Journal of Combinatorial Theory, Series B 86, 109–113 (2002) doi:10.1006/jctb.2002.2117

# Circular Colouring and Orientation of Graphs

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<u>Abstract:</u> This paper proves that if a graph **G** has an orientation **D** such that for each cycle **C** with **d**.  $|\mathbf{C}| (\mathbf{mod} \ \mathbf{k}) \in \{1, 2, ..., 2d - 1\}$  we have  $\frac{|\mathbf{C}|}{|\mathbf{C}^+|} \leq \frac{k}{d}$  and  $\frac{|\mathbf{C}|}{|\mathbf{C}^-|} \leq \frac{k}{d}$ , then **G** has a  $(\mathbf{k}, \mathbf{d})$ -colouring, and hence  $\chi_c(\mathbf{G}) \leq \frac{k}{d}$ .

# **Outline:**

- Basic Definitions
- > Some Results
- > Xuding Zhu's Result



- A graph G is a couple (V(G), E(G)) where V(G) is a finite, non empty set of elements, called vertices of G, and E(G) is set of pairs of vertices, called edges of G.
- A digraph *D* is a couple (V(D), E(D)) where V(D) is a finite nonempty set of elements, called vertices of *D*, and E(D) is subset of the set of ordered pairs of distinct elements of V(D), called arcs of *D*.
- For a digraph D = (V(D), E(D)), the **Underlying Graph** of D, denoted by G[D], is the undirected graph created using all of the vertices in V(D), and replacing all arcs in E(D) with undirected edges.
- An orientation of a graph G = (V(G), E(G)) is a digraph D = (V(D), E(D)) such that G[D] = G.

### **Basic Definitions**



For a cycle C of D with a chosen direction of traversal (each cycle has two different directions for traversal), let  $C^+$  be the set of positive edges of C (i.e., whose direction coincide with the direction of the traversal) and let  $C^-$  be the set of negative edges of C.

The parameter  $\tau(C) = \max\{\frac{|C|}{|C^+|}; \frac{|C|}{|C^-|}\}\$  measure the "**imbalance**" of the cycle *C*. Let  $\xi(D) = \max\{\tau(C): C \text{ is a cycle of } D\}$ . Then  $\xi(D)$  is a measure of the imbalance of the orientation *D* of *G*.

#### **Basic Definitions**

For any integer *n*;  $n \pmod{k}$  denotes the unique integer n' such that  $0 \le n' \le k - 1$  and  $n \equiv n' \pmod{k}$ .

Suppose G = (V, E) is a graph and  $1 \le 2d \le k$  are integers. A (k, d)-coloring of G is a mapping  $f : V \to \{0, 1, \dots, k - 1\}$  such that for every edge xy of G:

$$d \le |f(x) - f(y)| \le k - d$$

The circular chromatic number  $\chi_c(G)$  of G is defined as:

$$\chi_c(G) = \min\{\frac{k}{d} : \text{there exists a } (k, d) \text{-coloring of } G\}.$$

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### **Some Results**

> Theorem 1: (Minty [2]) For any finite graph G,

$$\chi(G) = \min\{ [\xi(D)] : D \text{ is an orientation of } G \}$$

The non-trivial direction of Minty's result asserts that if G has an orientation D with  $\xi(D) \le k$  where k is an integer then G is k-colorable.

Tuza's result says that to obtain the same conclusion, instead of requiring  $\xi(D) \le k$  which is equivalent to require that  $\tau(C) \le k$  for every cycle *C*, it suffices to require that  $\tau(C) \le k$  for those cycles *C* with |C|(mod k) = 1.

Goddyn, Tarsi and Zhang's result says that if  $\xi(D) \leq \frac{k}{d}$  for some fraction  $\frac{k}{d}$  then G is (k, d)-colorable.

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# Xuding Zhu's Result

Xuding Zhu in his paper proves that if a graph G has an orientation D such that for each cycle C with d.  $|C| \pmod{k} \in \{1, 2, ..., 2d - 1\}$  we have  $\frac{|C|}{|C^+|} \le \frac{k}{d}$  and  $\frac{|C|}{|C^-|} \le \frac{k}{d}$ , then G has a (k, d)-colouring, and hence  $\chi_c(G) \le \frac{k}{d}$ .

This result is a generalization of Tuza's result (if d = 1) and is an improvement of Goddyn et al's result.

# Proof

The proof of this result is parallel to Tuza's proof of that special case. Let r be a fixed vertex of D. Each spanning tree T of D is considered as rooted at r. Given such a spanning tree T, we define the weight  $w_T(x)$  of a vertex x of D(with respect to T) recursively as follows:

- $w_T(r) = 0;$
- If xy is an edge of T oriented from x to y and  $w_T(x)$  has already been defined, then  $w_T(y) = w_T(x) k + d$ ;
- If xy is an edge of T oriented from x to y and  $w_T(y)$  has already been defined, then  $w_T(x) = w_T(y) + d$ ;

Since T is a spanning tree, for each vertex x of D;  $w_T(x)$  is uniquely defined. Then we define the weight w(T) of T as

$$w(T) = \sum_{x \in V(D)} W_T(x)$$

Choose a rooted spanning tree T of D which has the maximum weight. Let  $f: V \to \{0,1,\ldots,k-1\}$  be defined as  $f(x) = w_T(x) \mod k$ : We shall show that f is a (k,d)-colouring of G.

Let  $xy \in E(G)$ . Without loss of generality we assume that the edge xy is oriented from x to y in D. We will study three different cases:

- If x is not on the y-r-path of T and y is not on the x-r-path of T;
- If y is on the x-r-path of T;
- If x is on the y-r-path of T.





If  $w_T(x) - w_T(y) < d$ ; then we delete the edge of T connecting x to its father, and add the edge xy



$$w_{T'}(x) = w_{T'}(y) + d = w_T(y) + d > w_T(x)$$

Then we obtain a spanning tree T' for which  $w_{T'}(v) \ge w_T(v)$ for each v; and  $w_{T'}(v) > w_T(v)$  for every descendents of x; including x itself. So w(T') > w(T), contrary to our choice of T. Thus  $w_T(x) - w_T(y) \ge d$ .

Similarly we can prove taht  $w_T(x) - w_T(y) \le k - d$ .

Those,  $d \leq w_T(x) - w_T(y) \leq k - d$ .





Those,  $d \leq w_T(x) - w_T(y) \leq k - d$ .

which implies that  $d \leq |f(x) - f(y)| \leq k - d$  (as  $f(x) - f(y) = w_T(x) - w_T(y) \mod k$ ), a contradiction.



Those, 
$$d \le w_T(x) - w_T(y) \le k - d$$
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Then ,  $d \le w_T(x) - w_T(y) (modk) \le k - d$ , and so:

$$d \leq |f(x) - f(y)| \leq k - d$$

As  $f(x) - f(y) = w_T(x) - w_T(y) \mod k$ .



• If *y* is on the *x*-*r*-path of *T*:

Assume  $w_T(x) - w_T(y) = ak + j$  for some integers a; j such that  $0 \le j \le k - 1$ : Then, f(x) - f(y) = j, And hence  $j \in \{0, ..., d - 1\} \cup \{k - d + 1, ..., k - 1\}$ .





Let p be the number of edges on the x-y-path of T oriented toward the root, and let q be the number of edges on this path oriented away from the root. By the definition of the weight, we know that  $w_T(x) - w_T(y) = pd - q(k - d)$ . Therefore, we have (pd + qd)(mod k) = j; and (p + q + 1).  $d(mod k) = j + d(mod k) \in \{1, 2, ..., 2d - 1\}$  note that the cycle C consisting the x-y -path and the edge xy is a cycle of length |C| = p + q + 1: Hence  $d|C|(modk) \in \{1, 2, ..., 2d - 1\}$ . By our assumption,  $\frac{|C|}{|C^-|} \leq \frac{k}{d}$ ; which implies  $\frac{|C^+|}{|C^-|} \leq \frac{(k-d)}{d}$  here we choose the direction of traversal of C so that those edge on the x-y -path oriented towards r belongs to  $C^+$ : Hence  $|C^+| = p$  and  $|C^-| = q + 1$ ; which implies that  $pd - q(k - d) \leq k - d$ . Therefore,  $d \leq w_T(x) - w_T(y) \leq k - d$ , a contradiction.





#### Similarly to the Case 3.

Thank you