A y \}$$

$$\varphi \subseteq R(\psi) \iff \forall x \in \varphi, y \in \psi, x \leq_A y \iff L(\varphi) \supseteq \psi$$

The completion keeps the pairs (φ, ψ) such that $\psi = L(\varphi)$ and $\varphi = R(\psi)$

The Isbell conjugacy

An adjunction



$$L(\varphi) : Y \mapsto \mathcal{P}\mathcal{A}(\varphi, Y) = \int_{x \in \mathcal{A}} \varphi(X) \Rightarrow \text{hom}(X, Y)$$

$$R(\psi) : X \mapsto \mathcal{Q}\mathcal{A}(X, \psi) = \int_{Y \in \mathcal{A}} \psi(Y) \Rightarrow \text{hom}(X, Y)$$

$$\mathcal{P}\mathcal{A}(\varphi, R(\psi)) \cong \int_{X, Y \in \mathcal{A}} \varphi(X) \times \psi(Y) \Rightarrow \text{hom}(X, Y) \cong \mathcal{Q}\mathcal{A}(L(\varphi), \psi)$$

Yoneda structures

An axiomatic approach by Street and Walters

General idea

Suppose a universe \mathcal{U}_1 inside a larger universe \mathcal{U}_2 .

- **Set** is the category of sets in the universe \mathcal{U}_1 ,
- \mathcal{CAT} is the 2-category of categories in the universe \mathcal{U}_2 .

In particular, the category **Set** is an object of \mathcal{CAT} .

Admissible functors

An object \mathcal{A} of \mathcal{CAI} is called **admissible** when its homsets

$$\mathcal{A}(A, A')$$

are all in the category **Set**. More generally, an arrow in \mathcal{CAI}

$$F : \mathcal{A} \longrightarrow \mathcal{B}$$

is called **admissible** when the homsets

$$\mathcal{B}(FA, B)$$

are all in the category **Set**.

Yoneda structures

A Yoneda structure in a 2-category \mathcal{K} is defined as

– a class of **admissible** arrows such that the composite arrow

$$\mathcal{A} \xrightarrow{F} \mathcal{B} \xrightarrow{G} \mathcal{C}$$

is admissible whenever the arrow

$$\mathcal{B} \xrightarrow{G} \mathcal{C}$$

is admissible.

Admissible objects

An object \mathcal{A} is called **admissible** when the identity arrow

$$id_{\mathcal{A}} : \mathcal{A} \longrightarrow \mathcal{A}$$

is admissible.

In the case of the 2-category \mathcal{CAT} equipped with its Yoneda structure:

admissible objects = locally small categories

Yoneda structures

– every admissible object

\mathcal{A}

induces an admissible arrow

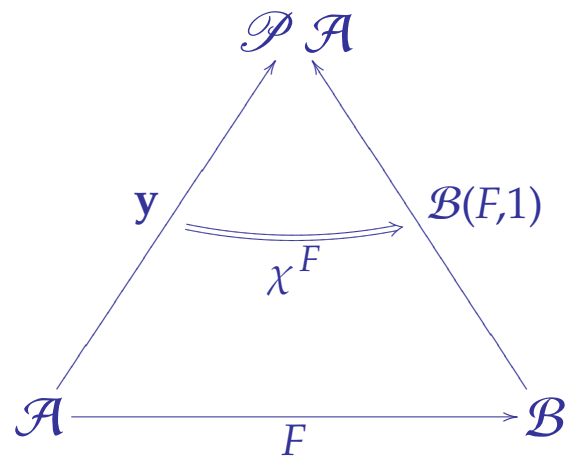
$$y_{\mathcal{A}} : \mathcal{A} \longrightarrow \mathcal{P}\mathcal{A}$$

Yoneda structures

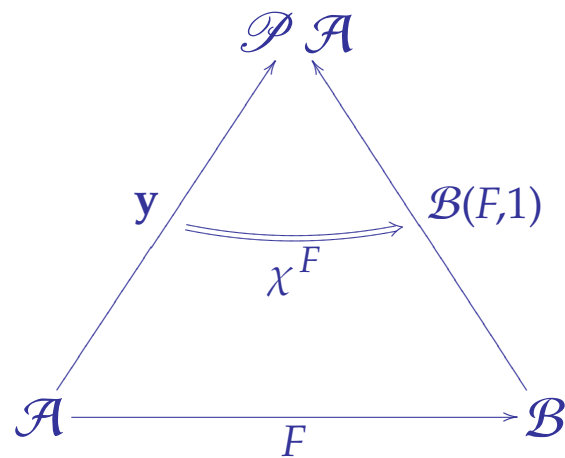
– every admissible arrow

$$F : \mathcal{A} \longrightarrow \mathcal{B}$$

from an admissible object \mathcal{A} induces a diagram



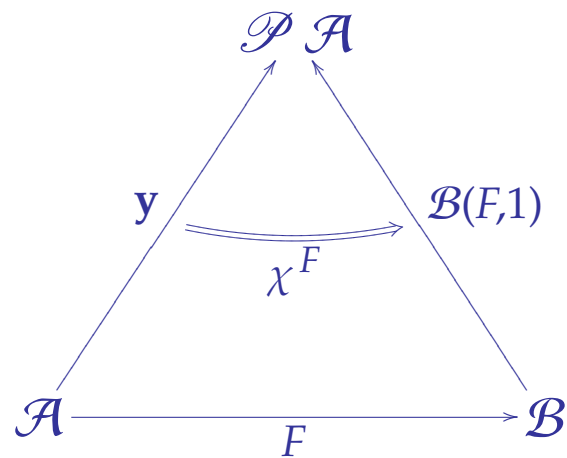
In the traditional case



$$\mathcal{B}(F,1) : b \mapsto \lambda a. \mathcal{B}(Fa, b)$$

$$\chi_{a_2}^F : \lambda a_1. \mathcal{A}(a_1, a_2) \longrightarrow \lambda a_1. \mathcal{B}(Fa_1, Fa_2)$$

In the traditional case

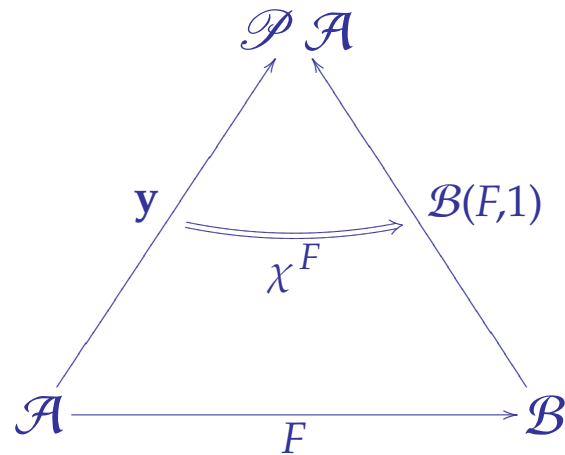


$$\mathcal{B}(F,1) : b \mapsto \lambda a. \mathcal{B}(Fa, b)$$

$$\chi_{a_1 a_2}^F : \mathcal{A}(a_1, a_2) \longrightarrow \mathcal{B}(Fa_1, Fa_2)$$

Axiom 1

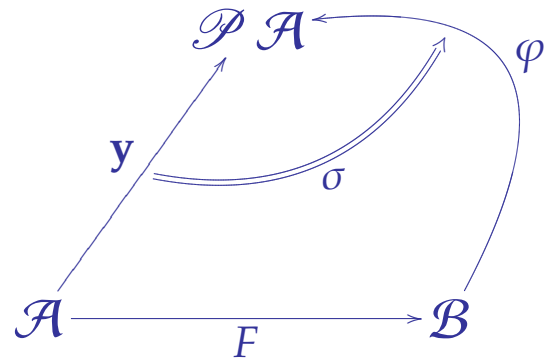
For \mathcal{A} and F accessible, the 2-cell



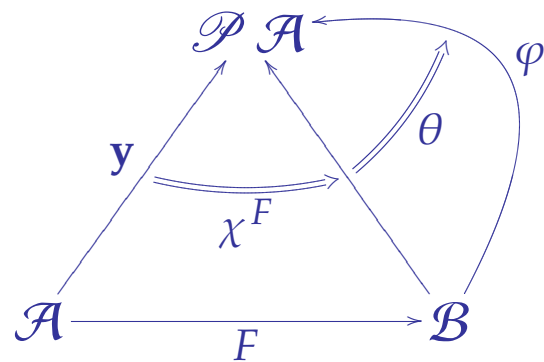
exhibits the arrow $\mathcal{B}(F,1)$ as a left extension of $y_{\mathcal{A}}$ along the arrow F .

Axiom 1

In other words, every 2-cell

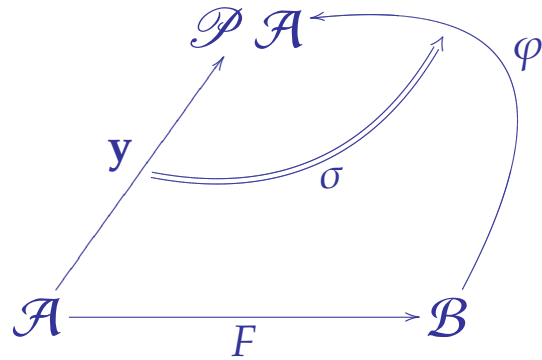


factors uniquely as



In the traditional case

A natural transformation



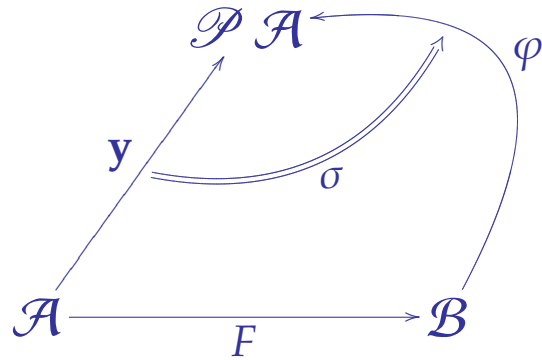
is the same thing as a family of functions

$$\sigma_{a_1 a_2} : \mathcal{A}(a_1, a_2) \longrightarrow \varphi(Fa_2)(a_1)$$

natural in a_1 contravariantly and in a_2 covariantly.

In the traditional case

A natural transformation



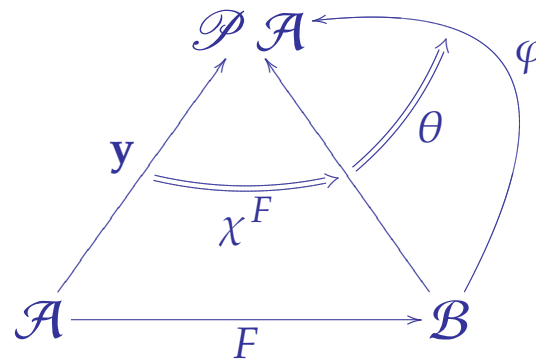
is the same thing as a family of functions

$$\sigma_{a_1 a_2} : \mathcal{A}(a_1, a_2) \longrightarrow \varphi(Fa_2)(a_1)$$

natural in a_1 contravariantly and in a_2 covariantly.

In the traditional case

A natural transformation θ in



is the same thing as a family of functions

$$\theta_{ab} : \mathcal{B}(Fa, b) \longrightarrow \varphi(b)(a)$$

natural in a contravariantly and in b covariantly.

In the traditional case

This means that every natural transformation

$$\sigma_{a_1 a_2} : \mathcal{A}(a_1, a_2) \longrightarrow \varphi(Fa_2)(a_1)$$

factors as

$$\chi_{a_1 a_2}^F : \mathcal{A}(a_1, a_2) \longrightarrow \mathcal{B}(Fa_1, Fa_2)$$

followed by

$$\theta_{a_1 Fa_2} : \mathcal{B}(Fa_1, Fa_2) \longrightarrow \varphi(Fa_2)(a_1)$$

for a unique natural transformation θ which remains to be defined.

Existence

Given the natural transformation σ , the natural transformation

$$\theta_{ab} : \mathcal{B}(Fa, b) \longrightarrow \varphi(Fa)(b)$$

transports every morphism

$$f : Fa \longrightarrow b$$

to the result of the action of f on the element

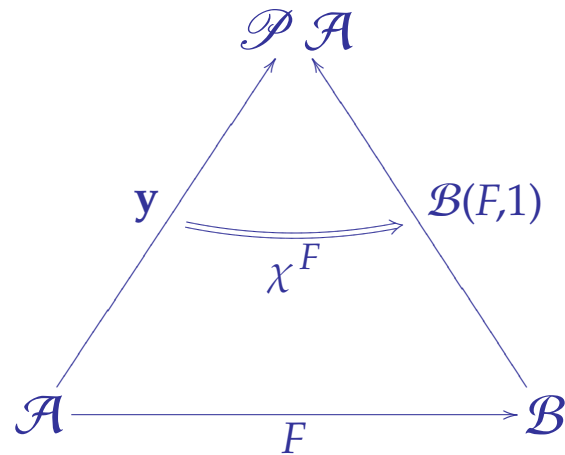
$$\sigma_{aa} \left(a \xrightarrow{id} a \right) \in \varphi(Fa)(a)$$

Uniqueness

...

Axiom 2

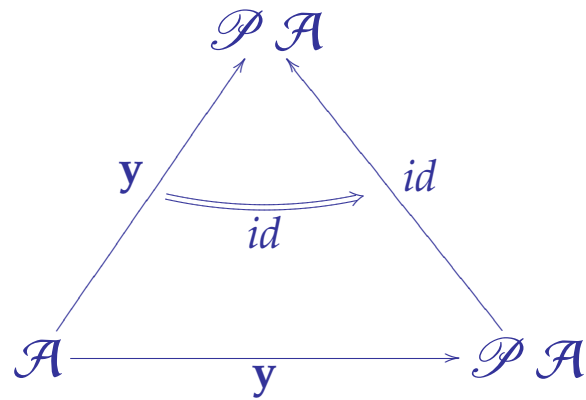
For \mathcal{A} and F accessible, the 2-cell



exhibits the arrow F as an absolute left lifting of $\mathbf{y}_{\mathcal{A}}$ through $\mathcal{B}(F, 1)$.

Axiom 3(i)

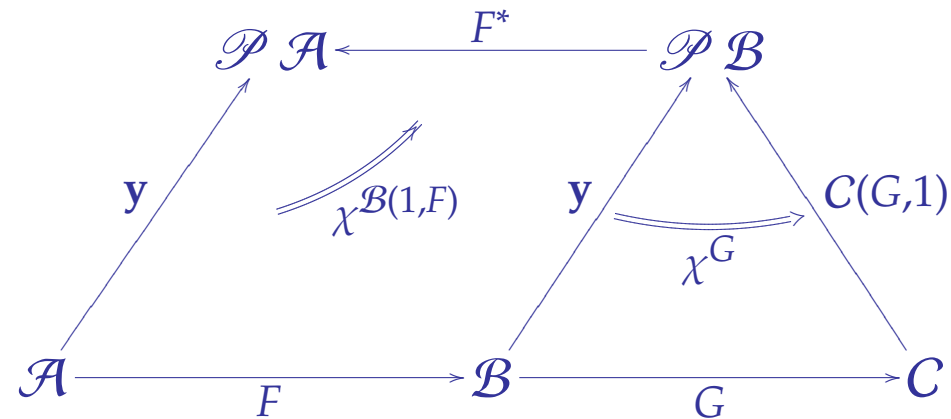
For \mathcal{A} accessible, the identity 2-cell



exhibits the identity arrow as a left extension of $\mathbf{y}_{\mathcal{A}}$ along $\mathbf{y}_{\mathcal{A}}$.

Axiom 3(ii)

For $\mathcal{A}, \mathcal{B}, F, G$ accessible, the 2-cell



exhibits the arrow $F^* \circ C(G,1)$ as a left extension of $y_{\mathcal{A}}$ along $G \circ F$.

The inverse arrow

Given an arrow between admissible objects \mathcal{A} and \mathcal{B}

$$F : \mathcal{A} \longrightarrow \mathcal{B}$$

the arrow

$$F^* : \mathcal{P}\mathcal{B} \longrightarrow \mathcal{P}\mathcal{A}$$

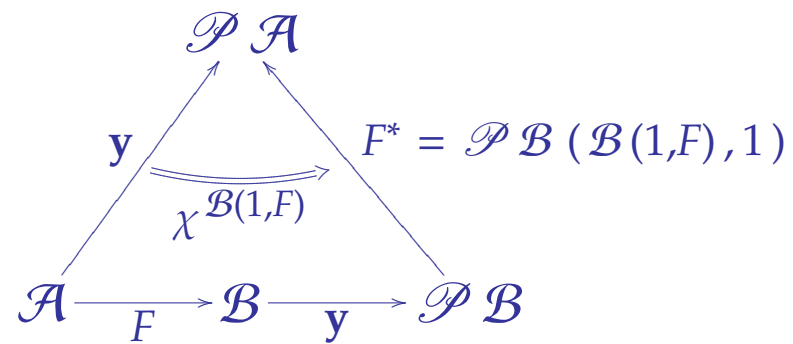
is defined as follows:

$$\begin{array}{ccccc}
 & & \mathcal{P}\mathcal{A} & & \\
 & \nearrow \mathbf{y} & & \nwarrow F^* = \mathcal{P}\mathcal{B}(\mathcal{B}(1,F), 1) & \\
 & \xrightarrow{\chi^{\mathcal{B}(1,F)}} & & & \\
 \mathcal{A} & \xrightarrow{F} & \mathcal{B} & \xrightarrow{\mathbf{y}} & \mathcal{P}\mathcal{B}
 \end{array}$$

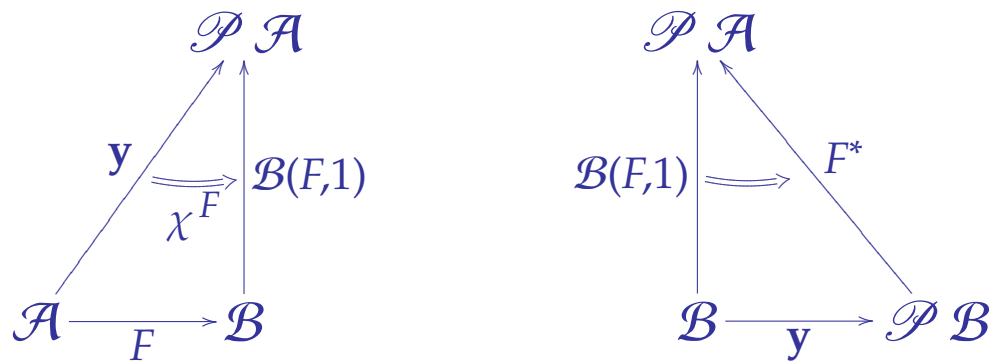
where $\mathcal{B}(1,F)$ denotes the composite arrow $\mathbf{y} \circ F$.

The inverse arrow

The 2-dimensional cell



factors as a pair of Kan extensions:



Existential image

Given an arrow between admissible objects \mathcal{A} and \mathcal{B}

$$F : \mathcal{A} \longrightarrow \mathcal{B}$$

the arrow

$$\exists_F : \mathcal{P}\mathcal{A} \longrightarrow \mathcal{P}\mathcal{B}$$

is defined as follows:

$$\begin{array}{ccc} & \mathcal{P}\mathcal{A} & \xrightarrow{\exists_F} & \mathcal{P}\mathcal{B} \\ & \uparrow y & & \uparrow y \\ \mathcal{A} & \xrightarrow{F} & \mathcal{B} & \end{array}$$

Universal image

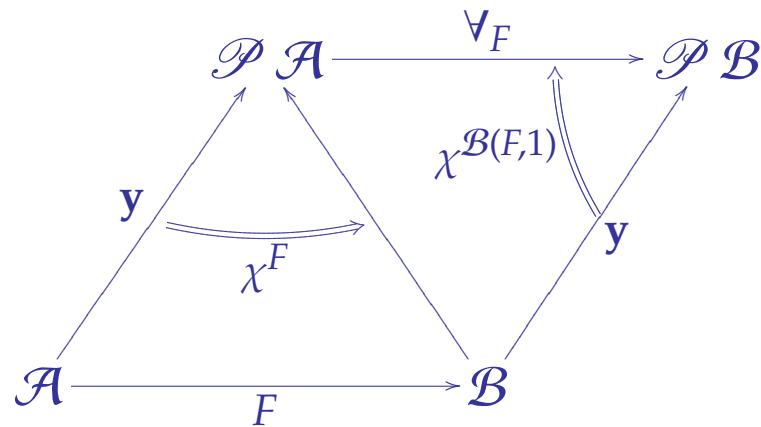
Given an arrow between admissible objects \mathcal{A} and \mathcal{B}

$$F : \mathcal{A} \longrightarrow \mathcal{B}$$

the arrow

$$\forall_F : \mathcal{P}\mathcal{A} \longrightarrow \mathcal{P}\mathcal{B}$$

is defined as follows:



Monads with arity

An idea by Mark Weber

Category with arity

A fully faithful and dense functor

$$i_0 : \Theta_0 \longrightarrow \mathcal{A}$$

where Θ_0 is a small category.

Category with arity

This induces a fully faithful functor

$$\mathcal{A}(i_0, 1) : \mathcal{A} \longrightarrow \mathcal{P}\Theta_0$$

which transports every object A of the category \mathcal{A} into the presheaf

$$\begin{array}{l} \mathcal{A}(i_0, A) : \Theta_0 \longrightarrow \mathbf{Set} \\ p \longmapsto \mathcal{A}(i_0 p, A) \end{array}$$

Monads with arity

A monad T on a category with arity (\mathcal{A}, i_0) such that:

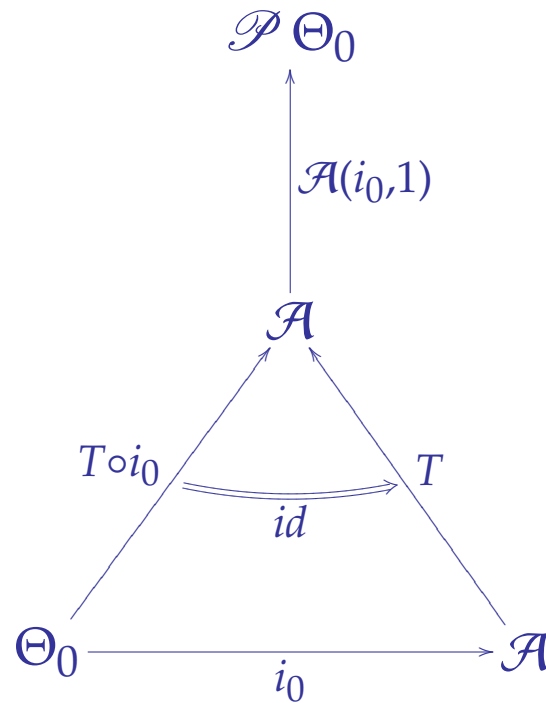
(1) the natural transformation exhibits the functor T

The diagram is a triangle with vertices Θ_0 at the bottom-left, \mathcal{A} at the bottom-right, and \mathcal{A} at the top. The bottom edge is a horizontal arrow from Θ_0 to \mathcal{A} labeled i_0 . The left edge is an arrow from Θ_0 to the top \mathcal{A} labeled $T \circ i_0$. The right edge is an arrow from the top \mathcal{A} to the bottom-right \mathcal{A} labeled T . A curved arrow labeled id points from the right edge T to the left edge $T \circ i_0$, indicating a natural transformation.

as a left kan extension of the functor $T \circ i_0$ along the functor i_0 ,

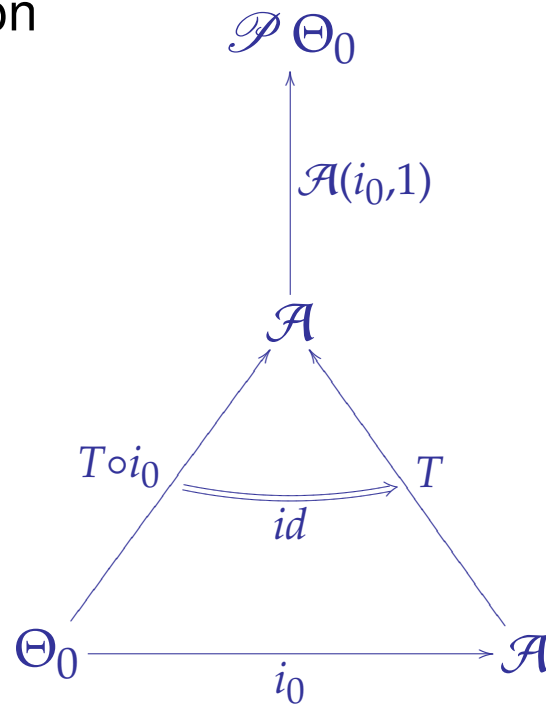
Monads with arity

(2) this left kan extension is preserved by the inclusion functor to $\mathcal{P} \Theta_0$.



Equivalently...

The natural transformation



exhibits $\mathcal{A}(i_0, 1) \circ T$ as a left kan extension of $\mathcal{A}(i_0, 1) \circ T \circ i_0$ along i_0 .

Equivalently...

For every object A , the canonical morphism

$$\int^{p \in \Theta_0} \mathcal{A}(i_0 n, T i_0 p) \times \mathcal{A}(i_0 p, A) \longrightarrow \mathcal{A}(i_0 n, T A)$$

is an isomorphism.

Unique factorization up to zig-zag

Every morphism

$$i_0 n \longrightarrow TA$$

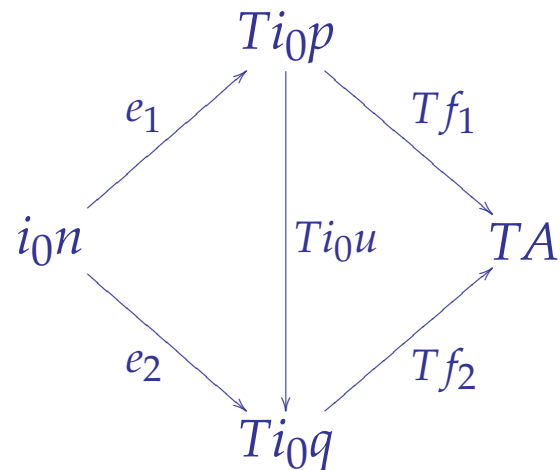
in the category \mathcal{A} decomposes as

$$i_0 n \xrightarrow{e} Ti_0 p \xrightarrow{Tf} TA$$

for a pair of morphisms $e : i_0 n \rightarrow Ti_0 p$ and $f : i_0 p \rightarrow A$.

Unique factorization up to zig-zag

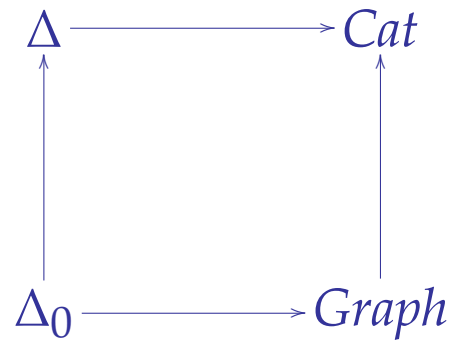
The factorization should be unique up to zig-zag of



An abstract Segal condition

A general theorem by Mark Weber
axiomatizing a theorem by Clemens Berger
on higher dimensional categories

Motivating example: the free category monad



Morphism between categories with arity

A morphism between categories with arity

$$(F, \ell) : (\mathcal{A}, i_0) \longrightarrow (\mathcal{B}, i_1)$$

is defined as a pair of functors (F, ℓ) making the diagram

$$\begin{array}{ccc} \Theta_1 & \xrightarrow{i_1} & \mathcal{B} \\ \ell \uparrow & & \uparrow F \\ \Theta_0 & \xrightarrow{i_0} & \mathcal{A} \end{array}$$

commute.

Morphism between categories with arity

This induces a commutative diagram

$$\begin{array}{ccc} \mathcal{P} \Theta_1 & \xleftarrow{i_1^*} & \mathcal{P} \mathcal{B} \\ \downarrow \ell^* & \swarrow id & \downarrow F^* \\ \mathcal{P} \Theta_0 & \xleftarrow{i_0^*} & \mathcal{P} \mathcal{A} \end{array}$$

which is required to be an exact square in the sense of Guitart.

Morphism between categories with arity

This means that the Beck-Chevalley condition holds, which states that the canonical natural transformation

$$\begin{array}{ccc} \mathcal{P} \Theta_1 & \xrightarrow{\forall_{i_1}} & \mathcal{P} \mathcal{B} \\ \ell^* \downarrow & \curvearrowright & \downarrow F^* \\ \mathcal{P} \Theta_0 & \xrightarrow{\forall_{i_0}} & \mathcal{P} \mathcal{A} \end{array}$$

is reversible.

Proposition A.

For every morphism (F, ℓ) between categories with arity

$$(F, \ell) : (\mathcal{A}, i_0) \longrightarrow (\mathcal{B}, i_1)$$

the adjunction $i_1^* \dashv \mathbb{V}_{i_1}$ induces an adjunction between

- the full subcategory of presheaves of \mathcal{B} whose restriction along F is representable in \mathcal{A} ,
- the full subcategory of presheaves of Θ_1 whose restriction along ℓ is representable along i_0 .

Moreover, this adjunction defines an equivalence when the functor F is essentially surjective.

Proposition B.

Every monad T with arity i_0 induces a commutative diagram

$$\begin{array}{ccc} \Theta_T & \xrightarrow{i_T} & \mathcal{A}_T \\ \ell \uparrow & & \uparrow F \\ \Theta_0 & \xrightarrow{i_0} & \mathcal{A} \end{array}$$

where the pair (F, ℓ) defines a morphism

$$(F, \ell) : (\mathcal{A}, i_0) \longrightarrow (\mathcal{A}_T, i_T)$$

of categories with arity.

Algebraic theories with arity

A 2-dimensional approach to Lawvere theories

Algebraic theory with arity

An algebraic theory \mathbb{L} with arity

$$i_0 : \Theta_0 \longrightarrow \mathcal{A}$$

is an identity-on-object functor

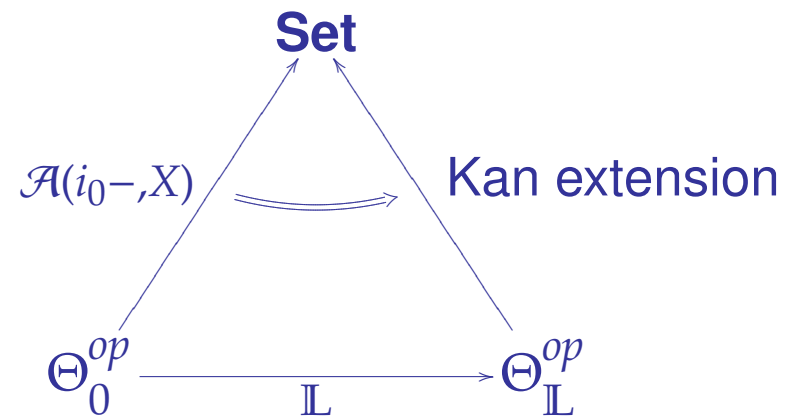
$$\mathbb{L} : \Theta_0 \longrightarrow \Theta_{\mathbb{L}}$$

such that the endofunctor

$$\mathcal{P} \Theta_0 \xrightarrow{\exists_{\mathbb{L}}} \mathcal{P} \Theta_{\mathbb{L}} \xrightarrow{\mathbb{L}^*} \mathcal{P} \Theta_0$$

maps a presheaf representable along i_0 to a presheaf representable along i_0 .

Algebraic theories with arity



Model of an algebraic theory

A model A of a Lawvere theory \mathbb{L} is a presheaf

$$A : \Theta_{\mathbb{L}}^{op} \longrightarrow \mathbf{Set}$$

such that the induced presheaf

$$\Theta_0^{op} \xrightarrow{i_0} \Theta_{\mathbb{L}}^{op} \xrightarrow{A} \mathbf{Set}$$

is representable along i_0 .

Main theorem

the category of algebraic theories with arity (\mathcal{A}, i_0)

is equivalent to

the category of monads with arity (\mathcal{A}, i_0)