# Adding machine, Automata and Julia sets

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**Objective**: study dynamical, spectral and probabilistic properties of stochastic adding machines. This work involves

Julia sets

Automata

Topological Dynamics on Banach spaces

Recurrence and transience of infinite state Markov chains.

#### Killeen and Taylor Stochastic adding machine

Let 
$$N \in \mathbb{N} = \{0, 1, 2, 3, ...\}.$$
  
$$N = \sum_{i=0}^{k(N)} \varepsilon_i(N) 2^i = \varepsilon_{k(N)} \dots \varepsilon_{0(N)}$$

with  $\varepsilon_i(N) = 0, 1$ .

**Example:**  $6 = 1 \cdot 2^2 + 1 \cdot 2^1 + 0 \cdot 2^0 = 110_2$ .

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Let  $N = \varepsilon_k \dots \varepsilon_1 \varepsilon_0 \in \mathbb{N}$ . The digits of N + 1 in base 2 are given by the following algorithm

If  $\varepsilon_0 = 0$ , then:

$$N = \begin{bmatrix} \varepsilon_k & \dots & \varepsilon_1 & 0 \\ + & & 1 \end{bmatrix}$$
$$N+1 = \begin{bmatrix} \varepsilon_k & \dots & \varepsilon_1 & 1 \end{bmatrix}$$

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If  $\varepsilon_{i-1} \dots \varepsilon_0 = 1 \dots 1$  and  $\varepsilon_i = 0$ , then:

$$N = \begin{bmatrix} \varepsilon_k & \dots & \varepsilon_{i+1} & {}^10 & {}^11 & \dots & {}^11 & 1 \\ + & & & 1 \\ \hline N+1 = \begin{bmatrix} \varepsilon_k & \dots & \varepsilon_{i+1} & 1 & 0 & \dots & 0 & 0 \end{bmatrix}$$

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0 transitions to 0 with probability 1 - p (if we don't add the carry): S(0,0) = 1 - p.

0 transitions to 1 with probability p (if we add the carry): S(0,1) = p.

1 transitions to 1 with probability 1 - p (if we don't add the carry): S(1,1) = 1 - p.

1 transitions to 0 with probability p(1-p) (if we add the first carry and we don't add the second one): S(1,0) = p(1-p).

1 transitions to 2 with probability  $p^2$  (if we add the first carry and also the second one):  $S(1,2) = p^2$ .

#### Transition graph of stochastic adding machine in base 2:



# Stochastic adding machine in base 2

Transition operator S of the stochastic adding machine in base 2: where  $s_{ij}$  is the probability of transition from state *i* to state *j*. Matrix S:

$$\begin{pmatrix} 1-p & p & 0 & 0 & 0 & 0 & 0 & 0 & \dots \\ p(1-p) & 1-p & p^2 & 0 & 0 & 0 & 0 & 0 & \dots \\ 0 & 0 & 1-p & p & 0 & 0 & 0 & 0 & \dots \\ p^2(1-p) & 0 & p(1-p) & 1-p & p^3 & 0 & 0 & 0 & \dots \\ 0 & 0 & 0 & 0 & 1-p & p & 0 & 0 & \dots \\ 0 & 0 & 0 & 0 & p(1-p) & 1-p & p^2 & 0 & \dots \\ 0 & 0 & 0 & 0 & 0 & 0 & 1-p & p & \dots \\ p^3(1-p) & 0 & 0 & 0 & p^2(1-p) & 0 & p(1-p) & 1-p & \dots \\ \vdots & \end{pmatrix}$$

#### Remarks:

- 1. S is bi-stochastic,  $S(X) \subset X$ , where  $X = I^{\infty}, I^{p}, p > 1$  or  $c_{0}$ ;
- 2. stochastic adding machine has application on psychology.

#### Stochastic adding machine in base 2

**Theorem**:(Killeen, Taylor)(2000) Acting in  $X = I^{\infty}(\mathbb{N})$ ,

$$\sigma(S) = \sigma_{pt}(S) = \{z \in \mathbb{C} : (f^n(z))_{n \ge 0} \text{ is bounded}\},\$$

where  $f(z) = z^2 + c$ ,  $c = \frac{p-1}{p^2}$ ,  $f^n = f \circ f \dots \circ f$ .

$$\sigma(S) = \{ z \in \mathbb{C} : S - zI : X \to X \text{ is not bijective} \},\$$

$$\sigma_{pt}(S) = \{ z \in \mathbb{C} : S - zI \text{ is not injective} \}.$$

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#### Stochastic adding machine in base 2





p = 0.6

p = 0.55

p = 0.51

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- What is the spectrum in other classical Banach spaces?
- What happens if we consider other bases instead of base 2?
- Is there exists a connection between topological properties of the spectrum, probabilistic properties of the Markov chain and also with dynamical properties of the transition matrix? Observe that if p ≥ 1/2, the spectrum is connected If 0

#### Fibonacci base

Constant base  $d \ge 2$ 

$$\sigma(S) = \sigma_{pt}(S) = \{z \in \mathbb{C} : (f^n(z))_{n \ge 0} \text{ is bounded}\},\$$

$$f(z) = z^d + c, \ c = \frac{p-1}{p^{d/d-1}}.$$

#### Fibonacci base

$$F_0 = 1, \ F_1 = 2, \ F_n = F_{n-1} + F_{n-2}, \ \forall n \geq 2.$$

$$N = \sum_{i=0}^{k(N)} \varepsilon_i(N) F_i = \varepsilon_{k(N)} \dots \varepsilon_{0(N)}$$

with  $\varepsilon_i(N) = 0, 1, \ \varepsilon_i \varepsilon_{i+1} \neq 11$  (greedy algorithm).

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Image: Image:

Exemple:  $F_n$ : 1, 2, 3, 5, 8, 13, 21,...

 $12 = 8 + 3 + 1 = F_0 + F_2 + F_4 = 10101$ 

Questions 1: How to define the adding machine?

Answer: by using automaton (Transducer): Classical

Questions 2: How to define the stochastic adding machine?

**Answer:** In collaboration with D. Smania, we defined a probabilistic transductor

**Automaton**: A finite automaton on A (finite alphabet): oriented graph set Au = (Q, A, E, I, T) where

- Q: finite set of states
- $I \subset Q$ : set of initial states
- $T \subset Q$ : set of terminal states.
- $E \subset Q imes A imes Q$  : set of edges.

The edges are labelled by the elements of A and

 $(p, a, q) \in E$  is also denoted by  $p \rightarrow q$ .

 $c = (p_0, a_0, p_1)(p_1, a_1, p_2) \dots (p_{n-1}, a_n, p_n)$  sequence of consecutive edges: finite path

**Transductor** : finite automaton T on  $A^* \times B^*$  labelled by elements of  $A^* \times B^*$ 

 $A^{\star}$  (resp.  $B^{\star}$ ) : set of finite words on A (resp. B).

An edge of T from a state p to a state q labelled by  $(u, v) \in A^* \times B^*$  will be denoted by  $p \to q$ . We say that T reads the word of input on  $A^*$  and transforms it to a word on  $B^*$ .

#### Trasductor

1. Transductor of base 2:



Example: 111 + 1 = 1000(0/1) (1/0) (1/0) (1/0)  $\rightarrow$  Output: 1000

2. Transductor of Fibonacci base:



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### Probabilistic Trasductor

1. Probabilistic Transductor of base 2:



**Example:** 111 + 1 = 1000 with probability  $p^4$ 

 $(0/0,1)(0/1,p) (1/0,p) (1/0,p) (1/0,p) \to 1000$  with probability  $p^4$ .

2. Probabilistic Transductor of Fibonacci base:



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#### Fibonacci base:

**Theorem:**(Messaoudi, Smania)(2010)  $\sigma_{pt}(S) \subset K \subset \sigma(S)$ 

$$\begin{split} & \mathcal{K} = \{ z \in \mathbb{C} : (f^n(z, z))_{n \geq 0} \text{ is bounded} \}, \\ & f : \mathbb{C}^2 \to \mathbb{C}^2, \ f(x, y) = (xy + c, x), \ \ c = \frac{p-1}{p^2}. \end{split}$$

**Remark:**  $K = \{z \in \mathbb{C} : (g_n(z))_{n \ge 0} \text{ is bounded}\}$ , where  $g_0(z) = z$ ,  $g_1(z) = z$ ,  $g_n(z) = g_{n-1}(z)g_{n-2}(z) + c$ ,  $n \ge 2$ .

**Conjecture**:  $K = \sigma(S)$ .

# Images of K



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## Images of K



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#### Idea of the proof

$$\lambda \in \sigma_{pt}(S)$$
 implies  $\exists v = (v_n)_{n \ge 0} \in I^{\infty}, \ Sv = \lambda v.$ 

We can prove that  $v_n = q_n v_0, \forall n \ge 1, q_n = q_n(p, \lambda).$ 

$$q_{F_n} = \frac{1}{p} q_{F_{n-1}} q_{F_{n-2}} - \frac{1-p}{p}, \ n \ge 2.$$

and if  $n = F_{n_1} + \cdots + F_{n_1}$ 

$$q_n = \prod_{i=1}^k q_{F_{n_i}}$$

 $\lambda \in \sigma_{pt}(S) \iff q_n \text{ bounded } \Longrightarrow (q_{F_n}, q_{F_{n-1}}) \text{ bounded }.$ 

Since  $(q_{F_n}, q_{F_{n-1}}) = g^n(q_{F_1}, q_{F_0}), g(x, y) = (\frac{1}{p}xy - \frac{1-p}{p}, x),$ we deduce that  $\sigma_{pt}(S) \subset K$ .

• Is 
$$\sigma_{pt}(S) = K$$
?  
This is equivalent to prove that

 $(q_{F_n})_{n\geq 0}$  bounded implies  $(q_n)_{n\geq 0}$  bounded .

**Conjecture**:  $K = \sigma(S)$ .

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• How about topological and dynamical properties of K?

$$K = \{z \in \mathbb{C} : (z, z) \in F(f)\},\$$

$$\begin{array}{l} F(f): \mbox{ Filled Julia set (in } \mathbb{C}^2) \mbox{ of } \\ f: \mathbb{C}^2 \to \mathbb{C}^2, \ f(x,y) = (xy+c,x), \ c = \frac{p-1}{p^2}. \end{array}$$

• How about topological and dynamical properties of F(f)?  $F(f) = \{(x, y) \in \mathbb{C}^2 : f^n(x, y) \text{ bounded } \}.$ 

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 $F(f) = \{(x, y) \in \mathbb{C}^2 : (f^n(x, y))_{n \ge 0} \text{ is bounded}\},\$ 

$$f: \mathbb{C}^2 \rightarrow \mathbb{C}^2, \ f(x,y) = (xy+c,x), \ c = rac{p-1}{p^2}.$$

Work of Julia sets in  $\mathbb{C}^n$ ,  $n \ge 2$  or in the projective spaces began very active in 80 with works of Hubbard, Fornaess-Sybony, Bedford, Smilie, Milnor among others.

Several results for Hénon diffeomorphisms (and its generalizations).

Our map is not bijective. It appeared in the literature in a work of Guedj. Few properties of F(f) are known.

# Properties of F(f)

If 
$$c = 0$$
,  
 $F(f) = \{(x, y) \in \mathbb{C}^2, |y| \le |x|^{-\beta}\},\$ 

 $\beta$ : Golden number.



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# Properties of F(f)

Proposition: (Bonnot, De Carvalho, Messaoudi):

- F(f) is a closed non compact subset of  $\mathbb{C}^2$ .
- If |c| < 1/4,  $\lambda(F(f)) = \infty$ .

We also show the existence of a trapping bidisk for F(f): in other words, a complex bidisk such that every point of F(f) eventually enters the bidisk under forward iteration.

**Conjecture**: For sufficiently large |c|,  $\lambda(F(f)) = 0$ .

**Theorem:**(Bonnot, De Carvalho, Messaoudi): If 0 < c < 1/4, then  $F(f) \cap \mathbb{R}^2$  is a finite union of stable manifolds.

$$F(f) \cap \mathbb{R}^2 = W^s(\alpha) \cup W^s(\theta) \cup W^s(p) \cup W^s(f(p)) \cup W^s(f^2(p))$$

 $\alpha$  = attractive fixed point,  $\theta$  = saddle fixed point, p = 3 periodic point

# Properties of F(f)

**Remark:** The last result is true in the case -1 < c < 0 (see the work of Danilo Caprio).

Image for c = 0.22:



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# Properties of $\overline{F(f)}$



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c) c = - 2,1

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 $K = \{z \in \mathbb{C} : (z, z) \in F(f)\}$  and  $K = \{z \in \mathbb{C} : (g_n(z))_{n \ge 0} \text{ is bounded}\}, \text{ where}$  $g_0(z) = z, g_1(z) = z^2, g_n(z) = g_{n-1}(z)g_{n-2}(z) + c, n \ge 2.$ 

Theorem:(A, B, M, S ):

- *K* is a non empty compact subset of  $\mathbb{C}$  such that  $\mathbb{C} \setminus K$  is connected.
- $\exists 0 < a < b$  such that if |c| > b, then K is disconnected. If |c| < a, then K is a quasi disk.

Questions: How about other topological Properties?

Can we have a finite number of connected components?

Big difficulty: K is not invariant by some map as classical Julia sets in  $\mathbb{C}$ .

## Properties of K

Slices of K where

(i) 
$$c = -0.72i$$
 and  $y = 0.17 + 0.2i$ ,

(ii) 
$$c = 0.2$$
 and  $y = 0.33$ ,

(iii) 
$$c = 0.33$$
 and  $y = 0.33$ .



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### Properties of K

**Idea of the proof:** We used. Technique of classical Julia sets, Riemann Hurwitz Formula.

K is a quasi disk: Hubbard technique in  $\mathbb{C}^2$ .



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 $\{c \in \mathbb{C} : (f^n(0,0))_{n \ge 0} \text{ is bounded}\}, f(x,y) = (xy + c, x):$ 



**Remark**: We can define a stochastic adding machine for a large class of recurrent sequences as done for Fibonacci base

For instance

If  $F_n = aF_{n-1} + bF_{n-2}$ ,  $a \ge b \ge 1$ . The spectrum is connected to the map

$$f: \mathbb{C}^2 \to \mathbb{C}^2, \ f(x,y) = (x^a y^b + c, x), \ c = \frac{p-1}{p^2}.$$

If  $F_n = a_1F_{n-1} + a_2F_{n-2} + \ldots + a_dF_{n-d}$ ,  $a_1 \ge a_2 \ge \ldots \ge a_d \ge 1$ ,

then the spectrum is connected to the map  $f: \mathbb{C}^d \longrightarrow \mathbb{C}^d$ ,

$$f(x_1,\ldots,x_d) = (x_1^{a_1} x_2^{a_2} \ldots x_d^{a_d} + c, x_1, x_2, \ldots, x_{d-1}).$$

Case of base d = 2, the Markov chain is irreducible, aperiodic and null recurrent for all 0 .

Case d = 2, where the probability change at each stage.

Let  $(p_i)_{i \ge 1}$ ,  $0 < p_i < 1$ .

0 transitions to 0 with probability  $1 - p_1$  if we don't add the carry  $S(0,0) = 1 - p_1$ .

0 transitions to 1 with probability  $p_1$  if we add the carry  $S(0,1)=p_1.$ 

1 transitions to 1 with probability  $1 - p_1$  if we don't add the carry  $S(1,1) = 1 - p_1$ .

1 transitions to 0 with probability  $p_1(1-p_2)$  if we add the first carry and we don't add the second one  $S(1,0) = p_1(1-p_2)$ .

1 transitions to 2 with probability  $p_1p_2$  if we add the first carry and also the second one  $S(1,2) = p_1p_2$ . **Theorem:**(M, V, S)(2013):

- The Markov chain is null recurrent if  $\prod_{i=1}^{\infty} p_i = 0$ , otherwise it is transient.
- Acting in  $X = I^{\infty}(\mathbb{N})$ ,

$$\sigma(S) = \sigma_{pt}(S) = \{z \in \mathbb{C} : (f_n \circ f_{n-1} \dots \circ f_1(z))_{n \ge 0} \text{ is bounded}\},$$
  
where  $f_i(z) = z^2 + c_i, \ c = \frac{p_i - 1}{p_i^2}.$ 

**Remark**: We obtain many results about the topology of  $\sigma(S)$  which can be connected, have finite or infinite number of connected components or to be a Cantor.

 $\exists 0 < a < 1$  such that if  $p_i \ge a$  for all *i*, then  $\sigma(S)$  is a quasi disk.

- Is there a difference of the spectrum between the null recurrent and transient cases?
- How about dynamical system associated to the transition matrix *S*?

Is there a difference of the spectrum between the null recurrent and transient cases?

**Answer:** If S is transient,  $\sigma(S)$  can be connected or have a finite number of connected components.

The point spectrum of S acting in  $l^1$  is empty.

If S is null recurrent,  $\sigma(S)$  can also have infinite connected components. The point spectrum of S acting in  $l^1$  is not empty.

How about dynamical system associated to the transition matrix S?

How about dynamical system associated to the transition matrix S?

Answer: If M is transient, there exists  $\lambda_0 > 1$  such that for all  $\lambda > \lambda_0$ , the dynamical system  $(\lambda S, l^1)$  is topologically transitive and Devaney chaotic.

This is false if M is null recurrent.

Bratteli diagrams are important objects in the theories of operator algebras and dynamical systems.

It was originally defined in 1972 by O. Bratteli for classification of C<sup>\*</sup>-algebras and was connected with symbolic dynamical systems.

A Bratteli diagram is an infinite directed graph (V, E) where the vertex set V and the edge set E can be partitioned into finite sets, i.e  $V = \bigcup_{k=0}^{\infty} V(k)$  and  $E = \bigcup_{k=1}^{\infty} E(k)$ , where  $\#V(k) < \infty$ and  $\#E(k) < \infty$  for every  $k \ge 0$ , such that there exist the source map  $s : E \longrightarrow V$  and the range map  $r : E \longrightarrow V$  such that srestricted to E(k) is an onto map from E(k) to V(k-1) and rrestricted to E(k) is an onto map from E(k) to V(k) for every  $k \ge 1$ . If  $\#V(k-1) = I_{k-1}$  and  $\#V(k) = I_k$ , then E(k) determines a  $I_k \times I_{k-1}$  incidence matrix M(k), where  $M(k)_{i,j}$  is the number of edges going from vertex j in V(k-1) to vertex i in V(k).

#### Definition

We say that (V, E) is a *simple Bratteli diagram* if for each nonnegative integer k, there exists an integer n > k such that the product  $M(n) \cdot M(n-1) \cdot \ldots \cdot M(k)$  have only non-zero entries.

#### Definition

An ordered Bratteli diagram  $(V, E, \geq)$  is a Bratteli diagram (V, E) together with a partial order  $\geq$  on E such that edges  $e, e' \in E$  are comparable if and only if r(e) = r(e'), in other words, we have a linear order on the set  $r^{-1}(\{v\})$  for each  $v \in V \setminus V(0)$ .

# Bratteli diagrams



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# The Vershik map $V_B$



 $X=(2,2,3,2,\epsilon_{5},\epsilon_{6},...)$ 

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# The Vershik map $V_B$



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x=(2,2,3,2,&,E\_a,...)

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 $x=(2,2,3,2,\epsilon_{g},\epsilon_{g},...)$ 

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x=(2,2,3,2,&,E\_a,...)

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 $x=(2,2,3,2,\epsilon_{g},\epsilon_{g},...)$ 

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Thus, by definition of stochastic Vershik map, we have a Markov chain which the states set is  $\mathbb{Z}_+$  and the transition operator  $S = (S_{m,n})$ , where

 $S_{m,n}$  is the probability that  $V_B^m(x_0)$  goes to  $V_B^n(x_0)$ .

Example: 
$$M = \begin{pmatrix} 2 & 1 \\ 3 & 1 \end{pmatrix}$$

#### Stochastic Bratteli diagrams



#### Theorem

(Caprio, Messaoudi, Valle) Let  $(p_i)_{i\geq 1}$  be a sequence of non-null probabilities such that  $\#\{i : p_i < 1\} = \infty$ . Thus, a) every stochastic Bratteli diagram associated to  $(p_i)_{i\geq 1}$  is an irreducible Markov chain; b) the Markov chain is transient if and only if  $\prod_{j=1}^{\infty} p_j > 0$ . c) if  $\prod_{j=1}^{\infty} p_j = 0$ , then the stochastic Bratteli diagram associated to  $(p_i)_{i\geq 1}$  is null recurrent.

#### Theorem

(C, M, V) Let  $B = (V, E, \ge)$  be a 2 × 2 stationary simple ordered Bratteli diagram with a = c = 1, b > 0 and d = 0. Then the BV stochastic adding machine associated to  $(p_j)_{j\ge 1}$  is positive recurrent if  $p_j$  decreases to zero sufficiently fast as  $j \to \infty$ .

#### Theorem

(C, M, V) Let S be the transition operator associated to the Bratteli diagram B and the stochastic Vershik map. Then, acting in  $I^{\infty}(\mathbb{N})$ , we have that the set of eigenvalues of S is

$$\sigma_{pt}(S) \subset$$

$$\mathcal{E} := \{ z \in \mathbb{C} : (g_n \circ \ldots \circ g_0(z, z))_{n \ge 0} \text{ is bounded} \} \subset \sigma(S),$$
where  $g_n : \mathbb{C}^2 \longrightarrow \mathbb{C}^2$  are polynomials defined by
$$g_n(x, y) = \left( \frac{1}{p_{n+1}} x^a y^b - \frac{1-p_{n+1}}{p_{n+1}}, \frac{1}{p_{n+1}} x^c y^d - \frac{1-p_{n+1}}{p_{n+1}} \right), \text{ for all } n \ge 1.$$

**Corollary:** If a + b = c + d, then

$$\sigma_{pt}(S) = \{\lambda \in \mathbb{C} : (f_n \circ \ldots \circ f_0(\lambda))_{n \ge 0} \text{ is bounded}\},\$$

where 
$$f_0(x) = \frac{x - (1 - p_1)}{p_1}$$
 and  $f_n(x) = \frac{1}{p_{n+1}} x^{a+b} - \left(\frac{1}{p_{n+1}} - 1\right)$ , for all  $n \ge 1$ .

 $\mathcal{E} = \{\lambda \in \mathbb{C} : (g_n \circ \ldots \circ g_0(\lambda, \lambda))_{n \ge 0} \text{ is bounded}\}$ 

Example: 
$$M = \begin{pmatrix} 2 & 1 \\ 3 & 1 \end{pmatrix}$$



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Example: 
$$M = \begin{pmatrix} 2 & 3 \\ 3 & 2 \end{pmatrix}$$



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 $\mathcal{E} = \{\lambda \in \mathbb{C} : (g_n \circ \ldots \circ g_0(\lambda, \lambda))_{n \ge 0} \text{ is bounded}\}$ 

Exemplo: 
$$M = \begin{pmatrix} 7 & 1 \\ 1 & 7 \end{pmatrix}$$



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# Topological properties of stochastic Bratteli diagrams

#### Theorem

Suppose that  $p_i = p$ , for all  $i \ge 1$ . If det M = ad - bc < 0 and  $bc > det M^2$ , then the set  $\mathcal{E}$  satisfies the following properties:

• 
$$\mathbb{C} \setminus \mathcal{E}$$
 is a connected set.

2) If 
$$p < \frac{1}{2}$$
, then  $\mathcal{E}$  is not connected.

# **OBRIGADO!**

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