

EQUATIONS: a dependent pattern-matching compiler

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- ▶ EPIGRAM/AGDA-style pattern-matching definitions with `with`
- ▶ Purely logical handling of recursion for inductive families
- ▶ Propositional equations for definitional equalities
- ▶ Elimination principle and support for applying it

Entirely elaborated to the vanilla kernel!

DEMO

Dependent Pattern-Matching

- ▶ Patterns = **well-typed** refinements of the signature
- ▶ We refine the **entire** context at each node (correct dependency tracking)
- ▶ Internalizes “Uniqueness of Identity Proofs” (axiom K)
- ▶ **Inaccessible** patterns + first-match semantics ensure operationality
- ▶ Empty nodes ensure decidability of coverage

Elaboration into CIC + K

Three phases:

- 1 *Generation* of a splitting tree from the clauses
- 2 *Translation* from the splitting tree to COQ terms with holes
- 3 *Proofs* of the obligations using a mix of ML and \mathcal{L}_{tac} code

term, type	t, τ	$::=$	$x \mid \lambda x : \tau, t \mid \Pi x : \tau, \tau' \mid \dots$
binding	d	$::=$	$(x : \tau) \mid (x := t : \tau)$
context	Γ, Δ	$::=$	\vec{d}
program	$prog$	$::=$	$f \Gamma : \tau := \vec{c}$
user clause	c	$::=$	$f \vec{u} \vec{p} n$
user pattern	up	$::=$	$x \mid \mathbf{C} \vec{u} \vec{p} \mid ?(t)$
user node	n	$::=$	$:= t \mid :=! x \mid \mathbf{with} t := \{ \vec{c} \}$

Searching for a splitting tree

For $f \Delta : \tau$ we define $f_{\text{comp}} \Delta := \tau$, so $f : \Pi \Delta, f_{\text{comp}} \bar{\Delta}$.

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pattern	p	$::=$	$x \mid \mathbf{C} \overrightarrow{p} \mid ?(t)$
context map	c	$::=$	$\Delta \vdash \overrightarrow{p} : \Gamma$
splitting	spl	$::=$	$\text{Split}(c, x, (spl?)^n) \mid \text{Compute}(c, rhs)$
node	rhs	$::=$	$\text{Program}(t) \mid \text{Refine}(c, t, \ell, spl)$

Goal Starting with $f \Delta : f_{\text{comp}} \overline{\Delta} := \overrightarrow{p} \dots$, find a covering of the context map $\text{idsubst}(\Delta) = \Delta \vdash \overline{\Delta} : \Delta$.

Proof search example

Overlapping clauses with first-match semantics.

```
Equations equal (n m : nat) : { n = m } + { n ≠ m } :=  
equal O O := left eq_refl ;  
equal (S n) (S m) with equal n m := {  
  equal (S n) (S ?(n)) (left eq_refl) := left eq_refl ;  
  equal (S n) (S m) (right p) := right _ } ;  
equal x y := right _.
```

```
Split(n m : nat ⊢ n m : n m : nat, n, [  
  Split(m : nat ⊢ O m : n m : nat, m, [  
    Compute(⊢ O O : n m : nat, Program(left eq_refl)),  
    Compute(m : nat ⊢ O (S m) : n m : nat, Program(right _))]),  
  Split(n m : nat ⊢ (S n) m : n m : nat, m, [  
    Compute(n : nat ⊢ (S n) O : n m : nat, ...),  
    Compute(n m : nat ⊢ (S n) (S m) : n m : nat,  
      Refine(equal n m,  
        idsubst(n m : nat, x : {n = m} + {n ≠ m}), ℓ, ...)))]])])
```


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- ▶ $\text{Program}(t)$: witnessed by the term.
- ▶ $\text{Refine}(t, c, \ell, s)$: witnessed by inserting a let-definition in the context, strengthening, abstracting and clearing its body, then applying the compiled term for label ℓ .

With nodes in detail

Consider a current problem $\Delta \vdash \vec{p} : \Gamma$ and a user clause $f \vec{u}\vec{p}$ **with** $t_{pre} := \{ e \}$ matching it. We typecheck t_{pre} into $t : \tau$ and use strengthening and abstraction to find a new context

$$\Delta^t, x_t : \tau, \Delta_t[t/x_t] \text{ such that } \Delta^t, \Delta_t \sim \Delta$$

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Using the clauses e we then build a subcovering s of the identity context map

$$c = \text{idsubst}(\Delta^t, x_t : \tau_\Delta, \Delta_t[t/x_t])$$

and return $\text{Refine}(t, c, l.n, s)$.

With nodes in detail

Consider a current problem $\Delta \vdash \overrightarrow{p} : \Gamma$ and a user clause $f \overrightarrow{up}$ with $t_{pre} := \{ e \}$ matching it. We typecheck t_{pre} into $t : \tau$ and use strengthening and abstraction to find a new context

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Compilation produces $\ell.n : \Pi \Delta^t (x_t : \tau_\Delta) \Delta_t[t/x_t], (f_{\text{comp}} \overrightarrow{p})[t/x_t]$, we build

$$(\lambda \Delta, \ell.n \overline{\Delta^t} t \overline{\Delta_t}) : \Pi \Delta, f_{\text{comp}} \overrightarrow{p}$$

1 Dependent pattern-matching compilation

2 Recursion

- The Below way
- Subterm relations

3 Reasoning support

- Equations
- Elimination principle
- Eliminating calls

- ▶ Syntactic guardness checks are too fragile (and buggy)
- ▶ Do not work well with abstraction/modularity
- ▶ Restricted to structural recursion on a single argument, with no currying allowed

Use the logic instead!

The Below way (McBride and McKinna)

```
Fixpoint Below_nat (P : nat → Type) (n : nat) : Type :=  
  match n with  
  | 0 ⇒ ()  
  | S n' ⇒ (P n' × Below_nat P n')  
end%type.
```

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Fixpoint below_nat (P : nat → Type)  
  (step :  $\prod n : \text{nat}, \text{Below\_nat } P n \rightarrow P n$ )  
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```

```
Definition rec_nat (P : nat → Type)  
  (step : Π n : nat, Below_nat P n → P n)  
  (n : nat) : P n := step n (below_nat P step n).
```

```
Equations unzip {A B n} (v : vector (A×B) n) : vector A n × vector B n :=
unzip A B n v by rec v :=
unzip A B ?(O) Vnil := (Vnil, Vnil) ;
unzip A B ?(S n) (Vcons (pair x y) n v) with unzip v := {
  | (pair xs ys) := (Vcons x xs, Vcons y ys) }.
```

- ▶ **by rec v** applies the elimination principle associated to the type of v (found using typeclass resolution).

Equations **unzip** $\{A\ B\ n\}$ $(v : \text{vector } (A \times B)\ n) : \text{vector } A\ n \times \text{vector } B\ n :=$
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- ▶ Each recursive occurrence of `f` is transformed to a trivial projection `fcomp_proj` : $\Pi\ \Delta\ \{p : f_{\text{comp}}\ \overline{\Delta}\}, f_{\text{comp}}\ \overline{\Delta}$.
- ▶ Proof search for `fcomp` goals appearing as obligations, unfolding `Below` hypotheses.

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- ▶ Extracts efficiently, but proof search a bit more complicated than [Below](#).

Subterm relation example: vectors

Derive Subterm for vector.

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Inductive vector_strict_subterm ($A : \text{Type}$)

: $\forall n m : \text{nat}, \text{vector } A n \rightarrow \text{vector } A m \rightarrow \text{Prop} :=$

vector_strict_subterm_1_1 : $\forall (a : A) (n : \text{nat}) (v : \text{vector } A n),$
vector_strict_subterm $A n (\text{S } n) v (\text{Vcons } a v).$

Check vector_subterm : $\forall A : \text{Type}, \text{relation } \{index : \text{nat} \ \& \ \text{vector } A \ index\}.$

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- ▶ Equations hold definitionally in $CCI + K$
- ▶ Equations for **with** nodes are just proxies to the helper function *f.l.*
- ▶ All put together in a rewrite database, *f* can now be opacified.
- ▶ For well-founded definitions, we use the unfolding lemma to prove the equations.

Elimination principle

```
Equations filter {A} (l : list A) (p : A → bool) : list A :=  
filter A nil p := nil ;  
filter A (cons a l) p with p a := {  
  | true := a :: filter l p ; | false := filter l p }.
```

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```

Check (filter_elim :

```
∀ (P : ∀ (A : Type) (l : list A) (p : A → bool), filter_comp l p → Prop),
```

```
let P0 := fun (A : Type) (a : A) (l : list A) (p : A → bool)
```

```
  (refine : bool) (res : filter_comp (a :: l) p) ⇒
```

```
  p a = refine → P A (a :: l) p res
```

in

```
(∀ (A : Type) (p : A → bool), P A [] p []) →
```

```
(∀ (A : Type) (a : A) (l : list A) (p : A → bool),
```

```
  P A l p (filter l p) → P0 A a l p true (a :: filter l p)) →
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The elimination principle can only be applied usefully to calls with solely variable arguments.

$$\prod A (l : \text{list } A), \text{app } l [] = l$$

Eliminating calls

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Use dependent elimination/abstraction by equalities again:

$$\begin{aligned} & (\lambda (l \ l' : \text{list } A) (r : \text{app}_{\text{comp}} l \ l'), \\ & \quad l' = [] \rightarrow r = \text{app } l [] \rightarrow \text{app } l [] = l) \\ & \quad l [] (\text{app } l []) \end{aligned}$$

Directly apply the elimination principle and simplify the equations.

A function definition package handling:

- ▶ Full, nested dependent pattern-matching
- ▶ Structural and well-founded recursion on inductive families
- ▶ Generation of useful support lemmas for reasoning a posteriori

Compared to `FUNCTION`, mainly adds support for inductive families and a more robust implementation.

Tested on a bit-fiddling library: less boilerplate, shorter proofs.

- ▶ Treatment of non-constructor indices and unsolved constraints, e.g.: $0 = x + y$, with a subsequent splitting on x .
- ▶ Mutual recursion (structural and well-founded)
- ▶ Move to `eq_dep` instead of `JMeq`? Necessary to use decidable instances of `K`.
- ▶ Efficiency, primitive handling of `K`.



<http://mattam.org/research/coq/equations.en.html>