# On the existence of Trott numbers relative to multiple bases

Pieter C. Allaart

University of North Texas (allaart@unt.edu)

One World Numeration Seminar November 2, 2021

(Joint work with S. Jackson, T. Jones and D. Lambert)

Preprint available at arxiv.org/abs/2108.03664

#### Question:

Is it possible for a number to have its continued fraction expansion agree with its decimal expansion? That is,

$$[0; a_1, a_2, a_3 \dots] = 0.a_1 a_2 a_3 \dots$$

If we require each  $a_i \in \{1, 2, \dots, 9\}$ , it is easy to see that the answer is no.

## Trott's Constant

In 2006, M. Trott published the following example in a blog post:



This number is now known as *Trott's constant* - see OEIS sequence A039662.

## Trott numbers

The use of the digit 0 is, of course, unnatural from a dynamical point of view. So we slightly tweak the definition:

#### Definition

Let  $b \in \mathbb{N}_{\geq 2}$ . A number  $x \in (0, 1)$  is a Trott number in base b if

$$x = [0; a_1, a_2, a_3, \dots] = (0.\hat{a}_1 \hat{a}_2 \hat{a}_3 \dots)_b,$$

where  $\hat{a}_i$  is the string of digits corresponding to the base b representation of  $a_i$ .

For example, in base 10,

$$[0; 3, 29, 545, 6, \ldots] = 0.3\,29\,545\,6\ldots$$

## Question: For which bases, if any, do Trott numbers exist? And how can we construct them?

## Two initial observations

• If x is a Trott number, that is,

$$x = [0; a_1, a_2, a_3, \dots] = (0.\hat{a}_1 \hat{a}_2 \hat{a}_3 \dots)_b,$$

then x is neither rational nor quadratic irrational.

• If  $x = [0; a_1, a_2, a_3, \dots]$  is Trott in base b, then  $a_1 = \lfloor \sqrt{b} \rfloor$ .

Proof. The intervals

$$\left[\frac{a_1}{b}, \frac{a_1+1}{b}\right] \quad \text{and} \quad \left[\frac{1}{a_1+1}, \frac{1}{a_1}\right]$$

must overlap.

This shows also that no Trott numbers exist in base b if b is a perfect square: If  $b = k^2$ , then  $a_1 = k$  and the two intervals only share an endpoint.

Theorem (A, Jackson, Jones & Lambert, 2021)

There exists a Trott number in base b if and only if

$$b \in \{3\} \cup \bigcup_{k=1}^{\infty} \{k^2 + 1, k^2 + 2, \dots, k^2 + k\} =: \Gamma.$$

Furthermore, if  $T_b$  is nonempty, then it is uncountable.

So the first few "good" bases are  $2,3,5,6,10,11,12,17,18,19,20, \ldots$ 

The proof of sufficiency is rather involved. See the slides of the talk by T. Jones on October 5, 2021 for the case b = 10. The construction is inductive, and the key is to choose  $a_2$  carefully!

## Proof of necessity

To prove necessity, suppose  $b \notin \Gamma$  and there is a Trott number

$$x = [0; a_1, a_2, a_3, \dots] = (0.\hat{a}_1 \hat{a}_2 \hat{a}_3 \dots)_b$$

Then there is a unique  $k \in \mathbb{N}$ ,  $k \ge 2$  such that

$$k^2 + k < b < (k+1)^2.$$

By our earlier observation,  $a_1 = \lfloor \sqrt{b} \rfloor = k$ . Suppose  $a_2 = j$ . Then

$$x \in \left[\frac{1}{k+\frac{1}{j}}, \frac{1}{k+\frac{1}{j+1}}\right]$$

Case 1. If j > (b-1)/k, then

$$\frac{1}{k+\frac{1}{j}} > \frac{b-1}{bk} \ge \frac{k+1}{b}$$

so x does not begin with digit k, contradiction.

## Proof of necessity (ctd)

Recall that  $a_1 = k$ ,  $a_2 = j$  and

$$x \in \left[\frac{1}{k+\frac{1}{j}}, \frac{1}{k+\frac{1}{j+1}}\right]$$

٠

Case 2. If  $j \leq (b-1)/k$ , then j < b and it can be checked easily that

$$\frac{1}{k + \frac{1}{j}} > \frac{k}{b} + \frac{j + 1}{b^2},$$

so the second digit of x is at least j + 1. Contradiction again!

## Some other results

Let  $T_b$  denote the set of Trott numbers in base b.

Theorem (A, Jackson, Jones & Lambert, 2021)

For each  $b \in \Gamma$ ,  $T_b$  is a complete  $G_\delta$  set. (That is,  $T_b$  is  $G_\delta$  but not  $F_{\sigma}$ .)

Theorem (A, Jackson, Jones & Lambert, 2021)

Let

$$T := \bigcup_{n=2}^{\infty} T_b$$

be the set of numbers that are Trott in some base. Then T is nowhere dense, and  $\dim_H T < 1$ .

(In fact we have explicit upper bounds for  $\dim_H T_b$  for each b, but they are almost surely very bad!)

#### Conjecture 1

 $\dim_H T_b > 0$  for all  $b \in \Gamma$ .

Conjecture 2

 $\dim_H T_b \to 0 \text{ as } b \to \infty.$ 

#### Conjecture 3

 $T_b \cap T_c = \emptyset$  whenever  $b \neq c$ .

We'll focus here on Conjecture 3, which says that no number is Trott in more than one base.

A (1) > A (2) > A

## Conjecture 3:

 $T_b \cap T_c = \emptyset$  whenever  $b \neq c$ .

æ

- 4 目 ト - 日 ト - 4

Let  $2 \leq b < c$ , and suppose

$$x = [0; a_1, a_2, \dots] \in T_b \cap T_c.$$

Since  $a_1 = \lfloor \sqrt{b} \rfloor = \lfloor \sqrt{c} \rfloor$ , we see at once that b and c must belong to the same interval  $[k^2 + 1, k^2 + k]$ .

From here, we tried (with only partial success) to obtain a contradiction by using only information about  $a_2$ .

### Two lemmas

#### Lemma 1

Let b < c. If  $a_2$  has l digits in base b and m digits in base c, then

 $2 \le m < l.$ 

Lemma 2

Let  $k \ge 2$  and  $k^2 + 1 \le b \le k^2 + k$ . Assume  $b \ge 6$ . Suppose  $x = [0; a_1, a_2, \ldots] \in T_b$  and  $a_2$  has  $l \ge 2$  digits in base b. Then

$$a_2 = \left\lceil \frac{b^l(b-k^2)}{k} - \frac{b}{k(b-k^2)} \right\rceil - 1.$$

*Proof.* Follows (after some straightforward but cumbersome algebra) from the fact that  $a_1 = k$ , so these intervals overlap:

$$\left[\frac{k}{b} + \frac{a_2}{b^{l+1}}, \frac{k}{b} + \frac{a_2 + 1}{b^{l+1}}\right] \quad \text{and} \quad \left[\frac{a_2}{ka_2 + 1}, \frac{a_2 + 1}{k(a_2 + 1) + 1}\right].$$

## Picking the low-hanging fruit

#### Corollary (Two special cases)

Let  $a_2$  have l digits in base b, where  $l \ge 2$ . (i) If  $b = k^2 + 1$  and  $k \ge 3$ , then

$$a_2 = \frac{(k^2+1)^l - (k^2+1)}{k} - 1.$$

(ii) If 
$$b = k^2 + k$$
 and  $k \ge 2$ , then  $a_2 = b^l - 2$ .

#### Proposition

No number is Trott in both base  $k^2 + 1$  and  $k^2 + k$ .

Proof. By the Corollary, we would have

$$\frac{(k^2+1)^l - (k^2+1)}{k} - 1 = (k^2+k)^m - 2.$$

This is impossible!

## Specific base pairs: using modular arithmetic

**Example.** Let  $b = 17 = 4^2 + 1$ ,  $c = 19 = 4^2 + 3$  (two "good" bases!). Let  $a_2$  have l digits in base b and m digits in base c. By the Corollary and Lemma 2,

$$a_2 = \frac{17^l - 21}{4} = \begin{cases} \frac{3 \cdot 19^m - 9}{4} & \text{if } m \text{ is odd}, \\ \frac{3 \cdot 19^m - 7}{4} & \text{if } m \text{ is even}, \end{cases}$$

Thus,

$$17^l - 3 \cdot 19^m = \begin{cases} 12 & \text{if } m \text{ is odd}, \\ 14 & \text{if } m \text{ is even}. \end{cases}$$

We can instantly rule out 12. This leaves us with

$$17^l - 3 \cdot 19^m = 14.$$

Does the equation

$$17^l - 3 \cdot 19^m = 14$$

have a solution (l,m) in  $(\mathbb{Z}_+ \times \mathbb{Z}_+)$ ? Consider first the equation modulo 6:

$$(-1)^l - 3 \cdot 1^m \equiv 2 \pmod{6}.$$

This shows l must be odd. Now consider the equation modulo 19:

$$(-2)^l \equiv -5 \pmod{19}.$$

Since l is odd, this becomes  $2^l\equiv 5 \pmod{19}$ . However, the odd powers of 2 modulo 19 are

 $2, 8, 13, 14, 18, 15, 3, 12, 10, 2, 8, 13, \ldots$ 

Hence,  $T_{17} \cap T_{19} = \emptyset$ .

### What the lemmas tell us

#### Recall:

#### Lemma 2

Let  $k \ge 2$  and  $k^2 + 1 \le b \le k^2 + k$ . Assume  $b \ge 6$ . Suppose  $x = [0; a_1, a_2, \ldots] \in T_b$  and  $a_2$  has  $l \ge 2$  digits in base b. Then

$$a_2 = \left\lceil \frac{b^l(b-k^2)}{k} - \frac{b}{k(b-k^2)} \right\rceil - 1.$$

Now let  $k \ge 3$  and  $k^2 + 1 \le b < c \le k^2 + k$ . Suppose  $a_2$  has l digits in base b and m digits in base c. By Lemmas 1 and 2,  $2 \le m < l$  and

$$\left\lceil \frac{b^l(b-k^2)}{k} - \frac{b}{k(b-k^2)} \right\rceil = \left\lceil \frac{c^m(c-k^2)}{k} - \frac{c}{k(c-k^2)} \right\rceil$$

## Two tight inequalities

We just saw that

$$\left\lceil \frac{b^l(b-k^2)}{k} - \frac{b}{k(b-k^2)} \right\rceil = \left\lceil \frac{c^m(c-k^2)}{k} - \frac{c}{k(c-k^2)} \right\rceil.$$

Thus,

$$\left|\frac{b^l(b-k^2)}{k} - \frac{b}{k(b-k^2)} - \left(\frac{c^m(c-k^2)}{k} - \frac{c}{k(c-k^2)}\right)\right| < 1.$$

Some simplification leads to

$$\frac{k^2(c-b)}{(b-k^2)(c-k^2)} - k < b^l(b-k^2) - c^m(c-k^2) < \frac{k^2(c-b)}{(b-k^2)(c-k^2)} + k,$$

and in particular,

$$\left| b^{l}(b-k^{2}) - c^{m}(c-k^{2}) \right| < k^{2}.$$

Recall  $k^2 + 1 \le b < c \le k^2 + k$ . From the inequality

$$\left| b^l (b - k^2) - c^m (c - k^2) \right| < k^2$$

we can deduce (after minor technicalities) that

 $l>m>k\log k.$ 

To make further progress, we need "Baker's theorem", really a collection of theorems (going back to the 1970s) giving lower bounds for expressions of the form

$$\Lambda := \beta_1 \log \alpha_1 + \beta_2 \log \alpha_2 + \dots + \beta_n \log \alpha_n$$

provided  $\Lambda \neq 0$ , where  $\alpha_1, \ldots, \alpha_n$  are algebraic numbers and  $\beta_1, \ldots, \beta_n$  are rational integers.

## Baker's theorem for our setting

The version of Baker's theorem that we'll use (stated in simplified form for our setting) is:

#### Lemma (Matveev, 2000)

Let  $\alpha_1, \ldots, \alpha_n \in \mathbb{Q}$ , not 0 or 1, and  $\beta_1, \ldots, \beta_n$  be integers. Let

$$\Lambda := \beta_1 \log \alpha_1 + \beta_2 \log \alpha_2 + \dots + \beta_n \log \alpha_n \neq 0.$$

#### Then

$$|\Lambda| \ge \exp\{-C_n h(\alpha_1) \cdots h(\alpha_n) \log(eB)\},\$$

where  $B := \max\{|\beta_1|, ..., |\beta_n|\}$ ,

$$C_n := \min\left\{\frac{e}{2} \cdot 30^{n+3} n^{4.5}, 2^{6n+20}\right\},$$

and for  $\alpha = p/q$  in lowest terms,  $h(\alpha) := \log \max\{|p|, |q|\}$  is the (logarithmic) Weil height of  $\alpha$ .

< /i>
< /i>
< /i>
< /i>
< /i>

## Applying Matveev's lemma

It's not too hard to show that  $b^l(b-k^2)-c^m(c-k^2)\neq 0$ . Assume WLOG that  $c^m(c-k^2)>b^l(b-k^2)$ . Let

$$\Lambda := m \log c - l \log b + \log \left(\frac{c - k^2}{b - k^2}\right),$$

so  $\Lambda > 0$ . Applying Matveev's lemma with  $\alpha_1 = c$ ,  $\alpha_2 = b$ ,  $\alpha_3 = (c - k^2)/(b - k^2)$ ,  $\beta_1 = m$ ,  $\beta_2 = -l$  and  $\beta_3 = 1$  we obtain

$$\begin{aligned} \frac{c^m(c-k^2)}{b^l(b-k^2)} - 1 &= e^{\Lambda} - 1 > \Lambda \\ &\geq \exp\{-C_3 \log b \log c \log(c-k^2) \log(el)\} \\ &\geq \exp\{-C_3 \left(\log(k^2+k)\right)^2 \log k \log(el)\} \\ &\approx \exp\{-4C_3 (\log k)^3 \log l\} \\ &=: A_{k,l}. \end{aligned}$$

## Applying Matveev's lemma (ctd)

Recall that 
$$\left| b^l(b-k^2) - c^m(c-k^2) \right| < k^2$$
. Hence,

$$k^{2} > c^{m}(c - k^{2}) - b^{l}(b - k^{2}) \gtrsim A_{k,l}b^{l}(b - k^{2})$$
$$\geq A_{k,l}(k^{2} + 1)^{l} > A_{k,l}k^{2l}.$$

Thus,

$$k^{2l-2} \lesssim A_{k,l}^{-1} = \exp\{4C_3(\log k)^3 \log l\}.$$

Taking logs,

$$2(l-1)\log k \lesssim 4C_3(\log k)^3\log l.$$

Since  $l > m > k \log k$ , we deduce that

$$k(\log k)^2 \lesssim 2C_3(\log k)^3 \log(k \log k),$$

or  $k \lesssim 2C_3 \log k \log(k \log k)$ . Since  $C_3 \approx 1.39 \times 10^{11}$ , this yields

$$k \lesssim 3.44 \times 10^{14}, \qquad b, c \le k^2 + k \lesssim 1.185 \times 10^{29},$$

The conjecture is now reduced to pairs (b,c) of bases below  $1.185\times 10^{29},$  finitely many!

If b and c are below this threshold (and hence  $k \leq 3.44 \times 10^{14}$ ), then the inequality

$$2(l-1)\log k \lesssim 4C_3(\log k)^3\log l$$

gives an implicit upper bound for l, and since m < l, there are only finitely many pairs (l,m) to try.

This reduces the verification of the conjecture to a finite number of operations.

Waiting for quantum computers...

#### Theorem

Let  $b, c \in \Gamma$  with b < c. Then  $T_b \cap T_c = \emptyset$  if at least one of the following holds: (i)  $|\sqrt{b}| \neq |\sqrt{c}|$ ; (ii) gcd(b,c) > 1;(iii) c = b + 1;(iv) c = b + 2 and there exists k > 3 such that  $k^{2} < b \le k^{2} + \sqrt{\frac{2k^{2}}{k-2}} + 1 - 1 \approx k^{2} + \sqrt{2k};$ There exists  $k \ge 2$  such that  $b = k^2 + 1$  and  $c = k^2 + k$ ; (v) (vi)  $b > 1.185 \times 10^{29}$ .

## Thank you!

æ